

Intro to Data Structures

Lecture #26 – Graphs
November 30, 2014

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Outline for Today

- HW 6 issues
- 2/3 of course grade is behind you...
 - HW7 (7%)
 - Q7 (1.5%)
 - Final (25%)
- Graphs

HW6 issues

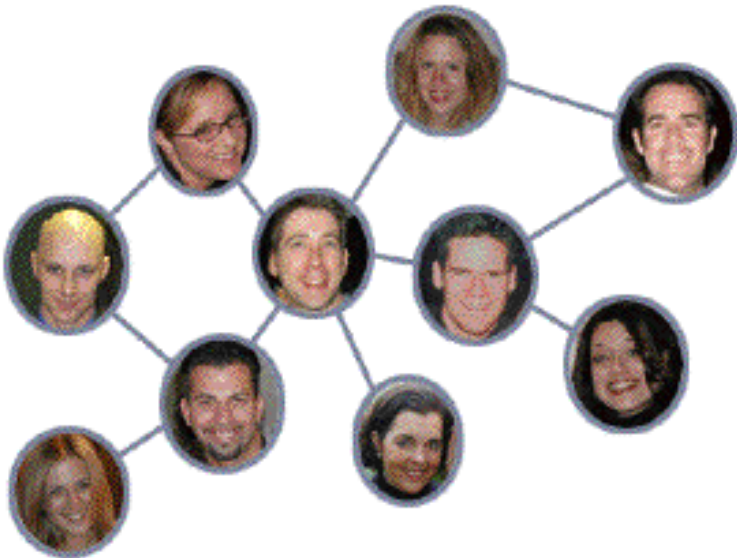
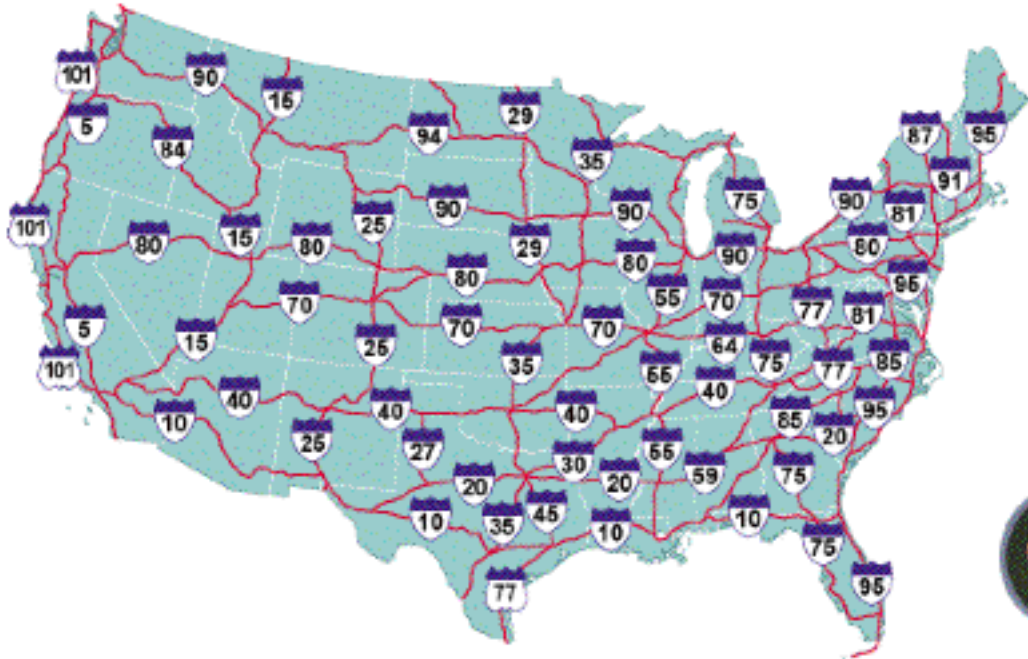
- Helper functions should be private
- *contains()* call in add method??
- $\text{if } ((\text{right} - \text{left} == 1) \parallel (\text{right} - \text{left} == -1) \parallel (\text{right} - \text{left} == 0)) \rightarrow \text{Math.abs}(\text{right} - \text{left}) \leq 1$
- *isBalanced()* needed to check subtrees as well!
- *mirrorTree()* was not supposed to damage/alter the original tree!!

