

Lecture 7: Spectral Graph Theory I

①

[SCRIBE] G : undir graph $\{ \text{self-loops, mult edges OK} \}$ $\{ \text{think regular to worm up} \}$
 Recall: $\{f: V \rightarrow \mathbb{R}\}$ dim. $n = |V|$ vec. space, w/ inner prod

$$\langle f, g \rangle := \sum_{u \in V} [f(u)g(u)].$$

π : invariant/stationary distrib. \rightarrow choose ^{unif.} rand edge $u \sim v$ ($= v \sim u$)
 • output u .
 • π_u prop. to $\deg(u)$: $\pi_u = \frac{\deg(u)}{2n}$. Unif. dist. iff G regular.

For $u \sim \pi$, $u_0 \rightarrow u_1 \rightarrow u_2 \rightarrow \dots \rightarrow u_t$ standard rand walk
 $\uparrow \quad \nearrow$
 all $\sim \pi$.

$$\mathcal{E}[f] := \frac{1}{2} \sum_{u \sim v} [(f(u) - f(v))^2] \quad , \quad \boxed{\mathcal{E}[1_S] = \Pr_{u \sim \pi}[u \in S, v \notin S]}.$$

Dirichlet form

local variance

$$\text{Var}[f] = \frac{1}{2} \sum_{\substack{u \sim v \\ \text{indep.}}} [(f(u) - f(v))^2]$$

✓ 1 on S ,
0 off S ,

min $\mathcal{E}[f]$?

→ If G 's conn. comps. are S_1, \dots, S_e , then $1_{S_1}, \dots, 1_{S_e}$ are lin. indep., their span is all f s.t. $\mathcal{E}[f] = 0$.

max $\mathcal{E}[f]$? $\{ \text{Not scale inv., } \because \mathcal{E}[cf] = c^2 \mathcal{E}[f] \}$, so we look at $\|f\|_2^2 = \mathcal{E}[f^2]$

$$\max \mathcal{E}[f] \quad \text{s.t. } \|f\|_2^2 = \mathcal{E}[f^2] \leq 1$$

$$\left(\text{equiv.: } \begin{array}{l} \mathcal{E}[f^2] = I, \\ \text{Var}[f] \stackrel{(\text{def.})}{\leq} 1 \end{array} \right)$$

fact: $\mathcal{E}[f] \leq 2 \|f\|_2^2$ w/ equality poss. iff G is bipartite.

(2)

Method of maximizing $E[f]$ with thinking is actually somewhat more interesting.
 stat: (Will look at max/mining $E[f]$) some more. Have to start by studying it.

$$\text{pf: } E[f] = \text{fing arithmetic} = \|f\|_2^2 - E_{u \sim \pi} [f(u) f(v)] \quad \text{Leave up.}$$

$\|f\|_2^2 \leq \|f\|_2^2$ by C.S.

{an interesting qty; let's study it}

$$= E_{u \sim \pi} [f(u) \cdot E_{v \sim u} [f(v)]]$$

\uparrow
avg val of f on nbrs
of u

$u \downarrow$
 \uparrow a nbr
on v

avg val of f on nbrs

{a number}

def:

$$Kf : V \rightarrow \mathbb{R}, Kf(u) = E_{v \sim u} [f(v)].$$

$(Kf)(u)$

Markov / transition / normalized adj. operator {think heat diffusion
(w/ self-loop)}
(w/ self-loop)

rem: K is linear operator {in vec space sense}, $K(f+g) = Kf + Kg$

\downarrow corresponds to matrix {why? Linearity of expectation}

$$u \rightarrow \begin{bmatrix} & & \\ & K & \\ & & \end{bmatrix} \begin{bmatrix} | \\ s \\ | \end{bmatrix} \xleftarrow{\text{adj.}} \begin{bmatrix} | \\ Kf \\ | \end{bmatrix} \xleftarrow{u} u.$$

$$K[u,v] = \begin{cases} \frac{1}{\deg(u)} & \text{if } (u,v) \in E \\ 0 & \text{else} \end{cases} = \Pr[u \rightarrow v | @u]$$

$\therefore K$ is adj. mtx, normalized so row sums 1 {stoch. mtx}.

$$= \frac{1}{d} A \quad \text{if } G \text{ d-regular.} \rightarrow \text{in which case } K \text{ is symm} \quad \text{if } K = K^T$$

