

## Usage of GPU for online data processing: The experience of ALICE and LHCb

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- **Overview of upgrades of LHC and the experiments:** 
  - What are the upcoming challenges for ALICE and LHCb?
  - What do the online processing approaches from ALICE and LHCb have in common?
- Short introduction to GPUs:
  - Why should we use GPUs and what can we gain?
- The experience of ALICE
- The experience of LHCb
- Conclusion

# COMPUTING FOR THE UPGRADES OF THE LHC EXPERIMENTS

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## LHC Upgrade Schedule





- LS2 LHC upgrade: heavy ion rate: >50 kHz in Run 3 (>10 kHz now), boost pp collision rate by small factor.
- LS3 LHC upgrade: HL-LHC era, boost pp collision rate by factor 5 7 in Run 4.
  - Highest pp luminosity only for ATLAS and CMS their detectors are upgraded for Run 4 accordingly.
  - ALICE and LHCb perform a major upgrade for Run 3 now.
- Run 3 is adiabatic increase for ATLAS / CMS, with no increase for ALICE.

## **Online / Offline Computing in ALICE / LHCb in Run 3**



#### LHCb

- First phase of trigger (HLT1) during data taking.
- Second phase of trigger (HLT2) when there is no beam.

#### Run 2



#### **Common strategy:**

- 2 phase processing with disk buffer
- Full processing in software



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## **Online / Offline Computing in ALICE / LHCb in Run 3**





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## **Comparison of processing, data rates and sizes**



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	AL	ICE (PD-PD)		LHCD	AILAS		CMS		
	Run 2	Run 3	Run 2	Run 3	Run 2 / 3	Run 4	Run 2 / 3	Run 4 (PU 140/200)	
Luminosity	~10 kHz	50 kHz	4*10 <sup>32</sup>	2*10 <sup>33</sup>	2.14*10 <sup>34</sup>	5-7.5 *10 <sup>34</sup>	2.14*10 <sup>34</sup>	5–7.5 *10 <sup>34</sup>	
Hardware trigger	500 Hz – 2 kHz	50 kHz continuous	1 MHz	- / Full 30 MHz bunch crossing rate	95 kHz	1 MHz (can evolve to 4)	100 kHz	500 / 750 kHz	
HLT Accept	No rejection	No HLT	12.5 kHz	>100 kHz	1 kHz (< 2)	10 kHz	1 kHz	5 / 7.5 kHz	
Raw Data Rate into HLT	45 GB/s (w. ZS)	3 TB/s (w.o. ZS)	55 GB/s	4 TB/s (w. ZS)	29 GB/s (260 GB/s L1)	2.6 TB/s (5.2 TB/s L1)	1.6 TB/s (event network)	23 / 44 TB/s (event network)	
Data stored	~10 GB/s	Up to 100 GB/s	0.6 GB/s	2-10 GB/s	2.4 GB/s	50 GB/s	5 GB/s	32 / 61 GB/s	
Data Buffer	~1 PB DAQ buffer to Tier0	~60 PB (one year of compressed data), up to 100 GB/s	~12 PT	~100 PB (two weeks of HLT1 accepted raw data, 150 + 150 GB/s read/write.	1.5 TB events + 48 hours to Tier 0	36 PB, 48 hours + L1 to HLT	12 TB (RAM disk, events before HLT, 60s)	171 / 333 TB (events before HLT, 60s)	

ALICE will take 100x more events, but only minimum bias, all other experiments collect 10x more statistics.

- ALICE investigates the option to run in a software triggered mode at higher rate during some time of a year.
- ALICE and LHCb process all data in software, use large disk buffers to hold large amount of (compressed raw data) for processing in the online farm when there is no beam.
- ATLAS and CMS have much higher luminosity, full readout of front-end at bunch-crossing rate and processing in software not feasible and not cost effective.
- ALICE features high data rate during Pb-Pb (due to TPC), collects large amount of data in only few weeks.

