

How do you describe concurrent code?  
`await, async, asyncio.run(), asyncio.to_thread()`

To run things concurrently: `asyncio.gather()`

How does concurrent code share state?

Until you call `await`, nothing will change

Mutable objects may be different after `await`

```
async def apollo():
    await asyncio.sleep(600)  A coroutine function
    # (may await!)
async def kay():
    await doordash_boba()  Give up control until
    # asyncio.sleep(600)    finishes
async def harry():
    await cupcake()
async def start():
    await asyncio.gather(  Run all of the
        apollo(),          awaitables (e.g.,
        kay(),             coroutines)
        harry(),           concurrently
    )
    # Start an event loop
asyncio.run(start())  (environment that knows how to
                      handle concurrent work)
```

Exceptions are raised with a `raise` statement.

```
raise <expr>
```

`<expr>` must evaluate to a subclass of `BaseException` or an instance of one.

```
try:
    <try suite>
except <exception class> as <name>:
    <except suite>
The <try suite> is executed first.
If, during the course of executing the
<try suite>, an exception is raised
that is not handled otherwise, and
If the class of the exception inherits from <exception class>, then
The <except suite> is executed, with <name> bound to the exception.
```

```
(append s t): list the elements of s and t; append can be called
on more than 2 lists
(map f s): call a procedure f on each element of a list s and list
the results
(filter f s): call a procedure f on each element of a list s and
list the elements for which a true value is the result
(apply f s): call a procedure f with the elements of a list as its
arguments
```

```
(define size 5) ; => size
(* 2 size) ; => 10
(if (> size 0) size (- size)) ; => 5
(cond ((> size 0) size) ((= size 0) 0) (else (- size))) ; => 5
((lambda (x y) (+ x y size)) size (+ 1 2)) ; => 13
(map (lambda (x) (+ x size)) (quote (2 3 4))) ; => (7 8 9)
(filter odd? (quote (2 3 4))) ; => (3)
(list (cons 1 nil) size 'size) ; => ((1) 5 size)
(list (equal? 1 2) (null? nil) (= 3 4) (eq? 5 5)) ; => (#f #t #f #t)
(list (or #f #t) (or) (or 1 2)) ; => (#t #f 1)
(list (and #f #t) (and) (and 1 2)) ; => (#f #t 2)
(list 'a 2) ; => (a 2)
(append '(1 2) '(3 4)) ; => (1 2 3 4)
(not (> 1 2)) ; => #t
(begin (define x (+ size 1)) (* x 2)) ; => 12
```

```
(define (factorial n)
  (if (= n 0) 1
      (* n (factorial (- n 1)))))
(define (fib n)
  (cond
    ((= n 0) 0)
    ((= n 1) 1)
    (else (+ (fib (- n 2)) (fib (- n 1))))))
(define (nines num)
  (if (= num 0)
      0
      (if (= (= (modulo num 10) 9)
            (+ 1 (nines (floor (/ num 10)))))
```

The way in which names are looked up in Scheme and Python is called **lexical scope** (or **static scope**).

**Lexical scope:** The parent of a frame is the environment in which a procedure was **defined**. `(lambda ...)`

**Dynamic scope:** The parent of a frame is the environment in which a procedure was **called**. `(mu ...)`

```
> (define f (mu (x) (+ x y)))
> (define g (lambda (x y) (f (+ x x))))
> (g 3 7)
```

A table has columns and rows

Latitude	Longitude	Name
38	122	Berkeley
42	71	Cambridge
45	93	Minneapolis

A column has a name and a type

A row has a value for each column

```
SELECT [expression] AS [name], [expression] AS [name], ...;
```

```
SELECT [columns] FROM [table] WHERE [condition] ORDER BY [order];
```

Tables **A** & **B** are joined by **JOIN** (or a comma) to form combos of an **A row** & a **B row**

A join often has some conditions for matching up the rows of two (or more) tables

**titles**

tconst	title	year
tt8267604	Capernaum	2018
tt8367814	The Gentlemen	2019
tt8404614	The Two Popes	2019
tt8503618	Hamilton	2020

**ratings**

tconst	avgRating	numVotes
tt8503618	8.3	134421
tt8579674	8.2	749471
tt8613070	8.0	123438
tt8404614	7.6	144516

```
SELECT * FROM titles JOIN ratings ON titles.tconst=ratings.tconst;
```

tconst	title	year	tconst	averageRating	numVotes
tt8404614	The Two Popes	2019	tt8404614	7.6	144516
tt8503618	Hamilton	2020	tt8503618	8.3	134421