③

fact: For $f, g: \mathbb{N} \rightarrow \mathbb{R}$,

$$\langle f, Kg \rangle = \underset{u \sim \pi}{\mathbb{E}} [f(u)g(u)] \quad \begin{aligned} &= \mathbb{E}_{u \sim \pi} [f(u) \cdot (Kg)(u)] \\ &= \mathbb{E}_{u \sim \pi} [f(u) \mathbb{E}_{v \sim u} [g(v)]] \\ &= \underset{u \sim \pi}{\mathbb{E}} [f(u)g(u)] \end{aligned}$$

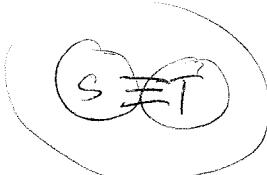
[almost think of this as def¹]

$$\langle Kf, g \rangle$$

" K is self-adjoint"

[this is the equiv. of " K is symmetric" in the G -is regular case.]

e.g.:



[it don't have to be disj. actually]

Let $S, T \subseteq V$, $f = \mathbf{1}_S$, $g = \mathbf{1}_T$.

$$\langle f, Kg \rangle = \underset{u \sim \pi}{\mathbb{E}} [\mathbf{1}_S(u) \mathbf{1}_T(u)]$$

$$(\Rightarrow \langle g, Kf \rangle) = \Pr_{u \sim \pi} [u \in S, v \in T]$$

[prob. you go from S to T in one step.]

~~Random walk~~: Aside: Let π_p be a prob. distr. on V :

$\pi_p = [p_v]$
[think a row vector]
rho, haha

• Draw $u \sim \pi_p$

• Do 1 step rand walk $\rightarrow v$.

Let π'_p be prob. dist. of v .

$$\text{Ex: } \pi'_p = \begin{bmatrix} p_v \\ \vdots \\ p_v \end{bmatrix} \quad \text{[trivial]}$$

\therefore [in particular] $\pi' K = \pi$.

[K is a mtx: can multiply mtxs / compose operators]

Q: What is $K^2 = K \circ K$

$$A: \langle K^2 f \rangle(u) = \underset{\substack{w \sim \pi \\ \text{2 steps}}} {\mathbb{E}} [f(w)]$$

$$\text{PF: } (K(Kf))(u) \triangleq \underset{(Kf)}{\underset{v \sim u}{E}} [f(v)] = \underset{v \sim u}{E} \left[\underset{w \sim v}{E} [f(w)] \right] \left(\underset{w \sim v}{E} [f(w)] \right) \quad (4)$$

con: $\forall t \in \mathbb{N}, \quad K^t f(u) = \underset{u \xrightarrow{+steps} w}{E} [f(w)]$
 { mtx power or t-fold composition }

$$t=0? \quad \checkmark \quad I f(u) = \underset{w=u}{E} [f(w)] = f(u). \quad \checkmark$$

i.e. id. mtx

Back to $E[f] = \frac{1}{2} \underset{u \sim v}{E} [(f(u) - f(v))^2]$ {as we saw by simple arith.}

$$\begin{aligned} &= \langle f, f \rangle - \underset{u \sim v}{E} [f(u) f(v)] \rightarrow \langle f, Kf \rangle \\ &= \langle f, f - Kf \rangle = \langle f, \underbrace{(I - K)f}_{L} \rangle \quad \text{{ah, the Laplacian!}} \end{aligned}$$

def: L : (normalized) Laplacian operator mtx is $L = I - K$.

i.e., $Lf: V \rightarrow \mathbb{R}$, $Lf(u) = f(u) - \underset{v \sim u}{E} [f(v)]$. {Aside: elsewhere you'll see all this tortured stuff about $D^{-1/2}$, D is diag(bj)}

Q: What is the "meaning" of L ?

