# <span id="page-0-0"></span>Dependence between Path Lengths and Size in Random Trees (joint with H.-H. Chern, H.-K. Hwang and R. Neininger)

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Hsinchu, Taiwan

Kraków, July 4, 2016

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 $\rightarrow$   $\rightarrow$   $\rightarrow$ Michael Fuchs (NCTU) [Dependence in Random Trees](#page-80-0) Kraków, July 4, 2016 1/30

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Proposed by Muntz and Uzgalis in 1971.

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Input: 6, 2, 4, 8, 7, 1, 5, 3, 10, 9
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If permutations are equally likely  $\longrightarrow$  random  $m$ -ary search trees

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# Size, KPL, and NPL

#### • Size (or Storage Requirement)

Number of nodes holding keys. Only random if  $m \geq 3$ .

 $S_n$  = size of a random m-ary search tree built from n keys.

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 $\langle \vert \bar{m} \vert \rangle$  ,  $\langle \vert \bar{m} \vert \rangle$  ,  $\langle \vert \bar{m} \rangle$  ,  $\langle \vert \bar{m} \rangle$ 

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Sum of all key-distances to the root.

 $K_n =$  KPL of a random  $m$ -ary search tree built from  $n$  keys.

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#### • Node Path Length (NPL)

Sum of all node-distances to the root.

 $N_n = \text{NPL}$  of a random m-ary search tree built from n keys.

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# Size: Mean

Knuth (1973):

$$
\mathbb{E}(S_n) \sim \phi n,
$$

where

$$
\phi:=\frac{1}{2(H_m-1)}
$$

and  $H_m$  are the Harmonic numbers.

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and  $H_m$  are the Harmonic numbers.

Mahmoud and Pittel (1989):

$$
\mathbb{E}(S_n) = \phi(n+1) - \frac{1}{m-1} + \mathcal{O}(n^{\alpha-1}),
$$

where  $\alpha$  is the real part of the second largest zero of

$$
\Lambda(z) = z(z+1)\cdots(z+m-2) - m!.
$$

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## Size: Phase Change for Variance

#### Mahmoud and Pittel (1989):

$$
\operatorname{Var}(S_n) \sim \begin{cases} C_S n, & \text{if } m \le 26; \\ F_1(\beta \log n) n^{2\alpha - 2}, & \text{if } m \ge 27, \end{cases}
$$

where  $\lambda = \alpha + i\beta$  is the second largest zero of  $\Lambda(z)$ .

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$$

where  $\lambda = \alpha + i\beta$  is the second largest zero of  $\Lambda(z)$ .

Here,  $F_1(z)$  is the periodic function

$$
F_1(z) = 2 \frac{|A|^2}{|\Gamma(\lambda)|^2} \left( -1 + \frac{m!(m-1)|\Gamma(\lambda)|^2}{\Gamma(2\alpha + m - 2) - m!\Gamma(2\alpha - 1)} \right) + 2\Re \left( \frac{A^2 e^{2iz}}{\Gamma(\lambda)^2} \left( -1 + \frac{m!(m-1)\Gamma(\lambda)^2}{\Gamma(2\lambda + m - 2) - m!\Gamma(2\lambda - 1)} \right) \right)
$$

with  $A = 1/(\lambda(\lambda-1)\sum_{0 \leq j \leq m-2} \frac{1}{j+1}$  $\frac{1}{j+\lambda}$ ).

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### Size: Phase Change for Limit Law

Theorem (Mahmoud & Pittel (1989); Lew & Mahmoud (1994)) For  $3 \le m \le 26$ ,  $S_n - \mathbb{E}(S_n)$  $\sqrt{\text{Var}(S_n)}$  $\stackrel{d}{\longrightarrow} N(0,1),$ 

where  $N(0, 1)$  is the standard normal distribution.

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where  $N(0, 1)$  is the standard normal distribution.

Theorem (Chern & Hwang (2001)) For  $m \geq 27$ ,  $S_n - \mathbb{E}(S_n)$  $\sqrt{\text{Var}(S_n)}$ does not converge to a fixed limit law.

