

# Principles of Software Construction

## Concurrency, part 4: In the trenches of parallelism

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# Administrivia

- Homework 5b due tonight
  - Commit by 9 a.m. tomorrow to be considered as a Best Framework
- Do you want to be a software engineer?

# The foundations of the Software Engineering minor

- Core computer science fundamentals
- Building good software
- Organizing a software project
  - Development teams, customers, and users
  - Process, requirements, estimation, management, and methods
- The larger context of software
  - Business, society, policy
- Engineering experience
- Communication skills
  - Written and oral

# SE minor requirements

- Prerequisite: 15-214
- Two core courses
  - 15-313 Foundations of SE (fall semesters)
  - 15-413 SE Practicum (spring semesters)
- Three electives
  - Technical
  - Engineering
  - Business or policy
- Software engineering internship + reflection
  - 8+ weeks in an industrial setting, then
  - 17-413

# To apply to be a Software Engineering minor

- Email [clegoues@cs.cmu.edu](mailto:clegoues@cs.cmu.edu)
  - Your name, Andrew ID, expected grad date, QPA, and minor/majors
  - Why you want to be a SE minor
  - Proposed schedule of coursework
- Spring applications due by Friday, 08 April 2016
  - Only 15 SE minors accepted per graduating class
- More information at:
  - <http://isri.cmu.edu/education/undergrad/>

# Key concepts from Thursday

# java.util.concurrent (j.u.c)

- I. Atomic vars: `java.util.concurrent.atomic`
  - Support various atomic read-modify-write ops
- II. Locks: `java.util.concurrent.locks`
  - Read-write locks, conditions, etc.
- III. Synchronizers
  - Semaphores, cyclic barriers, countdown latches, etc.
- IV. Concurrent collections
  - Shared maps, sets, lists

## V. Data exchange collections summary

- **BlockingQueue**: Supports blocking operations
  - `ArrayBlockingQueue`, `LinkedBlockingQueue`
  - `PriorityBlockingQueue`, `DelayQueue`
  - `SynchronousQueue`
- **BlockingDeque**: Supports blocking operations
  - `LinkedBlockingDeque`
- **TransferQueue**: **BlockingQueue** in which producers may wait for consumers
  - `LinkedTransferQueue`

# Summary of BlockingQueue methods

|                | <i>Throws exception</i> | <i>Special value</i>  | <i>Blocks</i>       | <i>Times out</i>                  |
|----------------|-------------------------|-----------------------|---------------------|-----------------------------------|
| <b>Insert</b>  | <code>add(e)</code>     | <code>offer(e)</code> | <code>put(e)</code> | <code>offer(e, time, unit)</code> |
| <b>Remove</b>  | <code>remove()</code>   | <code>poll()</code>   | <code>take()</code> | <code>poll(time, unit)</code>     |
| <b>Examine</b> | <code>element()</code>  | <code>peek()</code>   | n/a                 | n/a                               |

# Summary of BlockingDeque methods

- First element (head) methods

|                | <i>Throws exception</i>    | <i>Special value</i>       | <i>Blocks</i>            | <i>Times out</i>                       |
|----------------|----------------------------|----------------------------|--------------------------|----------------------------------------|
| <b>Insert</b>  | <code>addFirst(e)</code>   | <code>offerFirst(e)</code> | <code>putFirst(e)</code> | <code>offerFirst(e, time, unit)</code> |
| <b>Remove</b>  | <code>removeFirst()</code> | <code>pollFirst()</code>   | <code>takeFirst()</code> | <code>pollFirst(time, unit)</code>     |
| <b>Examine</b> | <code>getFirst()</code>    | <code>peekFirst()</code>   | n/a                      | n/a                                    |

|                | <i>Throws exception</i>   | <i>Special value</i>      | <i>Blocks</i>           | <i>Times out</i>                      |
|----------------|---------------------------|---------------------------|-------------------------|---------------------------------------|
| <b>Insert</b>  | <code>addLast(e)</code>   | <code>offerLast(e)</code> | <code>putLast(e)</code> | <code>offerLast(e, time, unit)</code> |
| <b>Remove</b>  | <code>removeLast()</code> | <code>pollLast()</code>   | <code>takeLast()</code> | <code>pollLast(time, unit)</code>     |
| <b>Examine</b> | <code>getLast()</code>    | <code>peekLast()</code>   | n/a                     | n/a                                   |

