

Lecture 3: (Chernoff/Tail/Large) Deviation Bounds

Massive naming confusion. CS people call them "Chernoff Bds" after a '52 paper. Or sometimes Hoeffding, who did slightly strongerish in '40s. Though Bernstein '20s was even strongerish. Let $H = \# \text{heads in } n \text{ fair flips}$.

rec: Berry-Esseen $\Rightarrow \Pr[H \geq \frac{n}{2} + \frac{\sqrt{n}}{2} \cdot t] \approx \Phi(\frac{t}{\sqrt{2}})$

$$\sim \frac{\phi(t)}{\sqrt{2}} \leq \phi(t) \quad [t \geq 1]$$

$$\text{Let } t = 10\sqrt{\ln n}, \quad = e^{-50\ln n} = \frac{1}{n^{50}} \rightarrow \leq e^{-t^2/2}.$$

$$\text{So } \Pr[H \geq \frac{n}{2} + 5\sqrt{n \ln n}] \leq \frac{1}{n^{50}} ? \quad \boxed{\text{Wrong: } \approx O(\frac{1}{\sqrt{n}})}$$

But actually, this claim is more-or-less (in fact, indeed) correct. But CLT doesn't give it; CLT is only good for "small devs" from mean. Need "Chernoff Bds" for "large" devs

Bounding r.v.'s: more info = better bounds

① Only know mean

② Say r.v. X is ≥ 0 , $X \neq 0$.

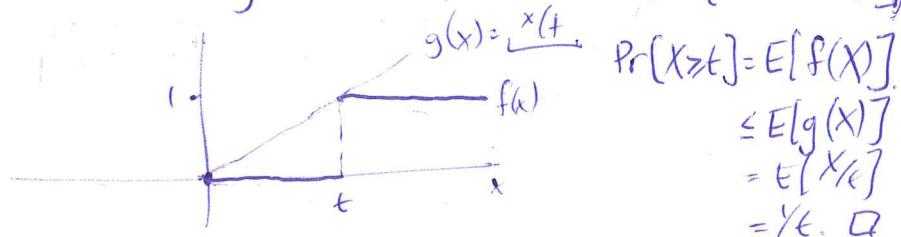
Markov Ineq. (Major workhorse when you know nothing.)

$$\Pr[X \geq t \cdot E[X]] \leq \frac{1}{t} \quad \forall t > 0 \quad [\geq 1]$$

pf 1: "words": WLOG $E[X] = 1$. (Why? Div X by its mean.)
 $\Pr[X \geq t] \leq \frac{1}{t}$?

Well, mean is 1. How could X be $\geq t$ with prob $> \frac{1}{t}$? That would already make X 's mean $> \frac{1}{t} \cdot t = 1$.

pf 2: "pictures"



② Similar fact: Say $0 \leq X \leq 1$ always, $E[X] = \varepsilon$. Then $\Pr[X \geq \varepsilon/2] \geq \varepsilon/2$. (??)

"Markov type"
"averaging drg": Words pf: [Sps $\Pr[X \geq \varepsilon/2] < \varepsilon/2$. Biggest $E(X)$ could be is if whenever it's $\geq \frac{\varepsilon}{2}$ it's 1, and rest of time it's $\frac{\varepsilon}{2}$. But that would still make its expect. $\langle \frac{\varepsilon}{2} \cdot 1 + 1 \cdot \frac{\varepsilon}{2} \rangle < \varepsilon$.]

② Know abt. mean & variance. (\sqrt{Var})

"AKA Second-moment Method"

Chebyshev Ineq: Let $E[X] = \mu$, std dev $[X] = \sigma \neq 0$.

Then $\forall t > 0$, $\Pr[|X - \mu| \geq t \cdot \sigma] \leq \frac{1}{t^2}$.

[words interp.]

pf: WLOG $\mu = 0$ [why?], $\sigma = 1$ [why?]
 $\Leftrightarrow E[X^2] = 1$.