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## KPL: Moments

### Mahmoud (1986):

$$
\mathbb{E}(K_n) = 2\phi n \log n + c_1 n + o(n),
$$

where  $c_1$  is an explicitly computable constant.

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### Mahmoud (1992):

$$
\text{Var}(K_n) \sim C_K n^2,
$$

where

$$
C_K = 4\phi^2 \left( \frac{(m+1)H_m^{(2)} - 2}{m-1} - \frac{\pi^2}{6} \right)
$$

with  $H_m^{(2)}=\sum_{1\leq j\leq m}1/j^2.$ 

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with  $H_m^{(2)}=\sum_{1\leq j\leq m}1/j^2.$ 

So, no phase change here for the variance!

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# KPL: Limit Law

Theorem (Neininger & Rüschendorf (1999)) We have.

$$
\frac{K_n - \mathbb{E}(K_n)}{n} \stackrel{d}{\longrightarrow} K,
$$

where  $K$  is the unique solution of

$$
X \stackrel{d}{=} \sum_{1 \le r \le m} V_r X^{(r)} + 2\phi \sum_{1 \le r \le m} V_r \log V_r
$$

with  $X^{(r)}$  an independent copy of  $X$  and

$$
V_r = U_{(r)} - U_{(r-1)},
$$

where  $U_{(r)}$  is the  $r\text{-}th$  order statistic of  $m$  i.i.d. uniform RVs.

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where  $c_2$  is an explicitly computable constant.

We have,

$$
\begin{cases}\nS_n \stackrel{d}{=} S_{I_1}^{(1)} + \dots + S_{I_m}^{(m)} + 1, \\
N_n \stackrel{d}{=} N_{I_1}^{(1)} + \dots + N_{I_m}^{(m)} + S_{I_1}^{(1)} + \dots + S_{I_m}^{(m)}.\n\end{cases}
$$

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$$

So, one expects a strong positive dependence between  $S_n$  and  $N_n!$ 

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# Size and NPL: Correlation (i)

## Theorem (Chern, F., Hwang, Neininger) We have,

$$
Cov(S_n, N_n) \sim \begin{cases} C_R n \log n, & \text{if } 3 \le m \le 13; \\ \phi F_2(\beta \log n) n^{\alpha}, & \text{if } m \ge 14, \end{cases}
$$

where  $C_R$  is a constant and  $F_2(z)$  is a periodic function. Moreover,

 $\text{Var}(N_n) \sim \phi^2 C_K n^2$ .

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# Size and NPL: Correlation (i)

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where  $C_R$  is a constant and  $F_2(z)$  is a periodic function. Moreover,

$$
Var(N_n) \sim \phi^2 C_K n^2.
$$

Thus (!),

$$
\rho(S_n, N_n) \begin{cases} \longrightarrow 0, & \text{if } 3 \le m \le 26; \\ \sim \frac{F_2(\beta \log n)}{\sqrt{C_K F_1(\beta \log n)}}, & \text{if } m \ge 27. \end{cases}
$$

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# Size and NPL: Correlation (ii)



Periodic function of  $\rho(S_n, N_n)$  for  $m = 27, 54, \ldots, 270$ .

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**Pearson:** for RVs  $X$  and  $Y$ 

$$
\rho(X, Y) = \frac{\text{Cov}(X, Y)}{\sqrt{\text{Var}(X)\text{Var}(Y)}}.
$$

Measures linear dependence between  $X$  and  $Y!$ 

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#### Refined correlation measures:

Distance correlation, Brownian covariance, mutual information, total correlation, dual total correlation, etc.

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**Question**: Can our counterintuitive result for  $\rho(S_n, N_n)$  be ascribed to the weakness of Pearson's correlation coefficient?

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Distance correlation, Brownian covariance, mutual information, total correlation, dual total correlation, etc.

**Question**: Can our counterintuitive result for  $\rho(S_n, N_n)$  be ascribed to the weakness of Pearson's correlation coefficient? NO!