# Today: In the trenches of parallelism

- A high-level view of parallelism
- Concurrent realities
  - ...and the j.u.c Executor framework

# Puzzler: “Racy Little Number”

```
import org.junit.Test;
import static org.junit.Assert.assertEquals;

public class LittleTest {
    int number;

    @Test
    public void test() throws InterruptedException {
        number = 0;
        Thread t = new Thread(() -> {
            assertEquals(2, number);
        });
        number = 1;
        t.start();
        number++;
        t.join();
    }
}
```



# How often does this test pass?

```
import org.junit.Test;
import static org.junit.Assert.assertEquals;

public class LittleTest {
    int number;

    @Test
    public void test() throws InterruptedException {
        number = 0;
        Thread t = new Thread(() -> {
            assertEquals(2, number);
        });
        number = 1;
        t.start();
        number++;
        t.join();
    }
}
```

- (a) It always fails
- (b) It sometimes passes
- (c) It always passes
- (d) It always hangs

# How often does this test pass?

- (a) It always fails
- (b) It sometimes passes
- (c) It always passes – but it tells us nothing
- (d) It always hangs

**JUnit doesn't see assertion failures in other threads**

# Another look

```
import org.junit.*;
import static org.junit.Assert.*;

public class LittleTest {
    int number;

    @Test
    public void test() throws InterruptedException {
        number = 0;
        Thread t = new Thread(() -> {
            assertEquals(2, number); // JUnit never sees the exception!
        });
        number = 1;
        t.start();
        number++;
        t.join();
    }
}
```

# How do you fix it? (1)

```
// Keep track of assertion failures during test
volatile Exception exception;
volatile Error error;

// Triggers test case failure if any thread asserts failed
@After
public void tearDown() throws Exception {
    if (error != null)
        throw error;
    if (exception != null)
        throw exception;
}
```

## How do you fix it? (2)

```
Thread t = new Thread(() -> {  
    try {  
        assertEquals(2, number);  
    } catch(Error e) {  
        error = e;  
    } catch(Exception e) {  
        exception = e;  
    }  
});
```

Now it sometimes passes\*

\*YMMV

# The moral

- JUnit does not support concurrency
- You must provide your own
  - If you don't, you'll get a false sense of security

# Puzzler: “Ping Pong”

```
public class PingPong {  
    public static synchronized void main(String[] a) {  
        Thread t = new Thread(()-> pong() );  
        t.run();  
        System.out.print("Ping");  
    }  
  
    private static synchronized void pong() {  
        System.out.print("Pong");  
    }  
}
```



# What does it print?

```
public class PingPong {  
    public static synchronized void main(String[] a) {  
        Thread t = new Thread(()-> pong() );  
        t.run();  
        System.out.print("Ping");  
    }  
  
    private static synchronized void pong() {  
        System.out.print("Pong");  
    }  
}
```

- (a) PingPong
- (b) PongPing
- (c) It varies

# What does it print?

(a) PingPong

(b) PongPing

(c) It varies

**Not a multithreaded program!**

## Another look

```
public class PingPong {  
    public static synchronized void main(String[] a) {  
        Thread t = new Thread(()-> pong() );  
        t.run(); // An easy typo!  
        System.out.print("Ping");  
    }  
  
    private static synchronized void pong() {  
        System.out.print("Pong");  
    }  
}
```

# How do you fix it?

```
public class PingPong {  
    public static synchronized void main(String[] a) {  
        Thread t = new Thread(()-> pong() );  
        t.start();  
        System.out.print("Ping");  
    }  
  
    private static synchronized void pong() {  
        System.out.print("Pong");  
    }  
}
```

**Now prints PingPong**

# The moral

- Invoke `Thread.start`, not `Thread.run`
  - Can be very difficult to diagnose
- `java.lang.Thread` should not have implemented `Runnable`
  - ...and should not have a public `run` method

# Today: In the trenches of parallelism

- A high-level view of parallelism
- Concurrent realities
  - ...and the j.u.c Executor framework

# Concurrency at the language level

- Consider:

```
Collection<Integer> collection = ...;  
int sum = 0;  
for (int i : collection) {  
    sum += i;  
}
```

- In python:

```
collection = ...  
sum = 0  
for item in collection:  
    sum += item
```

# Parallel quicksort in Nesl

```
function quicksort(a) =  
  if (#a < 2) then a  
  else  
    let pivot    = a[#a/2];  
        lesser   = {e in a | e < pivot};  
        equal    = {e in a | e == pivot};  
        greater  = {e in a | e > pivot};  
        result   = {quicksort(v): v in [lesser,greater]};  
    in result[0] ++ equal ++ result[1];
```

- Operations in `{ }` occur in parallel
- 210esque questions: What is the total work? What is the depth?