## **Comparison of ALICE and LHCb data processing in Run 3**



Similar

- Both experiments do full two-stage online processing at bunch crossing rate, with disk buffer and offline quality output.
  - Input data rates: >3TB/s for both zero-suppressed for LHCb (factor 4) raw for ALICE with zero-suppression in FPGA down to ~1TB/s.
  - Calibration: Full calibration available for Asynchronous stage / HLT2.
    - Online calibration with feedback loop tested in the ALICE HLT in Run 2, under study for Run 3, Velo alignment calibration in LHCb HLT1.
  - Event building: A set of input nodes (FLP / DAQ) receives the detector links via PCIe40 FPGA card.
    - ALICE sends time frames (TF) of 23ms from the FLP to Event Processing Nodes, where events are merged, build, and reconstructed.
    - LHCb first builds the events internally inside the DAQ via a fast network, then ships them to the Event Filter farm via a broad network.
      - Network transfers synchronized via software to avoid congestion.
      - Many events are coalesced (similar to TF).
      - Could switch to similar event building as ALICE if needed.
  - Cluster: Input: ~170 DAQ nodes @ LHCb, ~200 FEP nodes @ ALICE
     Processing: ~2000 Event Filter nodes @ LHCb, ~1500 GPUs in EPNs @ ALICE.
  - Disk buffer: LHCb buffers up to 3 weeks, exploits turnaround, TS, MD periods, runs MC at YETS. ALICE buffers 1 year, exploits also YETS.
    - ALICE compresses data in synchronous stage, compressed raw data stored to disk buffer and to tape is identical.
      - ALICE has highest data taking output rate of up to 100 GB/s, but only during beginning of Pb-Pb fill.
      - Pb-Pb data parked for processing on disk buffer for 1 year. Compute budget for 2 full reconstructions passes within the year.
      - pp data taking and asynchronous reconstruction in between / in parallel.
      - Part of asynchronous reconstruction can run on GRID.
      - Online farm must be capable to process this data in the synchronous stage.
    - LHCb stores raw data after HLT1 trigger to disk buffer, second trigger rejection in HLT2.
      - LHCb has higher data rate from HLT1 to disk buffer, but not for final storage.



# **POTENTIAL OF GPUS**

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## Introduction



- Moore's Law: ٠
- Manufacturing ٠
- size, frequency, •
- [GFlop/s] and performance ٠
- grow exponentially. Frequency began to stagnate 2003. [H5] / Land ٠
- ٠
- ٠





### Introduction



- Moore's Law: ٠
- Manufacturing ٠
- size, frequency, •
- [GFlop/s] and performance ٠
- grow exponentially. Frequency began to stagnate 2003. GPUs are faster than CPUs. ٠
- ٠
- ٠
- ٠
- ٠







- GPUs use their silicon for ALUs
- CPUs use their silicon mainly for caches, branch prediction, etc.





#### Intel Nehalem

**NVIDIA Kepler** 





- CPUs are designed for fast execution of serial programs.
  - Clocks have reached a physical limit.
    - → Vendors use parallelization to increase performance.
- GPUs are designed for parallel execution in the first place.
  - The "only" limit for GPU performance is heat dissipation.
  - GPU clocks are usually lower than they could be.
    - This saves power
    - Hence more hardware can be powered in parallel
      - → Better overall performance





## NVIDIA GTX280 GPU



## **GPU Programming example (stupid addition of 2 vectors)**





## **GPU Programming example (stupid addition of 2 vectors)**





## **TPC Tracking performance**





- Red curve: algorithm speed-up.
- Other curves: GPU v.s. CPU speed-up corrected for CPU resources.
  - How many cores does the GPU replace.
- Significant gain with newer GPU (blue v.s. green).
- GPU with Run 3 algorithm (2) 15 replaces > 800 CPU cores (3) Running Run 2 algorithm. (3) 10 (blue \* red). (at same efficiency / resolution).5
- We see ~30% speedup with new GPU generation (RTX 2080 v.s. GTX 1080)



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# EXPERIENCE OF THE ALICE EXPERIMENT

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## ALICE in Run 3: 50 kHz Pb-Pb

ALICE

- Record large minimum bias sample.
- All collisions stored for main detectors  $\rightarrow$  no trigger.
- Continuous readout  $\rightarrow$  data in drift detectors overlap.
- 100x more events, 100x more data.
- Cannot store all raw data  $\rightarrow$  online compression.
- $\rightarrow$  Use GPUs to speed up online processing.

- Overlapping events in TPC with realistic bunch structure @ 50 kHz Pb-Pb.

- Timeframe of 2 ms shown (will be 10 20 ms in production).
- Tracks of different collisions shown in different colors.

## Step 1 (Seeding)



Step 1: Combinatorial seeding Strategy: deal with the combinatorics as early as possible. Searches for three clusters composing straight line Seed everywhere, each track has at least • **Concatenates straight lines** some seedable part, no need to be 100% Only step with non-linear runtime. efficient. row r + 2row r + 4 $dx_{+}$ row r + dy\_ row r row r dx row r - 2 C row r - 2 11.12.2019 20 David Rohr, drohr@cern.ch

## Step 2 (Track Following)



- Step 2 (Simplified Kalman Fit):
- Track parameters are fit to the seed.
- Trajectory is extrapolated to adjacent TPC row.
- Cluster closest to extrapolated position is found.
- Fit is improved with new cluster.