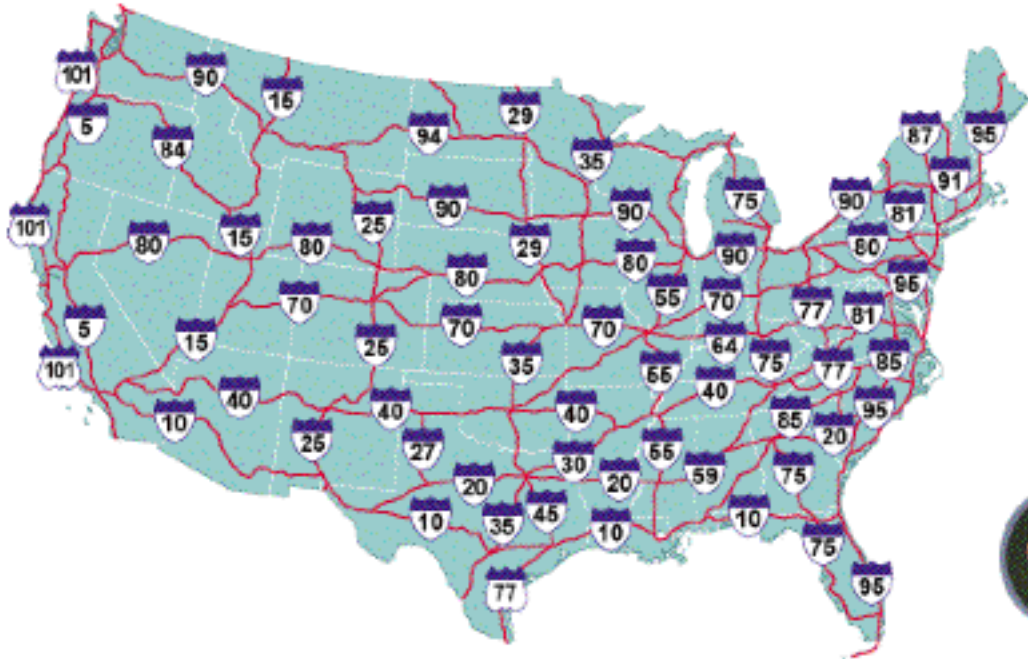
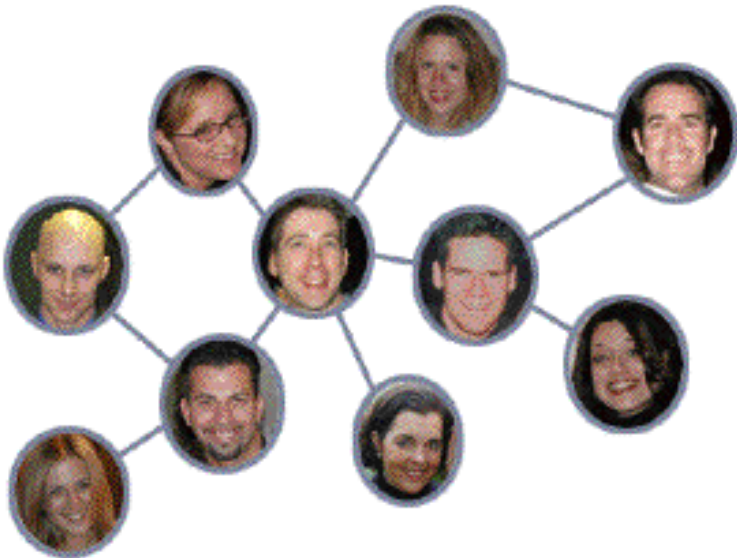
HW6 issues

