

ATLAS WORLD-cloud and networking in PanDA

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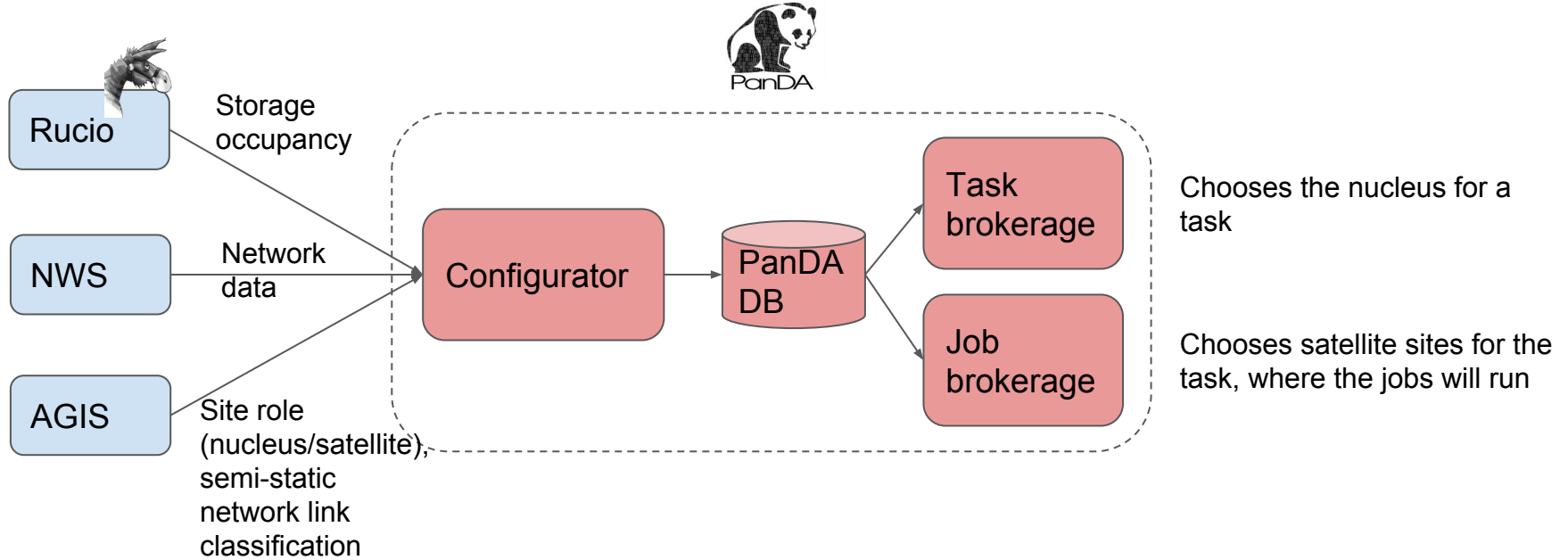
Rationale

- Original ATLAS Computing Model was designed as static clouds (=mostly national or geographical groupings of sites), setting data transfer perimeters
- Hierarchic model with clear distinctions in Tier 1-2-3 level
- The particular policies enforced in the workload management system
 - Output of tasks (=set of jobs) had to be aggregated in the O(10) Tier 1s
 - Tasks have to be inflexibly executed within a static cloud
- This model works, but is getting outdated and has a series of disadvantages
 - WLCG networks have evolved significantly in the last two decades and bandwidth has increased O(1000)
 - Limiting transfers within a cloud is no longer needed
 - Usage of sites is uneven. In particular Tier 2 storage was not optimally exploited
 - Small clouds stuck with large, high priority tasks
 - First try was to allow sites to belong to multiple clouds, now we are completely breaking the boundaries

WORLD cloud

- Dynamic model, where tasks are not confined to a cloud. A grouping of processing sites is defined dynamically for each task
- **Task nucleus:**
 - Task brokerage will choose a nucleus for each task based on data locality, queued work and available storage
 - The capability of a site to be a nucleus is defined manually in AGIS (ATLAS Grid Information System). Tier 1s and the bigger Tier 2s are defined as nuclei
 - The output will be aggregated in the task nucleus
- **Task satellites:**
 - Run jobs and ship the output to the nucleus
 - Job brokerage selects satellites for each task, based on usual criteria (e.g. number of jobs and data availability)
 - Satellites are selected across the globe: a network weight will match well connected nuclei and satellites

Configurator



Network data

- Rucio, FAX, PerfSonar events are collected in the ATLAS analytics platforms [1]
- The Network Weather Service[2] (DDM team) aggregates information from the platform.