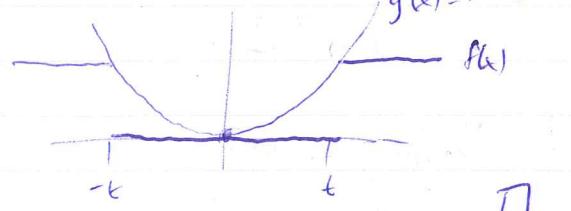
Need $\Pr[|X| \geq t] \leq \frac{1}{t^2}$.

Pf 1: (Markov:) $= \Pr[X^2 \geq t^2] \leq \frac{1}{t^2}$. \square

2. Pics:

$$E[f(X)]$$

$$\geq E[g(X)]$$



Ex cor: Say $X \geq 0$ always. Then $\Pr[X = 0] \leq \frac{\sigma^2}{\mu^2} = \frac{\text{Var}[X]}{E[X]^2} \leq \frac{E[X^2]}{E[X]^2}$

③ Common scenario: $X = X_1 + \dots + X_n$, X_i is indep., or "kinda indep."

E.g. coin flips: $H = X_1 + \dots + X_n$

$$\begin{cases} \text{1 w.p. } 1/2 \\ \text{0 w.p. } 1/2 \end{cases}$$

(this will occupy us for a while) [yuck!]

$$\Pr\left[H \geq \frac{n}{2} + \frac{\sqrt{n}}{2} \cdot 10\sqrt{\ln n}\right] \leq ? \text{ (Markov)} \stackrel{\text{ex}}{\Rightarrow} 1 - \delta\left(\frac{\sqrt{n}}{\sqrt{\ln n}}\right).$$

symm. $\Pr\left[|H - \mu| \geq 10\sqrt{\ln n} \cdot \sigma\right] \leq \frac{1}{100\ln n}$ $\rightarrow 0$ at least.
Still, truth is like $\frac{1}{n^{50}}$.

$$\Pr\left[|H - \mu| \geq 10\sqrt{\ln n} \cdot \sigma\right] \stackrel{\text{Chebyshev?}}{\approx} \frac{1}{100\ln n}$$

④ Chebyshev doesn't need independence, just "pairwise indep." \leftarrow much commoner in practice.

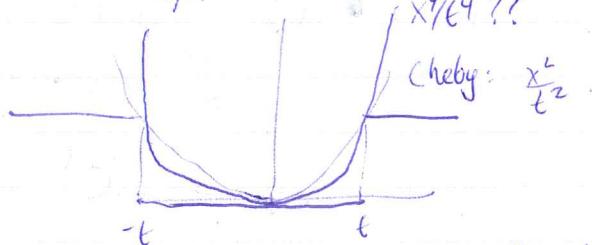
$$\begin{aligned} \text{Var}[X_1 + \dots + X_n] &= E[(X_1 + \dots + X_n)^2] - E[X_1 + \dots + X_n]^2 \\ &= \sum E[X_i^2] + \sum_{i \neq j} E[X_i X_j] - \dots \\ &\quad \sum_{i \neq j} E[X_i] E[X_j] \text{ if pairwise indep.} \\ &= \text{What } \text{Var}(X_1 + \dots + X_n) \text{ would be if } \underline{\text{fully indep.}} \end{aligned}$$

(Back to coins. How can we beat $\frac{1}{2}$ coin?)

Switch to $X = X_1 + \dots + X_n$, X_i indep. $\{-1 \text{ w.p. } \frac{1}{2}, 1 \text{ w.p. } \frac{1}{2}\}$

$$\Pr[|X| \geq 10\sqrt{\ln n} \cdot \sqrt{n}]$$

$\sim 0.5\mu$. $\xrightarrow{\text{standardize?}}$



"4th moment method":

[works w/ 4-wise indep. rvs]

[Markov ver.]

$$\begin{aligned} \Pr[|X| \geq t] &= \Pr[X^4 \geq t^4] \\ &\leq \frac{E[X^4]}{t^4 n^2} \end{aligned}$$

(4)

$$E[X^4]?$$

~~$(\sum X_i)^4 = \sum X_i^4 + \text{X}_i^3 \text{X}_j \text{ terms}$~~

\downarrow \downarrow \downarrow

$+ X_i^2 X_j^2 \text{ terms} + X_i^2 X_j X_k$
 $+ X_i X_j X_k X_l$

In expect: $n \cdot E[X_i^4] = n$.