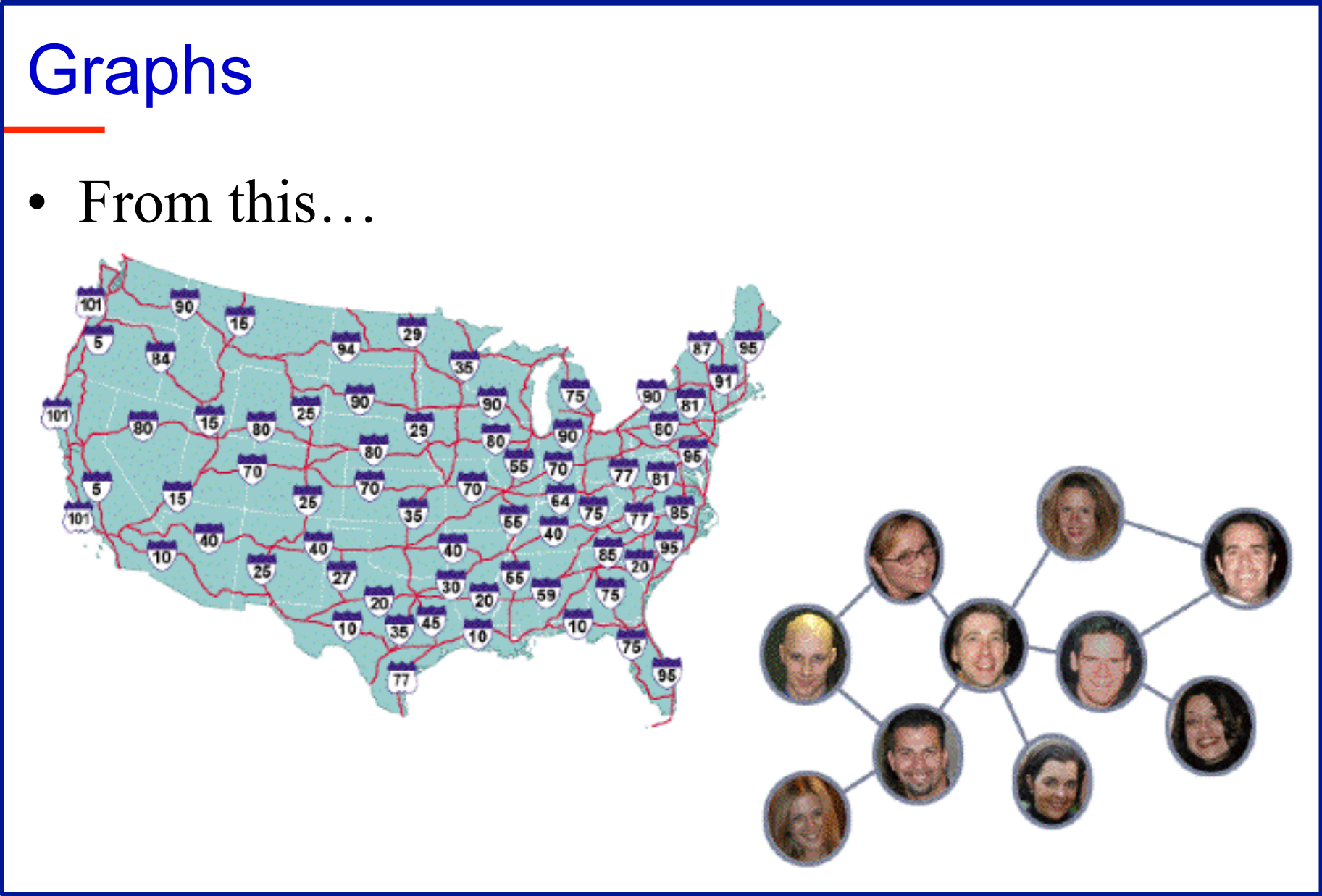
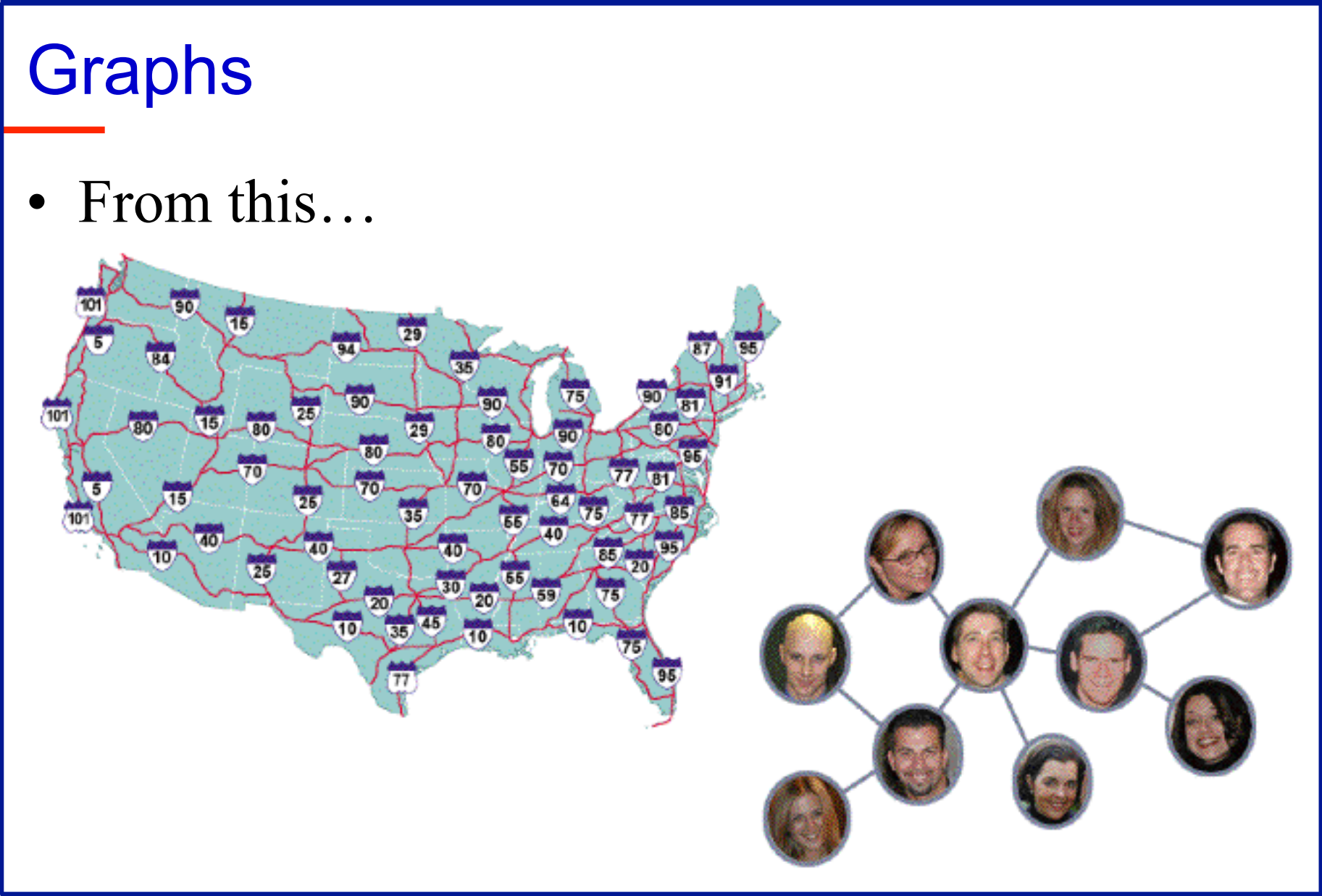
- No statistical output; mine looked like
Total number of words read: 172715
Number of words inserted (length ≤ 7): 51913
Number of nodes in the tree: 41121
- Many of yours looked like
- Test, test, test...

