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A new Variable Resolution Associative Memory for High Energy Physics

Data selection at hadronic colliders

- Search for rare SM or predicted BSM processes push the colliders intensity to new frontiers
 Rare processes are overwhelmed by well-known processes
 - 8/9 order of magnitudes between Higgs and total cross-section
- Need to prioritize the physics output and leave flexibility for the unexpected
 Need sophisticated techniques to maximize the bandwidth for interesting events
 Several strategies are used, track trigger
 showed good results at CDF (SVT), ATLAS
 implementation (FTK) can have similar success



Online tracking with FTK



PATTERN 1 PATTERN 2 PATTERN 3 PATTERN N PATTERN N PATTERN 4 PATTER

An associative memory (AM) chip allows a full parallel search in the pattern bank. Similar to commercial CAMs.

 $p_i = \sum_j C_{ij} \cdot x_j + q_i$

- In FTK tracking divided into two separated steps
- The pattern matching uses a large bank of patterns and special hardware
- The track's parameters are evaluated using a linear Principal Component Analysis algorithm (j.nima.2003.11.078)

The hits in a road are combined and the parameter evaluation is reduced to a linear problem, efficiently solved in an FPGA

ATLAS Inner detector layout



- FTK processor will read the ID for each Level 1 accept
- To improve the input bandwidth the ID is divided into 8 regions
 - Each region will have 16 processing units, all working in parallel
 - A total of about 1 Giga patterns
- The device is able to produce a complete list of tracks in tens of microseconds

Associative memory and granularity



³ patterns , 3 roads (1 fake), 4 tracks (2 fakes)



- A large segmentation (SS) improves the pattern recognition efficiency
 - More likely that noisy hits enter the road: fake tracks
 - Or random combinations seem to be aligned in pattern: fake roads
- Reducing the segmentation has benefits
 - Fake roads or tracks are less probable
 - But more hardware, more money

4 patterns, 2 roads, 2 tracks, no fakes

SS size, efficiency and bandwidth



Roads

Tree search processor idea

- Patterns can be generated at the finest resolution and grouped according to the low resolution granularity
 - A tree organization is natural
- AM chip size is a function of the low resolution granularity
 - Limited number of patterns
- Efficiency and fake rate driven by the granularity in the leaves
 - Final pattern recognition uses the full precision
- AM patterns usually have a limited number of leaves
 - Most of the internal volume unused



TSP rejection power

- TSP simulation implemented
 - The AM chip simulation finds the low resolution roads
 - For each AM road all the nodes are compared looking for matches
- WH events with 75 pileup-event simulation used
- Simulation results confirm TSP can reject large number of roads
 - Most of the fired roads have few subpatterns
 - Special attention to this category is justified



Variable resolution and don't-care bit

- To have a more flexible AM bank the SS are stored at maximum precision
 - Each additional bit halves the size of the SS
 - An additional bit is set if the match stops before the least significant bits, not using the full precision: "don't care" bit like modern CAM memories (K. Pagiamtzis and A. Sheikholeslami, Solid-State Circuits, IEEE Journal of, vol. 41, no. 3, 2006)
 - Where the full precision is used (no DC) a portion of the SS is vetoed
 - Where the DC is set the whole SS is active
- The use of DC allows a reduction in the active region of the patterns
 - This reduces the probability of fake roads and tracks
 - Reducing the number of roads in the output

Pattern in AM chip w/o the DC



Pattern in AM chip w/ the DC

Complete simulation studies

- FTK simulation improved to include DC and TSP
- TSP bank produced using 7 inner detector layers
- Thin patterns grouped within the large patterns doubling the SS size: 1 DC bit
 - Working point fixed at 90% TSP efficiency
 - Typical TSP \rightarrow AM reduction 1/3
 - Increment of area used in the AM chip 1/15=7%

- Physics sample used: WH with full ATLAS simulation
 - Multi-jet events with high number of tracks
- 3 pileup scenarios at 25 ns bunch spacing are evaluated:
 - 17.6, 40.5 and 75 pileup events
- For each pileup scenario the number of roads per processing unit is important
 - Number extrapolated from barrel only simulation



Road flux Bank 96x80x72 Road flux Bank 48x40x36 200000 12000 125 pilleup event 10000 150000 27.6 DM 8000 # roads 50% 100000 # roads 95% 6000 4000 50000 2000 0 0 AM TSP AM TSP AM@DC AM@TSP AM@DC AM@TSP Median •AM size 2 MP •AM size 32 MP TSP bank 48x40x72 TSP bank 48x40x72 •TSP size 5 MP

•TSP size 100 MP

Performance (2)

Roads flux Bank 12x10x36



Number of output roads Pileup AM AM w/ AM w/ TSP AM bank DC TSP size events 17.6 2 M 1060 5040 959 709 3880 6500 32 M 40 37000 5720 138 M 8250 75 53500 4430 5950

- Pattern bank reduction factor: 2.5 3
- A bank with DC capability reduces the fakes by a large factor: 5-7
- Good performance without adding further hardware

Conclusions

- Rapid tracking is an opportunity for the ATLAS experiment
- FTK based on a successful technology (SVT)
 - But moving from 2 TeV@10³² to 14 TeV@10³⁴ needs new ideas
- Variable resolution is a general idea
 - The patterns are able to change in shape and matching volume
 - Improves the precision only where needed

- Just using a single DC bit has a significant effect
- The number of roads out of the AM chip is reduced greatly by the DC
 - It has almost the same effect as a complete TSP
- 7% more information in the bank has the same power as a bank 3 times bigger
 - Use of several DC bits possible in the chip and a study is ongoing



FTK and ATLAS DAQ



ROD output sent to FTK at L1 rate

FTK:

- Operates in parallel with silicon readout after each L1 trigger
- Reconstructs tracks over entire inner detector (|η|<2.5) in ~25µs (for 3x10³⁴ cm² s⁻¹ Luminosity)
- Provides high-quality tracks by the beginning of L₂ trigger processing

Both trigger levels benefit:

- L1: events needing L2 tracking could have a higher fraction of the L1 rate
- L2: with tracking already performed, the extra CPU time can be used for more advanced algorithms

DC bit content construction

- After the construction of the tree the content of the AM patterns is used
- If in one layer all the nodes shares the same portion, the DC is not set
 - Shrinking the active part of the AM pattern
- If all the portions are used in the nodes, the DC is set

The smaller the number of TSP nodes, smaller the use of DC