\downarrow \downarrow \downarrow

$\underbrace{(1) \cdot (4)}_{= 6}$ \downarrow \downarrow

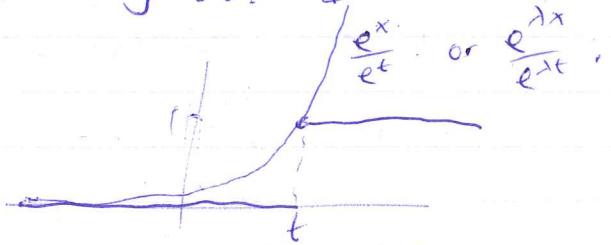
$= 6 \cdot 3n^2 - 3n$
 $= 3n^2 - 2n \leq 3n^2$

$$\therefore \Pr[|X| \geq t\sqrt{n}] \leq \frac{3n^2}{t^4 n^2} = \frac{3}{t^4}, \quad = \frac{3}{100 \ln^2 n} \quad \text{for } t = 10\sqrt{\ln n}.$$

Keep going? $E[X^{2s}] \leq C_s \cdot n^s$
grows with s, kind of pain.

$\therefore \Pr[|X| \geq t\sqrt{n}] \leq \frac{C_s}{t^{2s}}$. Optimize over s...
 [Pain. More flexible way...?]

"Chernoff method":



Markov: $\Pr[X \geq u] = \Pr[e^{\lambda X} \geq e^{\lambda u}]$ for any $\lambda > 0$
 $\leq \frac{E[e^{\lambda X}]}{e^{\lambda u}}$ [a func of λ , called M.G.F.]

[a number] $\rightarrow \frac{1}{e^{\lambda u}}$

easy to compute if $X = X_1 + \dots + X_n$ indep.

Plan: Compute/bound it, then optimize λ .

$$\begin{aligned}
 E[e^{\lambda X}] &= E[\exp(\lambda X_1 + \dots + \lambda X_n)] \\
 &= E[\exp(\lambda X_1) \cdot \dots \cdot \exp(\lambda X_n)] \\
 &= E[\exp(\lambda X_1)] \cdots E[\exp(\lambda X_n)] \quad (\text{by indep.}) \\
 &= E[\exp(\lambda X_1)]^n \\
 &= \left[\frac{1}{2} e^\lambda + \frac{1}{2} e^{-\lambda} \right]^n \\
 &\stackrel{x}{=} \left[1 + \frac{x^2}{2!} + \frac{x^4}{4!} + \dots \right]^n \\
 &= \left[1 + \frac{\lambda^2}{2!} + \frac{\lambda^4}{4!} + \dots \right]^n \\
 &\leq \left(e^{\lambda^2/2} \right)^n = 1 + \frac{\lambda^2}{2} + \frac{\lambda^4}{2^2 \cdot 2!} + \frac{\lambda^6}{2^3 \cdot 3!} + \dots
 \end{aligned}$$

$$\text{So } \Pr[X \geq u] \leq \frac{e^{\lambda^2/2}}{e^{u\lambda}} = \exp\left(\frac{1}{2}\lambda^2 - u\lambda\right) \quad \text{minimize over } \lambda$$

$\lambda := \frac{u}{\lambda}$

$$= e^{-\frac{u^2}{2\lambda}}$$

$$\text{So for } u = 10\sqrt{n} \cdot \sqrt{\lambda}, \quad \leq \frac{1}{n^{50}} \quad \circlearrowright$$

[Specific to sum of iid 50/50's]

Next: treat $X_i = \begin{cases} 1 & \text{w.p. } p_i \\ 0 & \text{w.p. } 1-p_i \end{cases}$. Must bound $E[\exp(\lambda X_i)]$, choose λ well.