# Prefix sums (a.k.a. inclusive scan)

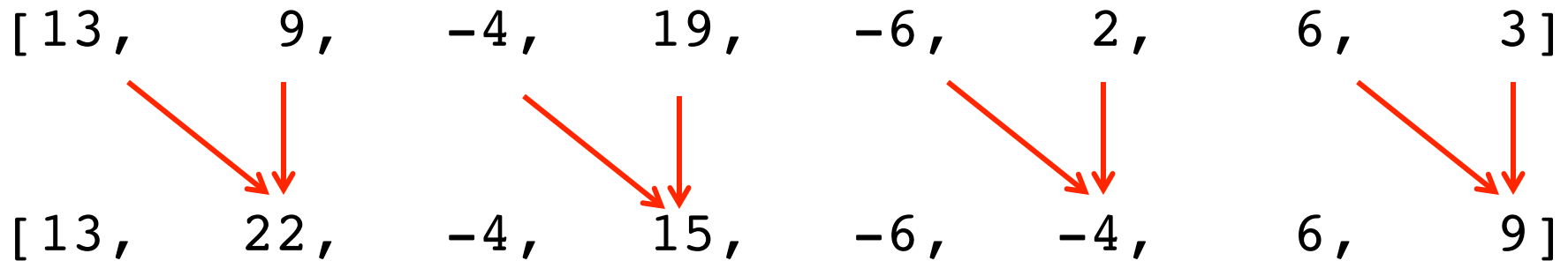
- Goal: given array  $x[0..n-1]$ , compute array of the sum of each prefix of  $x$   
[  $\text{sum}(x[0..0])$ ,  
     $\text{sum}(x[0..1])$ ,  
     $\text{sum}(x[0..2])$ ,  
    ...  
     $\text{sum}(x[0..n-1])$  ]
- e.g.,  $x = [13, 9, -4, 19, -6, 2, 6, 3]$   
prefix sums:  $[13, 22, 18, 37, 31, 33, 39, 42]$

# Parallel prefix sums

- Intuition: If we have already computed the partial sums  $\text{sum}(x[0\dots3])$  and  $\text{sum}(x[4\dots7])$ , then we can easily compute  $\text{sum}(x[0\dots7])$
- e.g.,  $x = [13, 9, -4, 19, -6, 2, 6, 3]$

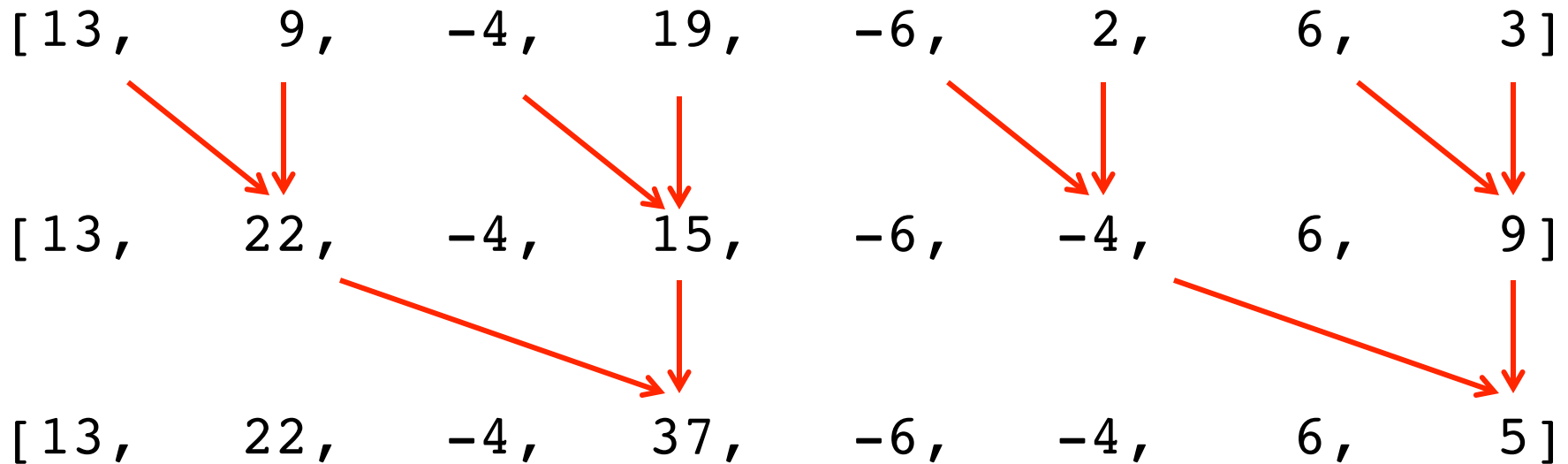
# Parallel prefix sums algorithm, upsweep