# Size and NPL: Limit Law for  $3 \le m \le 26$

Theorem (Chern, F., Hwang, Neininger)

Consider

$$
Q_n = (S_n, N_n).
$$

Then,

$$
Cov(Q_n)^{-1/2}\Big(Q_n - \mathbb{E}(Q_n)\Big) \stackrel{d}{\longrightarrow} \Big(N, C_K^{-1/2}K\Big),\,
$$

where N has a standard normal distribution.

Moreover,  $N$  and  $K$  are independent!

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$$

where N has a standard normal distribution.

Moreover,  $N$  and  $K$  are independent!

Thus, asymptotic independence for  $3 \le m \le 26$  is also observed in the bivariate limit law!

# Size and NPL: Limit Law for  $m \geq 27$

Theorem (Chern, F., Hwang, Neininger)

Consider

$$
Y_n = \left(\frac{S_n - \phi n}{n^{\alpha - 1}}, \frac{N_n - \mathbb{E}(N_n)}{n}\right).
$$

Then,

$$
\ell_2(Y_n, (\Re(n^{i\beta}\Lambda), \phi K)) \longrightarrow 0,
$$

where  $\ell_2$  is the minimal  $L_2$ -metric and  $\Lambda$  is the unique solution of

$$
W \stackrel{d}{=} \sum_{1 \le r \le m} V_r^{\lambda - 1} W^{(r)}
$$

with  $W^{(r)}$  independent copies of  $W.$ 

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# Size and KPL

Same results hold for size and KPL, e.g.,

$$
\rho(S_n, K_n) \begin{cases} \longrightarrow 0, & \text{if } 3 \le m \le 26; \\ \sim \rho(S_n, N_n), & \text{if } m \ge 27. \end{cases}
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$$

Theorem (Chern, F., Hwang, Neininger) We have  $\rho(K_n, N_n) \sim 1$  and

$$
||N_n - \phi K_n - (\mathbb{E}(N_n - \phi K_n))||_2 = o(n).
$$

In particular,

$$
\left(\frac{K_n-\mathbb{E}(K_n)}{n},\frac{N_n-\mathbb{E}(N_n)}{n}\right)\stackrel{d}{\longrightarrow}(K,\phi K).
$$

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Proposed by René de la Briandais in 1959.

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Name from the word data retrieval (suggested by Fredkin).

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Example:

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Proposed by René de la Briandais in 1959.

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Example:



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#### Two types of nodes:

- Internal nodes: only used for branching;
- External nodes: nodes which hold data.

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- Internal nodes: only used for branching;
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#### Random Model:

Bits are independent Bernoulli random variables with mean  $p$ .

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- $p = 1/2$ : symmetric trie;
- $\bullet$   $p \neq 1/2$ : asymmetric trie.

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Question: correlation between size and path-length in random tries?

#### Size

Number of internal nodes.

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#### Size

Number of internal nodes.

### External Path Length (EPL)

Sum of all distances between external nodes and root.

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#### Size

Number of internal nodes.

External Path Length (EPL)

Sum of all distances between external nodes and root.

### • Internal Path Length (IPL)

Sum of all distances between internal nodes and root.

We again use  $S_n, K_n, N_n$  and



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We use the following notation:

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We use the following notation:

 $\bullet$  Entropy:  $h = -p \log p - q \log q$ ;

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We use the following notation:

- $\bullet$  Entropy:  $h = -p \log p q \log q$ ;
- If  $\log p / \log q \in \mathbb{Q}$ , then

$$
\frac{\log p}{\log q} = \frac{r}{\ell}, \quad \gcd(r,\ell) = 1.
$$

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and

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$$

$$
\chi_k = \frac{2rk\pi i}{\log p}, \qquad (k \in \mathbb{Z}).
$$

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We use the following notation:

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$$
  
and  

$$
\chi_k = \frac{2rk\pi i}{\log p}, \qquad (k \in \mathbb{Z}).
$$

• For a sequence  $g_k$ :

$$
\mathscr{F}[g](z) = \begin{cases} \sum_{k \in \mathbb{Z}} g_k z^{-\chi_k}, & \text{if } \log p / \log q \in \mathbb{Q}; \\ g_0, & \text{if } \log p / \log q \notin \mathbb{Q}. \end{cases}
$$

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## Means and Variances



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## Means and Variances



Note that

$$
N_n \approx S_n \frac{\log n}{h}
$$

and for  $p \neq q$ 

$$
Var(D_n) \sim \frac{Var(K_n)}{n}.
$$

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#### **• Rice Method**

Exercise 54 in Section 5.2.2 of Knuth's book. Developed into a systematic tool by Flajolet and Sedgewick.

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#### **• Two-stage Approach**

Introduced by Jacquet and Régnier. Further developed by Jacquet and Szpankowski.

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#### **• Poisson Variance and Covariance**

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#### Other Approaches

Elementary (Schachinger), probabilistic (Devroye, Janson, etc.)

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## Size: Variance

Theorem (Régnier & Jacquet (1989); Kirschenhofer & Prodinger (1991); F., Hwang, Zacharovas (2014))

We have.

$$
\text{Var}(S_n) \sim \mathscr{F}[g^{(1)}](n)n,
$$

where

$$
g_k^{(1)} = \frac{\chi_k \Gamma(-1 + \chi_k)}{h} \left(1 - \frac{\chi_k + 3}{2^{1 + \chi_k}}\right)
$$
  

$$
- \frac{1}{h^2} \sum_{j \in \mathbb{Z}} \Gamma(\chi_j + 1) \Gamma(\chi_{k-j} + 1)
$$
  

$$
- \frac{2}{h} \sum_{j \ge 1} \frac{(-1)^j (j + 1 + \chi_k) \Gamma(j + \chi_k) (p^{j+1} + q^{j+1})}{(j-1)!(j+1)(1 - p^{j+1} - q^{j+1})}.
$$

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## Size and NPL: Variances and Covariance

Theorem (F., Hwang, Zacharovas (2014))  
\nWe have,  
\n
$$
Var(S_n) \sim \mathcal{F}[g^{(1)}](n)n,
$$
\nand  
\n
$$
Cov(S_n, N_n) \sim \mathcal{F}[g^{(1)}](n) \frac{n \log n}{h}
$$
\nand  
\n
$$
Var(N_n) \sim \mathcal{F}[g^{(1)}](n) \frac{n \log^2 n}{h^2}.
$$

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 $\left\{ \begin{array}{ccc} 1 & 0 & 0 \\ 0 & 1 & 0 \end{array} \right.$ 

# Size and NPL: Variances and Covariance

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\n
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Var(N_n) \sim \mathscr{F}[g^{(1)}](n) \frac{n \log^2 n}{h^2}.
$$
\nCorollary  
\nWe have,  
\n
$$
\rho(S_n, N_n) \longrightarrow 1.
$$

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# Size and NPL: Limit Law

#### Theorem (F. & Lee (2015))

We have,

$$
\left(\frac{S_n - \mathbb{E}(S_n)}{\sqrt{\text{Var}(S_n)}}, \frac{N_n - \mathbb{E}(N_n)}{\sqrt{\text{Var}(N_n)}}\right) \xrightarrow{d} \mathcal{N}(0, E_2),
$$

where  $E_2$  is the  $2 \times 2$  unit matrix

$$
E_2 = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}.
$$

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# Size and NPL: Limit Law

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NPL was also investigated by Nguyen-The (2004) in his PhD-thesis, but his result is incorrect.

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# Size and NPL: Limit Law

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NPL was also investigated by Nguyen-The (2004) in his PhD-thesis, but his result is incorrect.

Question: same result for size and KPL?