Graphs

- From this...

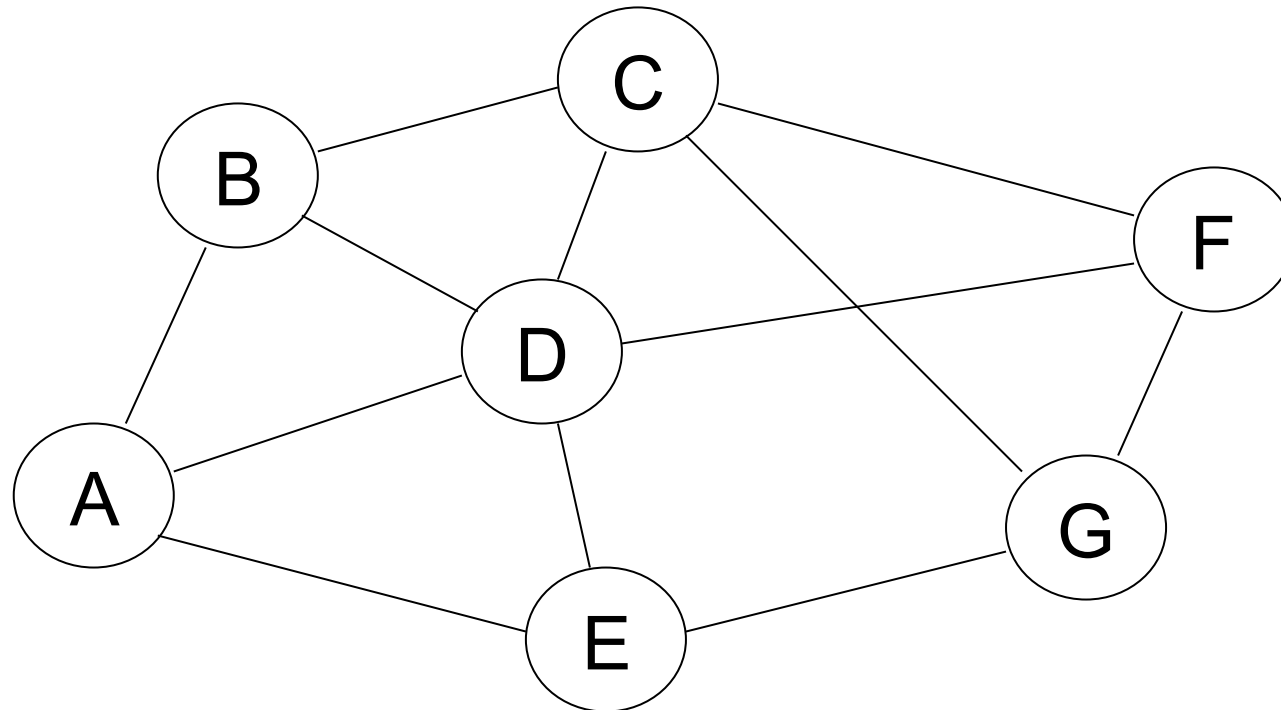


- # Graphs
- From this...
- 
- 



Graphs

- To this...