Per source-destination pair:

- #files transferred in last hour
- #files queued
- Throughput according to FTS, based on 1 week data
- Throughput according to FAX
- PerfSonar metrics (latency, packet loss, throughput)
- AGIS also provides semi-static link classification to be used as a backup
- Configurator agent downloads and processes this information every 30 min. Data is cached in a key-value table in PanDA DB
 - Table is extensible for any new metric without modifying the structure

Presentations in this conference:

[1] Ilija Vukotic et al., “Big Data Analytics Tools as Applied to ATLAS Event Data”

[2] Mario Lassnig et al., “Using machine learning algorithms to forecast network and system load metrics for ATLAS Distributed Computing”

Task brokerage: nucleus selection

- One nucleus is chosen for each task
- Nuclei must fulfil all conditions:
 - Be in active state and be able to execute the workload
 - Have over 5TB free space - the free space includes an estimation of the space to be filled by pending tasks
- The number of output files transferring to the nucleus must be below 2k
- Nuclei compete through a combined weight, based on data locality, total RW (remaining work) and available storage size in the nucleus.

$$RW = (nEvents - nEventsUsed) \times cpuTime$$

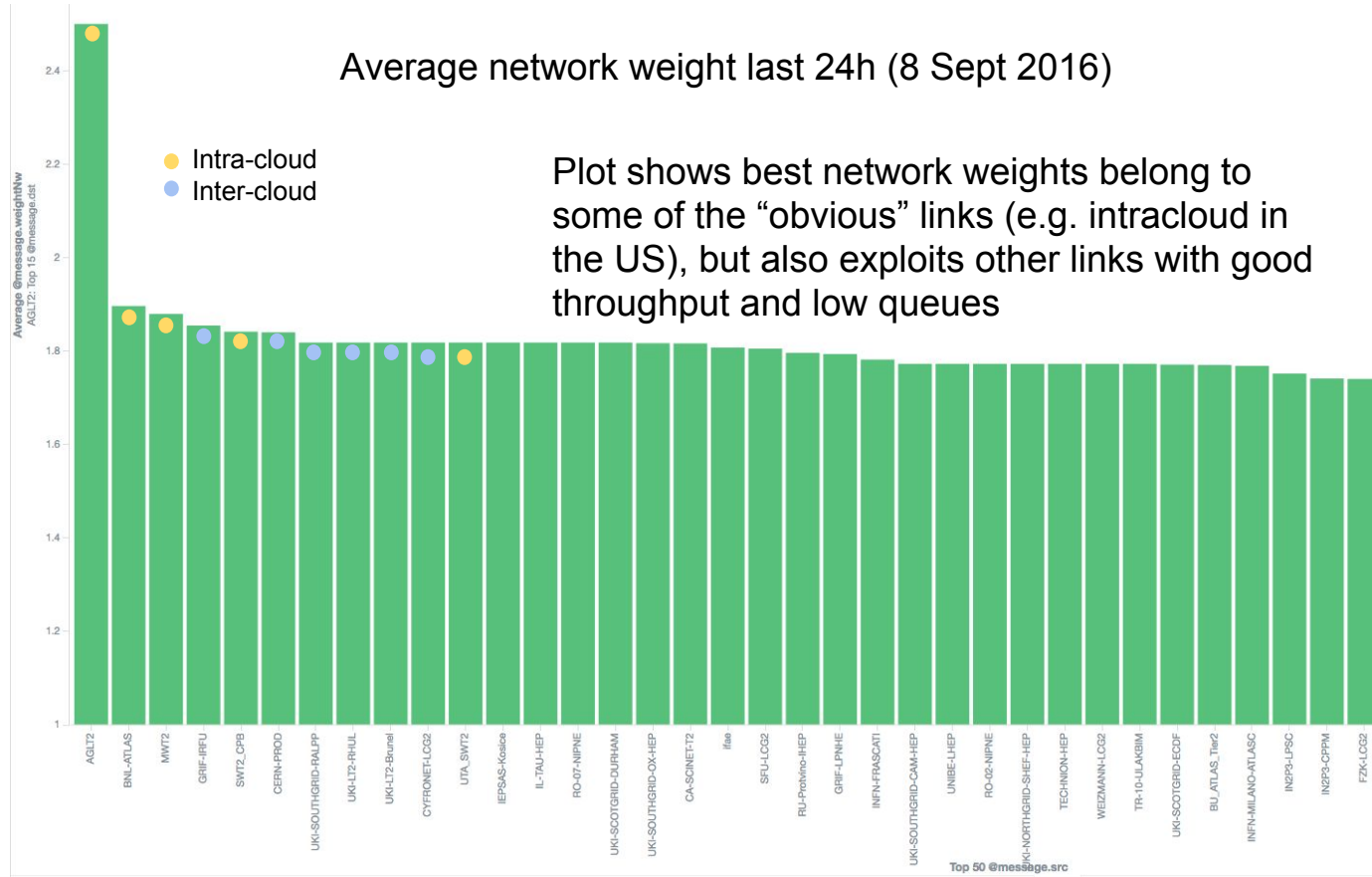
$$weight = \frac{1}{total\ RW} \times \frac{availableInputSize}{totalInputSize} \times tapeWeight \times \frac{freeSpace}{totalSpace}$$