A: $\langle f, Lf \rangle = E[f]$ awesomely meaningful. $= \frac{1}{2} \underset{u \sim v}{E} [(f(u) - f(v))^2]$ use $\langle \cdot, \pi \cdot \rangle \approx 1$

[not] $\langle f, Lg \rangle = E[f, g]$

Let $S \subseteq V$, let $f = 1_S$. $\langle f, Lf \rangle = E[f] = \Pr_{u \sim v} [u \in S, v \notin S] \quad (= \frac{1}{2} \text{"vol"}(\partial S))$

$$\langle f, f \rangle = E[f^2] = \Pr_{u \sim v} [u \in S]$$

$u \sim v$
 ∴ ratio is $\Pr_{u \sim v} [v \notin S | u \in S]$

$= \Pr$ [pick rand $u \sim S$, do 1 step,
 ↗ that you go outside S]
] choose w. prob. prop'l to $\deg(u)$

def: $\Phi(S) = \frac{\text{Conductance of } S}{\text{Vol}(S)} \in [0, 1]$ "escape prob. of S "

Sparset Cut: Given G , find S with $\text{Vol}(S) \leq \frac{1}{2}$ minimizing $\Phi(S)$
 [NP-hard, majorly important for divide & conquer algs on graphs
 approximability wide open.]

[Though we're somewhat more interested in minimizing $E[f]$, let's think about maximizing (cf. Max-Cut), as eigenvalues will finally enter the pic.]

$$\max E[f] = \langle f, Lf \rangle = \frac{1}{2} \sum_{uv} E[(f(u) - f(v))^2]$$

$$\text{s.t. } \left\| f \right\|_2^2 = \langle f, f \rangle = 1. \quad \begin{array}{l} \text{← compact set } \{ \text{ellipsoid in } \mathbb{R}^n \} \\ \text{continuous fn } (f: \mathbb{R}^n \rightarrow \mathbb{R}) \end{array}$$

\therefore a maximizer $\phi: V \rightarrow \mathbb{R}$ exists. [also, roughly speaking, efficiently findable]

Claim: $L\phi = \lambda\phi$ for some $\lambda \in \mathbb{R} \iff L\phi$ parallel to ϕ

[assuming that:] cor: $E[\phi] = \langle \phi, \lambda\phi \rangle = \lambda \|\phi\|_2^2 = \lambda \stackrel{(>0)}{\leq} 2.$

fact: $E[\phi] = 0$. Pf: maximizer will have $\text{Var}[\phi] = 1$ [replace ϕ with ϕ/μ , rescale].

Proof of claim: [Lagrange mults] Suppose

$L\phi$ (not parallel)
 $\exists \lambda$

Let γ be unit vector in this dir. [exists, \because not \parallel]

let $\varepsilon \neq 0$ be small. Let $f = \phi + \varepsilon\gamma$. $\|f\|_2^2 = 1 + \varepsilon^2$ (Pythagorus)

$$\begin{aligned} \langle f, Lf \rangle &= \langle \phi + \varepsilon\gamma, L\phi + \varepsilon L\gamma \rangle \\ &= \langle \phi, L\phi \rangle + 2\varepsilon \langle \gamma, L\phi \rangle + O(\varepsilon^2) \end{aligned}$$

[used self-adj.]

$$\frac{\langle f, Lf \rangle}{\langle f, f \rangle} = \langle \phi, L\phi \rangle + 2\varepsilon \langle \phi, L\phi \rangle + O(\varepsilon^2)$$

\uparrow not 0 [see pic]

$> \langle \phi, L\phi \rangle$ for small enough $\pm \varepsilon$, $\Rightarrow \Leftarrow$. \square

Great, so ϕ is a maximizer. Let's disallow it now — restrict to subspace orthog.

Now consider $\max \langle f, Lf \rangle$
 s.t. $\langle f, f \rangle = 1$, $f \perp \phi$ $\left(\begin{matrix} \phi \perp f \\ \text{compact set} \end{matrix} \right)$

Repeat: Gives ϕ' with $L\phi' = \lambda'\phi'$, $E[\phi'] = 0 \Leftrightarrow \langle \phi', 1 \rangle = 0$
 $\phi' \perp \phi$. $\lambda' \leq \lambda$.