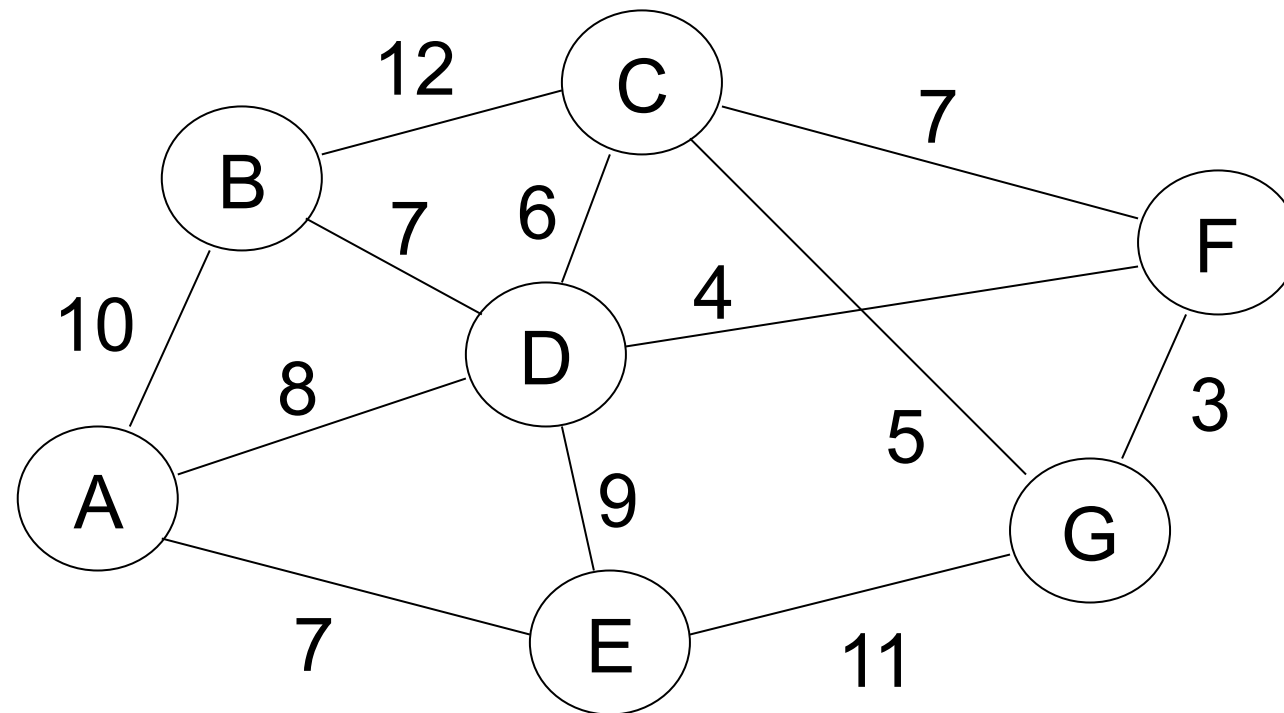
Graphs

- A graph is a collection of vertices (nodes) and edges (connecting the nodes); $G = \langle V, E \rangle$
- Edges can be
 - directed/undirected
 - non-weighted/weighted (have different values)
- A *path* is a sequence of vertices connected by edges; the length of the path is the number of edges in the path