Job brokerage: satellite selection

- Up to 10 satellites are selected to execute the task
- The satellites must be able to run the jobs (RAM, walltime, core count, queues, SW releases and other settings)
- Sites must have <150 files in the transfer queue to the nucleus
- Eligible sites compete on a weight basis (#jobs, available & missing data)
- We have added a new network weight, looking for balance between good throughput and queue length of output files
 - Throughput calculated from FTS transfers in the last hour
 - Queued files in Rucio/FTS. In practice the number of files queued for a link is very low; having $O(100)$ files starts to indicate the channel is not keeping up

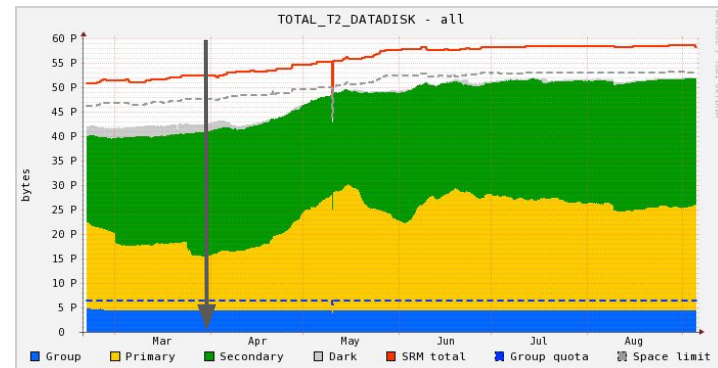
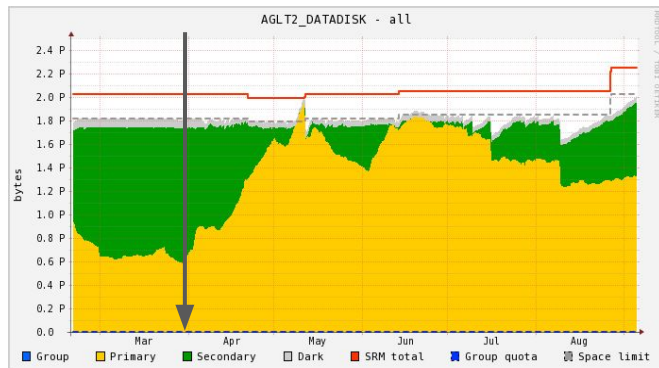
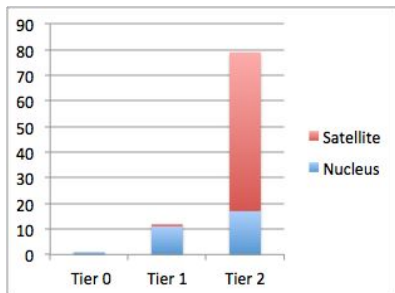
Example: Top connected sites to Nucleus **AGLT2** (Michigan)

Average network weight last 24h (8 Sept 2016)