Now do s.t. $f \perp \phi, \phi'$. repeat etc. (All will be \perp to 1; last one will be 1 then.)

thm: Given G , \exists orthonormal basis $\phi_0, \phi_1, \dots, \phi_{n-1} : V \rightarrow \mathbb{R}$ ("eigenfunctions")
 $\langle \phi_i, \phi_j \rangle = \begin{cases} 1 & \text{if } i=j \\ 0 & \text{else} \end{cases}$ $\| \phi \|$ form basis ϕ_0, \dots, ϕ_n (Change w.r.t.
 ϕ_0, \dots, ϕ_n)
 & reals ("eigenvalues") $0 = \lambda_0 \leq \lambda_1 \leq \dots \leq \lambda_{n-1} \leq 2$
 s.t. $L\phi_i = \lambda_i \phi_i$. $\phi_0, \dots, \phi_{n-1}$ our first 1, our first ϕ

rem/
 ex: λ is uniquely def'd but ϕ 's not uniquely def'd when \exists "multiplicity";
 e.g., $\lambda_5 = \lambda_6 = 2$ ϕ_5, ϕ_6
 Even w/o mult.: $\pm \phi_i$. any rotation
 [I will always be one, we'll stick w/ it.]

What is the importance of this?

The basis $\{\phi_0, \dots, \phi_{n-1}\}$ makes understanding L, K, E easy!!

def: Given $\phi_0, \dots, \phi_{n-1}$ (o.n. basis),

can write any $f: V \rightarrow \mathbb{R}$ uniquely as lin comb.

$$f = \sum_{i=0}^n \hat{f}(i) \phi_i + \text{real coeffs} = \sum_{i=0}^n \hat{f}(i) \phi_i$$

for real consts $\hat{f}(0), \dots, \hat{f}(n-1) \in \mathbb{R}$.

$\{\phi_i\}$ are like world's most interesting funcs \mathbb{I} .

$$\text{prop: } Lf = L\left(\sum_{i=0}^n \hat{f}(i) \phi_i\right) = \sum_{i=0}^n \hat{f}(i) L\phi_i = \sum_{i=0}^n \lambda_i \hat{f}(i) \phi_i$$

$\hat{L}\hat{f}(i)$: L multiplies $\hat{f}(i)$ by λ_i .

$$\text{prop: } \langle f, g \rangle = \sum_{i=0}^n \hat{f}(i) \hat{g}(i). \quad \text{[dot prod formula]}$$

$$\text{pf: } \langle \sum_i \hat{f}(i) \phi_i, \sum_j \hat{g}(j) \phi_j \rangle = \sum_{ij} \hat{f}(i) \hat{g}(j) \langle \phi_i, \phi_j \rangle \text{ : orthog.} \quad (1)$$

$$\text{cor: } \|f\|_2^2 = \langle ff \rangle = \sum_{i=0}^n \hat{f}(i)^2. \quad E[f] = \langle f, 1 \rangle = \langle f, \phi_0 \rangle = \hat{f}(0). \quad (\because \text{Var}[f] = \sum_{i=1}^n \hat{f}(i)^2)$$

$$E[f] = \langle f, Lf \rangle = \sum_{i=0}^n \lambda_i \hat{f}(i)^2 \quad \text{e.g. } 0 \cdot \hat{f}(0)^2 + 1 \cdot \hat{f}(1)^2 + 2 \cdot \hat{f}(2)^2 + 9 \cdot \hat{f}(3)^2 + \dots + 1.8 \cdot \hat{f}(n-2)^2 + 2 \cdot \hat{f}(n)^2$$

\hookrightarrow bipartite

$$Q: \text{What is } \min_{\text{over } f} \frac{E[f]}{\text{Var}[f]} \text{ s.t. } \text{Var}[f] \neq 0 \quad \text{Var}[f] \neq 0 \quad \text{f not const}$$

A: λ_1 : follows immmed. from (1), (2)

achieved by ϕ_1

"Poincaré Ineq": $\lambda_1 \text{Var}[f] \leq E[f]$.

If $\phi_1: V \rightarrow \mathbb{R}$ had range $\{0, 1\}$, then $E[\phi] = \text{edge bdy size}$

$$\text{Var}[\phi] = E[\phi^2] - E[\phi]^2 = \text{vol}(S)(1 - \text{vol}(S)). \\ = \text{vol}(S)\text{vol}(S^c)$$

$$\frac{E[f]}{\text{Var}[f]} = \frac{E[1_S]}{\text{Var}[1_S]} = \frac{\text{Var}[1_S]}{\|1_S\|_2^2} \approx \Phi(S) \text{ or } \Phi(S^c)$$

Would get S with $\Phi(S) = O(1)$. Const factor approx to sparsity. What if not? for whichever has $\text{vol} \leq \frac{1}{2}$. Cheeger...

Random walks