**Explicit join syntax:** Use `FROM [table] JOIN [table] ON [condition]`

```
SELECT title, averageRating FROM titles JOIN ratings
  ON titles.tconst=ratings.tconst LIMIT 3;
```

**Implicit join syntax:** Use a comma (or just **JOIN**) and put all conditions in the **WHERE** clause

```
SELECT title, averageRating FROM titles, ratings
  WHERE titles.tconst=ratings.tconst LIMIT 3;
```

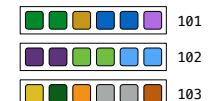
Leaving out the **ON** or **WHERE** clause is allowed and creates all pairs of rows  
A table can be joined by itself to compare one row to another row  
Aliases are required

E.g., create a table of remakes: two movies that have the same title

```
SELECT old.title, old.year AS first, new.year AS second
  FROM titles AS old JOIN titles AS new
    ON old.title=new.title AND old.year < new.year;
```

**CREATE TABLE** lift AS

```
SELECT 101 AS chair, 2 AS single, 2 AS pair UNION
SELECT 102      , 0      , 3      UNION
SELECT 103      , 4      , 1      ;
```



String values can be combined to form longer strings

```
sqlite> SELECT "hello," || " world";
hello, world
```

Basic string manipulation is built into SQL, but differs from Python

```
sqlite> CREATE TABLE phrase AS SELECT "hello, world" AS s;
sqlite> SELECT substr(s, 4, 2) || substr(s, instr(s, " ")+1, 1)
  FROM phrase;
low
```

The number of groups is the number of unique values of an expression  
A **having** clause filters the set of groups that are aggregated

```
SELECT weight/legs, count(*)
  FROM animals
 GROUP BY weight/legs
 HAVING COUNT(*)>1;
```

kind	legs	weight
dog	4	20
cat	4	10
ferret	4	10
parrot	2	6
penguin	2	10
t-rex	2	12000

An aggregate function in the **[columns]** clause computes a value from a group of rows:

- `MAX([expression])` evaluates to the largest value of `[expression]` for any row in a group
- `COUNT(*)` evaluates to the number of rows in a group
- `MIN, SUM, & AVG` are also aggregate functions similar to `MAX`

With no **GROUP BY** clause, aggregation is performed over all rows:

```
select max(legs) from animals;
```

max(legs)
4

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient, ...
- Combinations: (quotient 10 2), (not true), ...

Numbers are self-evaluating; *symbols* are bound to values. Call expressions have an operator and 0 or more operands.

A combination that is not a call expression is a *special form*:

- If expression: (if <predicate> <consequent> <alternative>)
- Binding names: (define <name> <expression>)
- New procedures: (define (<name> <formal parameters>) <body>)

```

> (define pi 3.14)           > (define (abs x)
> (* pi 2)                  (if (< x 0)
6.28                         (- x)
                           x)
> (abs -3)                  3

```

Lambda expressions evaluate to anonymous procedures.

```
(lambda (<formal-parameters>) <body>)
```



Two equivalent expressions:

```
(define (plus4 x) (+ x 4))
(define plus4 (lambda (x) (+ x 4)))
```

An operator can be a combination too:

```
((lambda (x y z) (+ x y (square z))) 1 2 3)
```

In the late 1950s, computer scientists used confusing names.

- **cons**: Two-argument procedure that **creates a Link**
- **car**: Procedure that returns the **first element** of a Link
- **cdr**: Procedure that returns the **second element** of a Link
- **nil**: The empty list

They also used a non-obvious notation for linked lists.

- A (linked) Scheme list has a first element and the rest, which may either be a Scheme list or nil (the empty Scheme list)
- Scheme lists are written as space-separated combinations.

```

> (define x (cons 1 nil))
> x
(1)
> (car x)
1
> (cdr x)
()
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))
(1 2 3 4)

```

Symbols normally refer to values; how do we refer to symbols?

```

> (define a 1)
> (define b 2)
> (list a b)
(1 2)

```

No sign of "a" and "b" in the resulting value

Quotation is used to refer to symbols directly in Lisp.

```

> (list 'a 'b)
(a b)
> (list 'a b)
(a 2)

```

Symbols are now values

Quotation can also be applied to combinations to form lists.

```

> (car '(a b c))
a
> (cdr '(a b c))
(b c)

```

```

(car (cons 1 nil)) -> 1
(cdr (cons 1 nil)) -> ()
(cdr (cons 1 (cons 2 nil))) -> (2)