## **TPC Data Compression**



**Unassigned clusters** 

**Removed Clusters** 

**Reconst/ucted Tracks** 

**Fit failed** 

- TPC Data compression involves 3 steps:
  - 1. Entropy reduction (Track model, variable precision, etc.)
  - 2. Entropy encoding (Huffman, Arithmetic, ANS)
  - 3. Removal of tracks not used for physics.
- Steps 1 + 2 implemented for Run 2.
  - Current compression factor 8.3x.
- Prototype for Run 3 achieves factor 9.1x (TDR assumed 10x).
- Step 3 must close the gap to the required compression in Run 3.
  - Remove clusters from background / looping tracks.
    - Adjacent to low- $p_{\rm T}$  track < 50 MeV.
    - Adjacent to secondary leg of low- $p_{T}$  track < 200 MeV.
    - Adjacent to any track with  $\varphi > 70^\circ$  in the fit.
  - Protect clusters of physics tracks.
    - Not Adjacent to any physics-track (except  $\varphi$  > 70°).
- In addition:
  - Use reconstructed track quantities to reduce entropy.

**Noisy TPC pads** 

## **Online / Offline Computing in ALICE in Run 3**





## **Tracking in ALICE in Run 3**

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Bulk of computing workload:

#### **Synchronous**

- >90% TPC tracking / compression
- Low load for other detectors

#### Asynchronous

- TPC among largest contributors
- Other detectors also significant



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#### ALICE GPU processing strategy

Baseline solution (almost available today): TPC + part of ITS tracking on GPU

Mandatory solution to keep up with the data rate online.

Defines number of servers / GPUs.

#### Optimistic solution (what could we do in the ideal case): Run most of tracking + X on GPU.

- Extension of baseline solution to make best use of GPUs.
  - Ideally, full barrel tracking without ever leaving the GPU.
  - In the end, we will probably be somewhere in between.

Asynchronous phase should make use of

Available in the O<sup>2</sup> farm anyway.

Future HPC / grid sites may have GPUs.

(inner tracking system)

152 pad rows TPC

nsition radiation det

the available GPUs.



- Status of reconstruction steps on GPU:
  - All TPC steps during synchronous reconstruction are required on the GPU.
  - Synchronous ITS tracking and TPC dE/dx in good shape, thus considered baseline on the GPU.
  - Remaining steps in tracking chain part of optimistic scenario, being ported step by step to GPU.
    - Porting order follows topology of chain, to avoid unnecessary data transfer for ported steps current blocker is TPC ITS matching.









- Status of reconstruction steps on GPU:
  - Different reconstruction steps enabled in synchronous and asynchronous reconstruction.



TRD tracking / TPC calibration: see poster of Ole Schmidt Space point calibration of the ALICE TPC with track residual



- Status of reconstruction steps on GPU:
  - Different reconstruction steps enabled in synchronous and asynchronous reconstruction.





# TRD tra Space

#### Status of reconstruction steps on GPU:

- TPC steps during synchronous reconstruction are required or Strategy:
- Start with standalone TPC and ITS tracking.
  - Standalone ITS tracking needed since TPC tracks lack absolute time.
    - nin ITS tracking uses vertexer as first step. cessary data transfer for ported steps current blocker is TPC ITS matching
      - TPC tracking has no vertex constraint, starts with segment tracking in individual TPC sectors, than merges the segments and refits.
- ITS and TPC tracks are matched, fixing the time for the TPC.
- The afterburner propagates unmatched TPC tracks into the ITS and tries to find matching hits of short tracks not found in ITS standalone tracking.
- Tracks are extrapolated outwards into the TRD, once the time is fixed.
  - TRD standalone tracking and matching (like for ITS) is less efficient due to many fake TRD tracklets.
- Optionally, after TRD tracks can be extrapolated to TOF.
- Global refit uses the information from all detectors.
- V0 finding
- In the synchronous phase, the TPC compression chain starts after the TPC standalone tracking in parallel:
  - Clusters not used in physics are removed, depending on the strategy (see later) this might require extra steps for identification and rejection of very low p<sub>T</sub> clusters below 10 MeV/c.
  - Track model (and other steps) reduce the entropy for the final entropy encoding.

Components:

 Final entropy encoding using ANS. Not clear yet whether this will run on GPU efficiently. Alternatively, transport entropy-reduced clusters to host and run entropy encoder there.