- *degree*: number of edges incident to a vertex

Graphs (facts)

- An undirected graph with n vertices will have at most ? edges
 - $\binom{n}{2} = n * (n-1) / 2$
- A directed graph with n vertices will have at most ? Edges
 - $n * (n-1)$

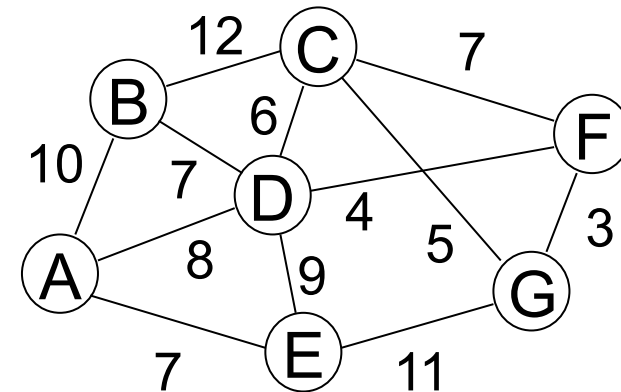
A weighted graph



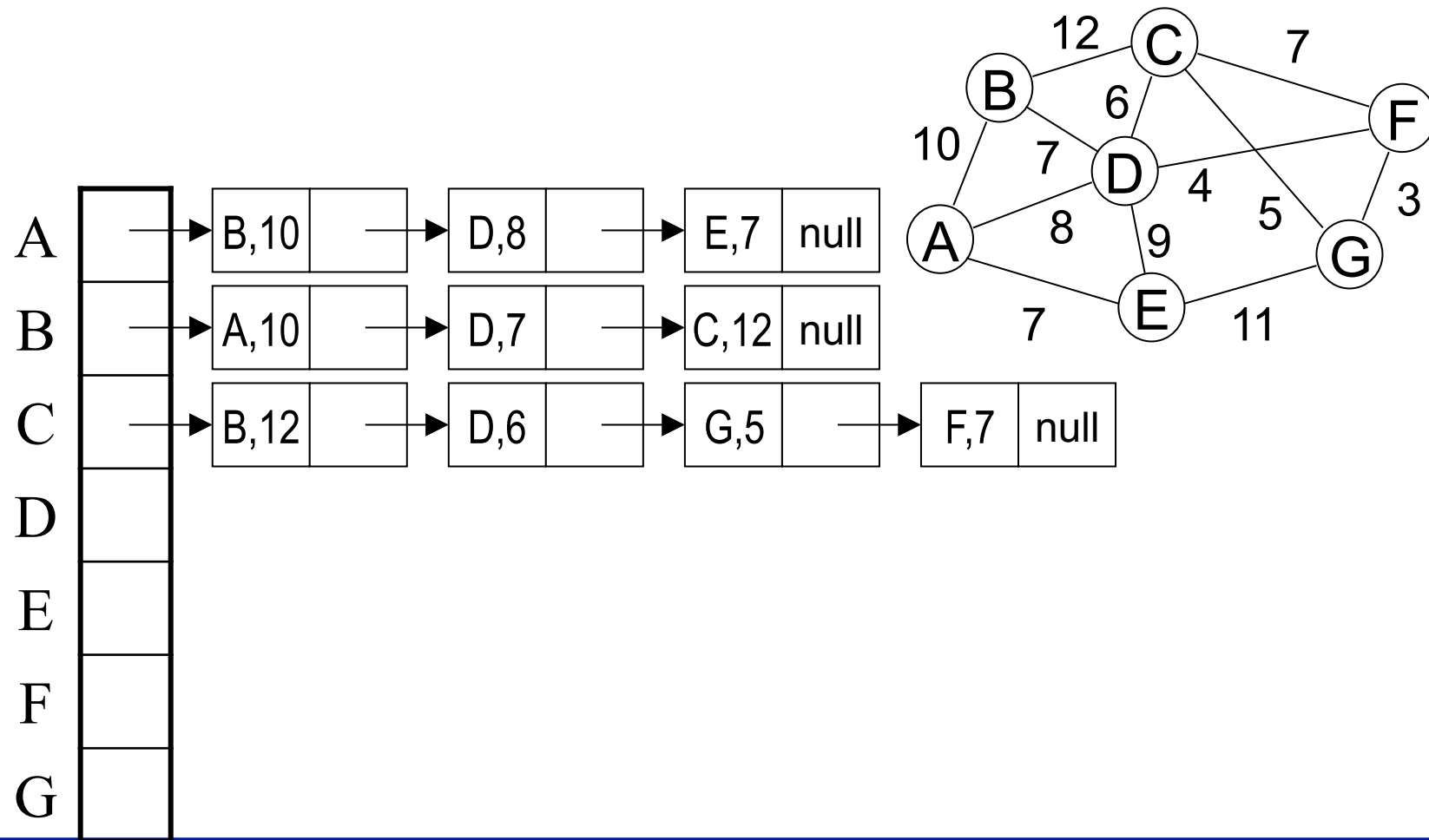
What does this look like?