```

The built-in Scheme list data structure (which is a linked list) can represent a Scheme expression

```

scm> (list 'quotient 10 2)      scm> (eval (list 'quotient 10 2))
(quotient 10 2)                  5
scm> (define (fact n)
      (if (= n 1) 1 (* n (fact (- n 1)))))
fact
scm> (fact 5)
120
scm> (define (fact-expr n)
      (if (= n 1) 1 (list '* n (fact-expr (- n 1)))))
fact-expr
scm> (fact-expr 5)
(* 5 (* 4 (* 3 (* 2 1))))
scm> (define (fact-expr-quasiquoted n)
      (if (= n 1) 1 `(* ,n ,(fact-expr-quasiquoted (- n 1)))))
fact-expr-quasiquoted
scm> (fact-expr-quasiquoted 5)
(* 5 (* 4 (* 3 (* 2 1))))

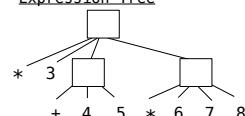
```

The Calculator language has primitive expressions and call expressions

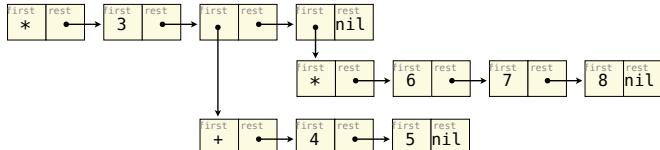
Calculator Expression

Expression Tree

```
(* 3
  (+ 4 5)
  (* 6 7 8))
```



Representation as Link objects



A Scheme list is written as elements in parentheses:

```
(<element> <element> ... <element>))
```

A Scheme list

Each <element> can be a combination or atom (primitive).

```
(+ (* 3 (+ (* 2 4) (+ 3 5))) (+ (- 10 7) 6))
```

def reduce(f, s, initial):

"""Combine elements of s pairwise using f, starting with initial.  
 E.g., reduce(mul, [2, 4, 8], 1) is equivalent to  
 mul(mul(mul(1, 2), 4), 8).

```
>>> reduce(mul, [2, 4, 8], 1)
```

64

"""

```
    for x in s:
        initial = f(initial, x)
    return initial
```

f is ...

a two-argument function that returns a first argument

s is ...

a sequence of values that can be the second argument

initial is ...

a value that can be the first argument

The Scheme version of reduce doesn't have an initial argument:

```
scm> (reduce * (list 2 4 8))
```

64

Base cases:

- Primitive values (numbers)
- Look up values bound to symbols

**Eval**

Recursive calls:

- Eval(operator, operands) of call expressions
- Apply(procedure, arguments)
- Eval(sub-expressions) of special forms

The structure of the Scheme interpreter

Creates a new environment each time a user-defined procedure is applied

Requires an environment for name lookup

Base cases:

- Built-in primitive procedures

Recursive calls:

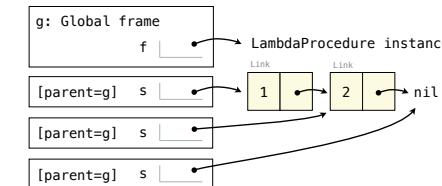
- Eval(body) of user-defined procedures

**Apply**

To apply a user-defined procedure, create a new frame in which formal parameters are bound to argument values, whose parent is the **env** of the procedure, then evaluate the body of the procedure in the environment that starts with this new frame.

```
(define (f s) (if (null? s) '() (cons (car s) (f (cdr s)))))

(f (list 1 2))
```



There are two ways to quote an expression

Quote: '(a b) => (a b)

Quasiquote: `(a b) => (a b)

Parts of a quasiquoted expression can be unquoted with ,

(define b 4)

Quote: '(a ,(+ b 1)) => (a (unquote (+ b 1)))

Quasiquote: `(a ,(+ b 1)) => (a 5)

Quasiquotation is convenient for generating Scheme expressions:

(define (make-add-lambda n) `(lambda (d) (+ d ,n)))

(make-add-lambda 2) => (lambda (d) (+ d 2))

A macro is an operation performed on the source code of a program before evaluation. Macros exist in many languages, but are easiest to define correctly in a language like Lisp. Scheme has a **define-macro** special form that defines a source code transformation

Constructs & evaluates the expression:  
(begin (print 2) (print 2))

```
(define-macro (twice expr)      > (twice (print 2))
  (list 'begin expr expr))      2

```

Evaluation procedure of a macro call expression:

- Evaluate the operator sub-expression (evaluates to a macro)

- Call the macro procedure on the operand expressions *without evaluating them first*

- Evaluate the expression returned from the macro procedure