Sorting

ation

lemorv Reuse

## Approach if Run 2 HLT TPC / ITS Tracking Components





## **Approach for Run 3**





## **Approach for Run 3**







- ALICE reconstructs timeframes (TF) independently (~10 ~20 ms; 128 256 orbits; ~500 ~1000 collisions).
  - One TPC drift time of data not reconstructible at TF border (~ 90 us) → < 1 % of statistics lost (< 0.5 % for 20 ms).
  - Timeframe should fit in GPU memory. If not, could use kind of ring buffer, or reduce TF length to 128 orbits.
- Trying to avoid the ring buffer approach, could be added later if needed.
- Custom allocator: grabs all GPU memory, gives out chunks manually, memory will be reused when possible.
  - Classically: reuse memory between events, collisions are not that large.
  - ALICE reuses memory between different algorithms in a TF, possibly also between independent collisions.
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Pers	stent da	ata	1				1	Non-p	ersistent scratcl	h data for	algorithms	Non-persisting i	nput data
TP0 Hits	C TF 1 Hit	PC is 2	TPC Hits 3	TPC Hits 4	ITS Hits	TPC Tracks	ITS Tracks	Matches	lemory	Scratch			TPC Raw 1
							TPC I <sup>-</sup> Matchi	TS ng			F be	Preload TPC raw data of next TF fore current TF is finished.	

#### Work in Progress



- ALICE reconstructs timeframes (TF) independently (~10 ~20 ms; 128 256 orbits; ~500 ~1000 collisions).
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Persiste	nt data	3				- And	Non-p	ersistent scratch data for algorithms	Von-persisting in	put data		
TPC Hits 1	TPC Hits 2	TPC Hits 3	TPC Hits 4	ITS Hits	TPC Tracks	ITS Tracks	Matches	lemory	TPC Raw 2	TPC Raw 1		
Estimated maximum memory needed during important for 10 ms TF (*2 for 20 ms):												
•	TPC Clu	ster finde	er: ~	3 GB	( + input /	scratch d	ata, which	is pipelined)				
•	TPC Transformation: 12.1 GB											
TPC Sector tracker: ~ 14.6 GB (including persistent memory												
	TPC Me	rger / tra	ck fit:	14.1 GB from previous steps)								
	TPC Cor	mpressic	on:	12.9 GB								
<ul> <li>Later steps do not scale their scratch memory with TPC input → less memory intensive.</li> </ul>												
→ 16 GB GPU will suffice for 10 ms TF (unclear for 12 GB after optimizations).												
8 GB insufficient for 10 ms TF, 20 ms TF needs 32 GB, alternatively ring buffer.												

## **Performance Tuning**



- GPUs of different vendor's / generation's might favor different tuning.
  - Many algorithms have tunable parameters (for processing speed).
  - We implemented most features such, that they can be switched off.
    - Worst case, at compile time via preprocessor definition.
- One example: Distribution of tracks among GPU threads during track following:
  - Illustration of active GPU threads over time (time on y-axis).



- Blue : Track Fit
- Green : Track Extrapolation





- Number of average idle threads reduced by factor ~3, but large overhead for rescheduling.
- Yields ~50% speedup on some GPUs, but becomes even slower on others.
- For new GPUs:
- Run a benchmark with a parameter range scan to find best settings.
- After 3 iterations (GPU generations), we got good results out of the box.

## **Performance Tuning**



- Handling of asynchronous computation / data transfers
  - 1<sup>st</sup> iteration (Run 1 HLT): Split event in chunks, to pipeline CPU processing, GPU processing, and PCIe transfer.

DMA			STATE OF	
GPU				
CPU 1				
CPU 2				
CPU 3				
		Time		
outine:	Initialization 📃 Neighbor Finding	Tracklet Construction	Tracklet Selection	Tracklet Output

- 2<sup>nd</sup> iteration (Run 2 HLT): Processing of two events in parallel on the GPU concurrently.
  - ~20% faster than first version GPUs have become wider and this exploits the parallelism better.
  - Not possible during Run 1 due to GPU limitations at that time.
  - We still kept the pipeline-scheme within each event, to maximize performance.
- 3<sup>rd</sup> iteration (Run 3): Go back to the old scheme from Run 1 with time frames instead of events.
  - Time frames are large  $\rightarrow$  avoid keeping multiple in memory.
  - Enough parallelism inside one time frame.

## **Compatibility with several GPU frameworks**



- Generic common C++ Code compatible to CUDA, OpenCL, HIP, and CPU (with pure C++, OpenMP, or OpenCL).
  - OpenCL needs clang compiler (ARM or AMD ROCm) or AMD extensions (TPC track finding only on Run 2 GPUs and CPU for testing).
  - Certain worthwhile algorithms have a vectorized code branch for CPU using the Vc library.
  - All GPU code swapped out in dedicated libraries, same software binaries run on GPU-enabled and CPU servers.



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