- Computes the partial sums in a more useful manner



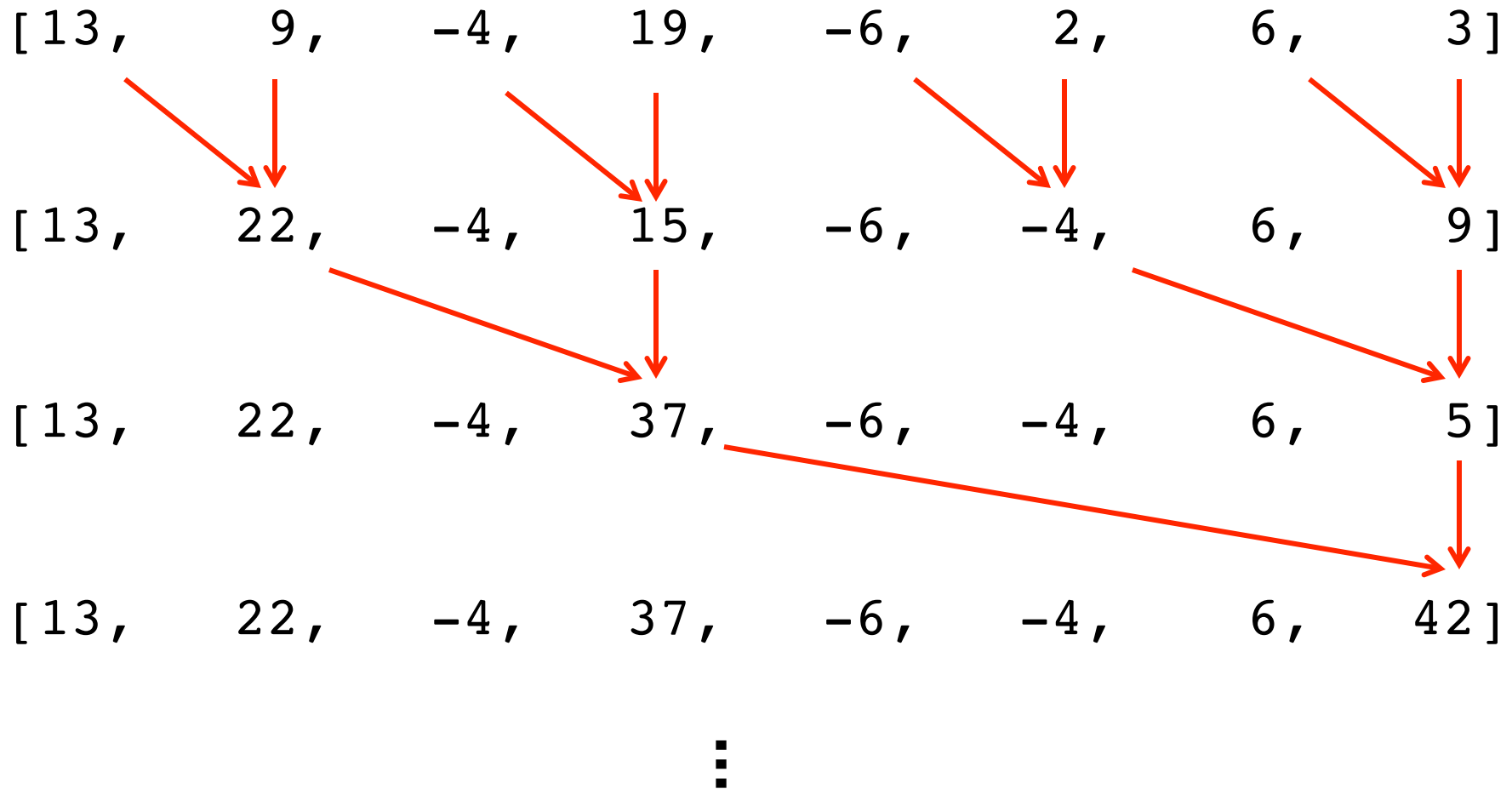
# Parallel prefix sums algorithm, upsweep

