

Intro to Data Structures

Lecture #21 – BSTs (finale)

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Outline for Today

- HW6 due 11/18
- HW5 grading coming along...
- Some unfinished business – BST wrapup
 - an alternative add method
 - remove(value)
 - an alternate traversal
 - some vocabulary

Implementing a generic BST class

- What methods do we implement on a data structure?
 - ✓ constructor
 - ✓ isEmpty
 - add (today, an alternative implementation)
 - ✓ traversal [inOrder, the rest are variants]
 - ✓ size/count [$O(n)$, for practice]
 - ✓ contains/find [$O(?? \log n)$, but $O(n)$ if tree is not balanced]
 - ✓ toString [a rotated tree]
 - remove [discuss algorithm]

Add (an alternative implementation)

- Add looks a little different from, say, `inOrder()` or `size()` – where's the (simple) base case?
- You don't need to (and probably shouldn't) look ahead in a recursive method
- **But**, you need a way to communicate a change in the parameter across a method call
- **And**, the only way to do that is???

Add (an alternative implementation)

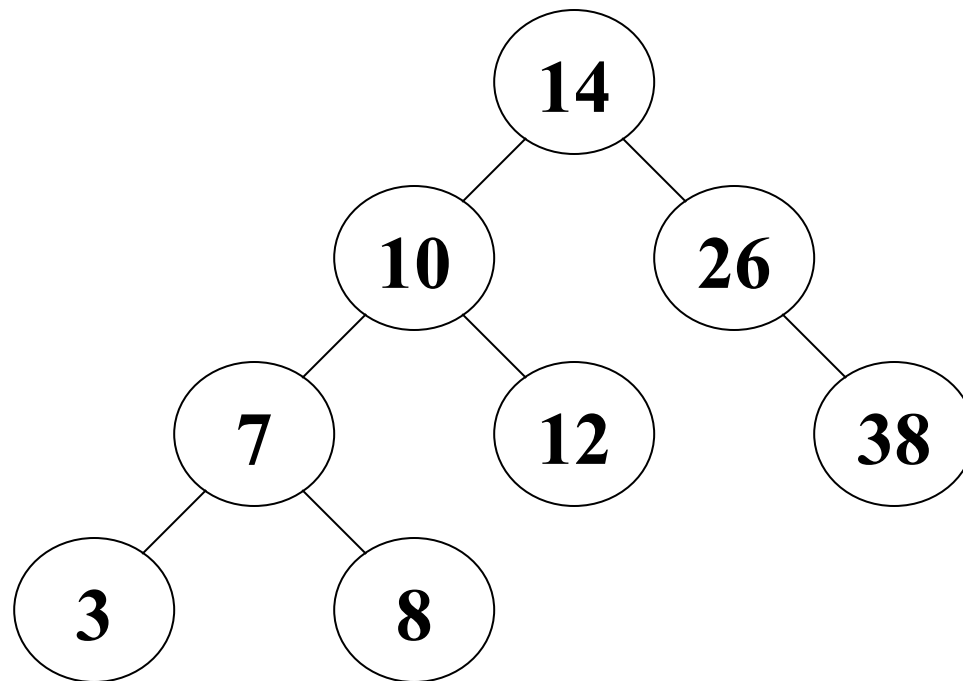
- How would I write a method to double an int?
- What if I wanted to only double odd int's?
- And how would I call that method to effect the change in a particular int variable?
- The programming idiom is
$$x = \text{changed}(x);$$
- Let's write an alternative add without lookahead using this idiom.

Tree traversal (revisited)

- All the traversals we've seen so far (pre-, in-, post-order) are *depth-first*, that is, they start at the root and go as far down a single branch as possible before *backtracking* up that branch and continuing forward as specified.
- Thinking about the array-based implementation of a tree (which is a good quiz question) made me think about a traversal we haven't discussed yet:
 - *breadth-first*

Tree traversal (revisited)

- Given the following BST



Tree traversal (revisited)

- And its array representation

	1	2	3	4	5	6	7	8	9
	14	10	26	7	12		38	3	8

- What would a natural traversal of this representation be?

Tree traversal (breadth-first)

- In a *breadth-first* search, you start at the root and look at all the nodes at the next level, and from each of those, the next level...
- How to do this?
- Queues!
 enqueue the root
 while (queue not empty)
 dequeue node // print (visit)
 enqueue children

Remove

- Three cases to consider for the node to remove.
It can have
 - no children (a leaf – just null that node)
 - one child (loop over it to the child)
 - two children (here's where it gets complicated):
 - could just delete value and reinsert left/right child, but...
 - instead, replace value in node with smallest value in right subtree or largest value in the left subtree, and then delete the node that contained that value (either 1-child or leaf)

Some "final" vocabulary

- *Perfect* binary tree
 - all internal nodes have 2 children
 - all leaf nodes are at same depth
 - a perfect binary tree of height k has $2^{k+1} - 1$ nodes
 - alternatively, you can optimally place n nodes into a binary tree of minimum height ?? $\log n$
- *Complete* binary tree
 - every level, except possibly the last is completely filled
 - last (deepest) level has all leaves as far left as possible