Graph representation (Adjacency Matrix)

	A	B	C	D	E	F	G
A	0	10	0	8	7	0	0
B	10	0	12	7	0	0	0
C	0	12	0	6	0	7	5
D							
E							
F							
G							



Graph representation (Adjacency List)



Graph representation cost/benefits

- Adjacency matrix
 - Advantage?
 - $O(1)$ access to edge
 - Disadvantage?
 - $O(|V|^2)$ space
- Adjacency list
 - Advantage
 - $O(|V| + |E|)$ space (dominated by $O(|E|)$)
 - good if not too many edges (sparse graph)
 - Disadvantage?
 - $O(|E|)$ access to edge

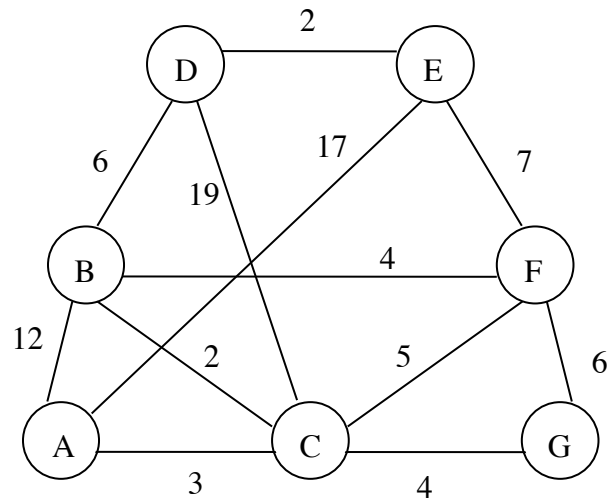
Graph algorithms (a selection)

- Traversals
 - Breadth-first
 - Depth-first
- Others
 - Minimum Spanning Tree (Prim; Kruskal)
 - Tree -> no cycles
 - Spanning -> connects all nodes in the graph
 - Minimum -> lowest overall cost
 - Single-source Shortest Path (Dijkstra)
 - Shortest path from the source to all other nodes in the graph

Graphs (Dijkstra's algorithm)

- Assign every node an initial distance (0 to the source, ∞ for all others); mark all nodes as unvisited
- While there are unvisited nodes:
 - select unvisited node with smallest distance (current)
 - consider all unvisited neighbors of current node:
 - compute distance to each neighbor from current node
 - if less than current distance, replace with new distance
 - mark current node as visited (and is never evaluated again)

Dijkstra example



	A	B	C	D	E	F	G
A->	0	∞	∞	∞	∞	∞	∞
C->		5		22	17	8	7
B->				11	17	8	7
G->				11	17	8	
F->				11	15		
D->					13		
E	0	5	3	11	13	8	7