Moderately annoying calcs.

\Rightarrow

(6)

Let $X = X_1 + \dots + X_n$, X_i 's indep.
 Let $\mu = E[X] = \sum E[X_i]$

"Hoeffding": Say $a_i \leq X_i \leq b_i$ always.

$$\Pr[X \geq \mu + t], \Pr[X \leq \mu - t] \leq \exp\left(-\frac{2t^2}{\sum (b_i - a_i)^2}\right) \quad t > 0$$

[Jibes w/ our result for $X = \pm 1$]

[better if μ is surprisingly small]

"Chernoff": Say $0 \leq X_i \leq 1$

$$t > 0, \Pr[X \leq (1-\varepsilon)\mu] \leq \exp\left(-\frac{\varepsilon^2}{2}\mu\right)$$

$$\Pr[X \geq (1+\varepsilon)\mu] \leq \exp\left(-\frac{\varepsilon^2}{2+\varepsilon}\mu\right) \quad \text{[can't delete]} \\ \leq \exp\left(-\frac{\varepsilon^2}{3}\mu\right) \quad \text{if } \varepsilon \leq 1$$

Fact: If $\mu_L \leq \mu \leq \mu_H \dots$

(Memorize this one like a poem)
 Look up the rest
 [Give "sampling hints" there]

Bonus results | Neg. dep. r.v.s : do first

Def: X_1, \dots, X_n are negatively associated (NA) if:

$$E[f(X_i; i \in A)g(X_i; i \in B)] \leq E[f(X_i; i \in A)]E[g(X_i; i \in B)]$$

If nondecr f,g, $A, B \subseteq [n]$.

$$\text{ex: } \Rightarrow E[e^{\lambda X_1 + \dots + \lambda X_n}] \leq E[e^{\lambda X_1}] \cdots E[e^{\lambda X_n}]$$

\Rightarrow Hoeffding/Chernoff holds

$$\Rightarrow \text{Corr}(X_i, X_j) \leq 0$$



Facts: • indep \Rightarrow N.A.

- ~~• closed under unions, intersections, subsets~~
- closed under subsets, ^{indep} unions
- " " applying nondecr fns to disjoint subsets; e.g.: $X_1 + X_2 + X_3, X_4 + \dots + X_{10}, X_{11} + X_{12}, \dots$

eg: [see Jong-Dev - Proschan 83]

- ① Put n pts randomly on unit circle, X_1, \dots, X_n = arc lengths
- ② Throw N balls into n bins [[diff probs for each bin]]
 X_1, \dots, X_n occupancies " " $"$ ball ok]
- ③ Let x_1, \dots, x_n be fixed #s, X_1, \dots, X_n a rand permutation. \Rightarrow sampling from finite population WITHOUT replacement

All N.A. \rightarrow Intuitive meaning: knowing that some values are "large" makes it more likely other values are "small".

Martingales: Let X_1, \dots, X_n be ~~any~~ any ^{discrete} r.v.'s. (not nec indep)

let $f: \mathbb{R}^n \rightarrow \mathbb{R}$

Let $Y_i = E[f(X_1, \dots, X_n) | X_1, \dots, X_{i-1}]$, so

$Y_0 = E[f(X)]$, (constant nu.) $Y_n = f(X_1, \dots, X_n)$

Then $\mu := Y_0, Y_1, \dots, Y_n$ is a "martingale" (w.r.t. X_1, \dots, X_n), meaning $E[Y_i | X_1, \dots, X_{i-1}] = Y_{i-1} \quad \forall i \in \{1, \dots, n\}$.



"Azuma" / "Method of Bdd Diffs"

thm: Suppose X_1, \dots, X_n indep, f sats

$$|f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n) - f(x_1, \dots, x_{i-1}, x'_i, x_{i+1}, \dots, x_n)| \leq c_i$$

Then $\Pr[f(X) \geq \mu + t] \leq \exp\left(-\frac{2t^2}{\sum c_i^2}\right)$. always