Status: Impact on T2 disk space usage

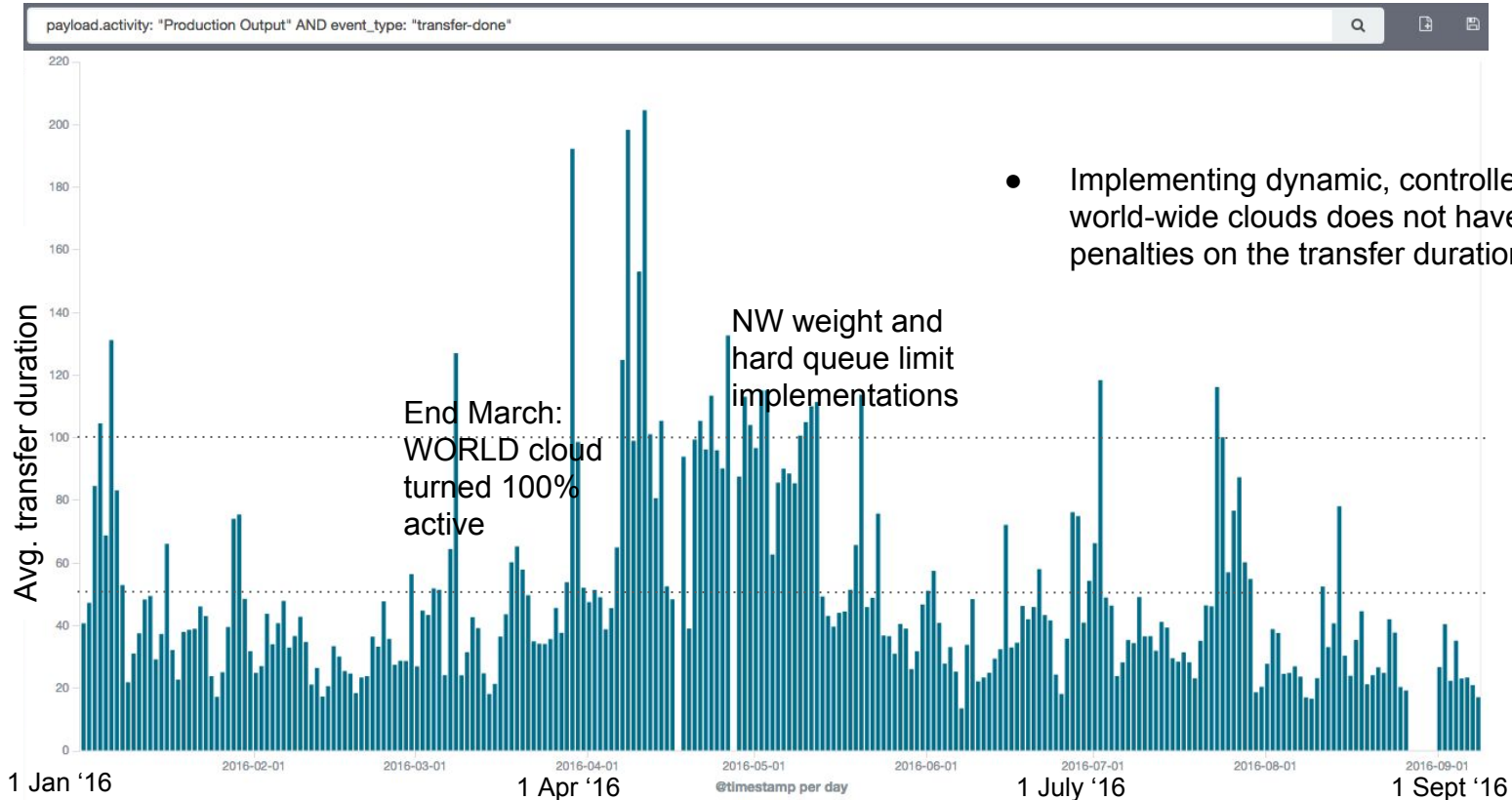
- WORLD cloud was fully activated end March 2016
- Nuclei being added progressively
 - Currently T1s and ~20% of T2s (tentatively more T2s will be added)
- Extending task output recipients to T2s is starting to have a positive impact on the overall disk usage (more primary data on T2s)



Storage plots from David Cameron's space monitoring:

<http://adc-ddm-mon.cern.ch/ddmusr01/plots/>

Impact on output file transfer duration



Observations and future work

- Some sites suffered initially under transfer load during heavy campaigns
 - Not because of inter-cloud transfers, but because they had not enough bandwidth for the Nucleus role
- Hard limit queue controls worked fine to alleviate the issue and deviate the traffic from blocked sites
 - Limits trigger also during unrelated, “accidental” massive transfers
- Further downstream controls could be implemented to e.g. avoid already assigned jobs to run while their nucleus is stuck
 - Pause overloaded nucleus in job brokerage (using a higher queued file threshold to avoid waves)
- Reduced operational effort to rebroker tasks with more resources
- Need to improve analytics data for gridwise analysis
- We have tried to optimize output file transfer, but still need to include some optimization for input file transfers
 - This case needs to be solved together with the DDM team, since it involves further uncertainties (multiple copies, tape staging, etc)