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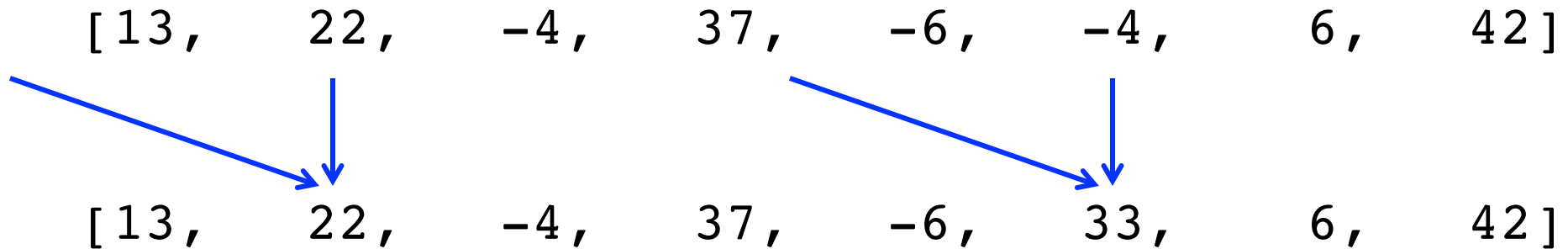
# Parallel prefix sums algorithm, upsweep

- Computes the partial sums in a more useful manner



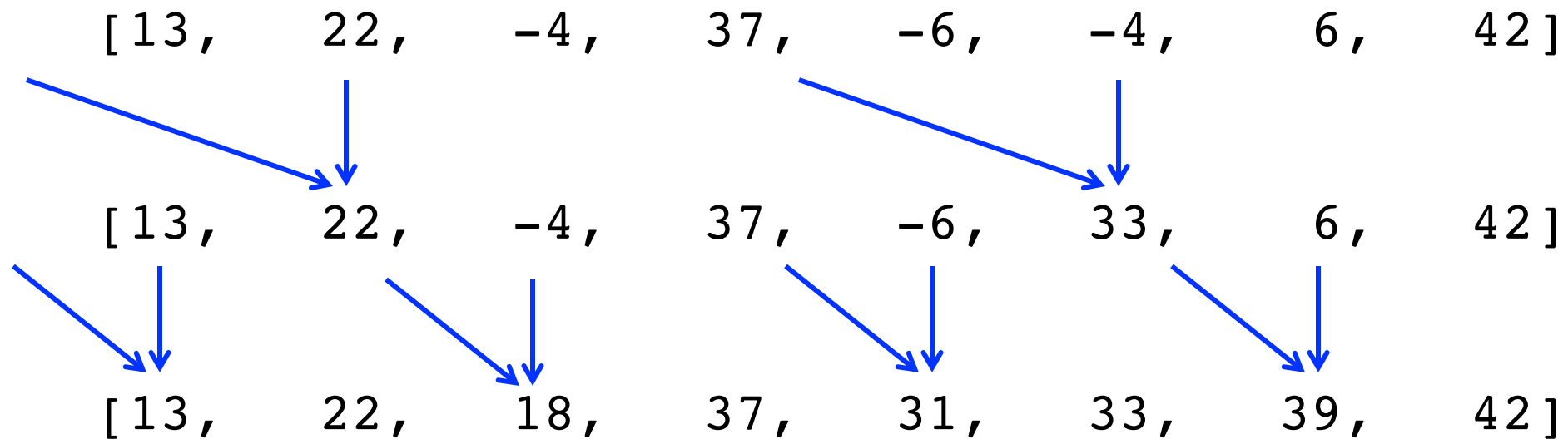
# Parallel prefix sums algorithm, downsweep

- Now unwinds to calculate the other sums



# Parallel prefix sums algorithm, downsweep

- Now unwinds to calculate the other sums



- Recall, we started with:

$[13, 9, -4, 19, -6, 2, 6, 3]$

# Parallel prefix sums

- Intuition: If we have already computed the partial sums  $\text{sum}(x[0..3])$  and  $\text{sum}(x[4..7])$ , then we can easily compute  $\text{sum}(x[0..7])$
- e.g.,  $x = [13, 9, -4, 19, -6, 2, 6, 3]$