# Size and KPL: Covariance

#### Theorem (F. & Hwang)

We have,

$$
Cov(S_n, K_n) \sim \mathscr{F}[g^{(2)}](n)n,
$$

where

$$
g_k^{(2)} = \frac{\Gamma(\chi_k)}{h} \left( 1 - \frac{\chi_k + 2}{2\chi_{k+1}} \right)
$$
  
- 
$$
\frac{1}{h^2} \sum_{j \in \mathbb{Z} \setminus \{0\}} \Gamma(\chi_{k-j} + 1)(\chi_j - 1) \Gamma(\chi_j)
$$
  
- 
$$
\frac{\Gamma(\chi_k + 1)}{h^2} \left( \gamma + 1 + \psi(\chi_k + 1) - \frac{p \log^2 p + q \log^2 q}{2h} \right)
$$
  
+ 
$$
\frac{1}{h} \sum_{j \ge 2} \frac{(-1)^j (2j^2 - 2j + 1 + (2j - 1)\chi_k) \Gamma(j - 1\chi_k) (p^j + q^j)}{j!(1 - p^j - q^j)}.
$$

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### Correlation Coefficient

 $p = q = 1/2$ :



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# Size and KPL: Correlation Coefficient

#### Theorem (F. & Hwang)

We have.

$$
\rho(S_n, K_n) \sim \begin{cases} 0, & \text{if } p \neq q; \\ F(n), & \text{if } p = q, \end{cases}
$$

where

$$
F(n) = \frac{\mathscr{F}[g^{(2)}](n)}{\sqrt{\mathscr{F}[g^{(1)}](n)\mathscr{F}[g^{(3)}](n)}}
$$

is a periodic function with

average value =  $0.927\cdots$  and amplitude  $\leq 1.5\times 10^{-5}$ .
# Size and KPL: Correlation Coefficient

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Question: is now this behavior due to the weakness of Pearson's correlation coefficient?

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# Size and KPL: Correlation Coefficient

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is a periodic function with

average value =  $0.927\cdots$  and amplitude  $\leq 1.5\times 10^{-5}$ .

Question: is now this behavior due to the weakness of Pearson's correlation coefficient? Again NO!

## Size and KPL: Limit Laws

Theorem (F. & Hwang)\n
$$
\bullet \ p \neq q:
$$
\n
$$
\left(\frac{S_n - \mathbb{E}(S_n)}{\sqrt{\text{Var}(S_n)}}, \frac{K_n - \mathbb{E}(K_n)}{\sqrt{\text{Var}(K_n)}}\right) \xrightarrow{d} \mathcal{N}(0, I_2),
$$

where  $I_2$  is the  $2 \times 2$  identity matrix.

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### Size and KPL: Limit Laws

### Theorem (F. & Hwang)

$$
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\n
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\left(\frac{S_n - \mathbb{E}(S_n)}{\sqrt{\text{Var}(S_n)}}, \frac{K_n - \mathbb{E}(K_n)}{\sqrt{\text{Var}(K_n)}}\right) \xrightarrow{d} \mathcal{N}(0, I_2),
$$

where  $I_2$  is the  $2 \times 2$  identity matrix.

$$
\bullet \ \ p = q.
$$

$$
\Sigma_n^{-1/2}\Big(S_n - \mathbb{E}(S_n), K_n - \mathbb{E}(K_n)\Big) \stackrel{d}{\longrightarrow} \mathcal{N}_2(0, I_2),
$$

where  $\Sigma_n$  is the (asymptotic) covariance matrix:

$$
\Sigma_n := n \begin{pmatrix} \mathscr{F}[g^{(1)}](n) & \mathscr{F}[g^{(2)}](n) \\ \mathscr{F}[g^{(2)}](n) & \mathscr{F}[g^{(3)}](n) \end{pmatrix}.
$$

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### Joint Distribution of  $S_n$  and  $K_n$





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Similar results for fringe-balanced binary search trees and quadtrees: H.-H. Chern, M. Fuchs, H.-K. Hwang, R. Neininger. Dependence and phase changes in random  $m$ -ary search trees, arxiv:1501.05135.

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- Better explanation of our results?