- Pseudocode:

```
prefix_sums(x):
```

```
    for d in 0 to (lg n)-1:           // d is depth
```

```
        parallelfor i in  $2^d-1$  to n-1, by  $2^{d+1}$ :
```

```
             $x[i+2^d] = x[i] + x[i+2^d]$ 
```

```
    for d in (lg n)-1 to 0:
```

```
        parallelfor i in  $2^d-1$  to n-1- $2^d$ , by  $2^{d+1}$ :
```

```
            if  $(i-2^d \geq 0)$ :
```

```
                 $x[i] = x[i] + x[i-2^d]$ 
```

# Parallel prefix sums algorithm, in code

- An iterative Java-esque implementation:

```
void iterativePrefixSums(long[] a) {  
    int gap = 1;  
    for ( ; gap < a.length; gap *= 2) {  
        parfor(int i=gap-1; i+gap<a.length; i += 2*gap) {  
            a[i+gap] = a[i] + a[i+gap];  
        }  
    }  
    for ( ; gap > 0; gap /= 2) {  
        parfor(int i=gap-1; i<a.length; i += 2*gap) {  
            a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);  
        }  
    }  
}
```

# Parallel prefix sums algorithm, in code

- A recursive Java-esque implementation:

```
void recursivePrefixSums(long[] a, int gap) {
    if (2*gap - 1 >= a.length) {
        return;
    }

    parfor(int i=gap-1; i+gap<a.length; i += 2*gap) {
        a[i+gap] = a[i] + a[i+gap];
    }

    recursivePrefixSums(a, gap*2);

    parfor(int i=gap-1; i<a.length; i += 2*gap) {
        a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);
    }
}
```

# Parallel prefix sums algorithm

- How good is this?

# Parallel prefix sums algorithm

- How good is this?
  - Work:  $O(n)$
  - Depth:  $O(\lg n)$
- See `PrefixSums.java`,  
`PrefixSumsSequentialWithParallelWork.java`

# Goal: parallelize the PrefixSums implementation

- Specifically, parallelize the parallelizable loops

```
parfor(int i=gap-1; i+gap<a.length; i += 2*gap) {  
    a[i+gap] = a[i] + a[i+gap];  
}
```
- Partition into multiple segments, run in different threads

```
for(int i=left+gap-1; i+gap<right; i += 2*gap) {  
    a[i+gap] = a[i] + a[i+gap];  
}
```

# Recall the Java primitive concurrency tools

- The `java.lang.Runnable` interface

```
void          run();
```

- The `java.lang.Thread` class

```
Thread(Runnable r);
```

```
void          start();
```

```
static void    sleep(long millis);
```

```
void          join();
```

```
boolean        isAlive();
```

```
static Thread  currentThread();
```

# Recall the Java primitive concurrency tools

- The `java.lang.Runnable` interface  
`void run();`
- The `java.lang.Thread` class  
`Thread(Runnable r);`  
`void start();`  
`static void sleep(long millis);`  
`void join();`  
`boolean isAlive();`  
`static Thread currentThread();`
- The `java.util.concurrent.Callable<V>` interface
  - Like `java.lang.Runnable` but can return a value  
`V call();`

# A framework for asynchronous computation

- The `java.util.concurrent.Future<V>` interface

```
V          get();
V          get(long timeout, TimeUnit unit);
boolean    isDone();
boolean    cancel(boolean mayInterruptIfRunning);
boolean    isCancelled();
```

# A framework for asynchronous computation

- The `java.util.concurrent.Future<V>` interface:

```
V          get();  
V          get(long timeout, TimeUnit unit);  
boolean    isDone();  
boolean    cancel(boolean mayInterruptIfRunning);  
boolean    isCancelled();
```

- The `java.util.concurrent.ExecutorService` interface:

```
Future<?>  submit(Runnable task);  
Future<V>  submit(Callable<V> task);  
List<Future<V>>  
    invokeAll(Collection<? extends Callable<V>> tasks);  
Future<V>  
    invokeAny(Collection<? extends Callable<V>> tasks);  
void shutdown();
```

# Executors for common computational patterns

- From the `java.util.concurrent.Executors` class  
`static ExecutorService newSingleThreadExecutor();`  
`static ExecutorService newFixedThreadPool(int n);`  
`static ExecutorService newCachedThreadPool();`  
`static ExecutorService newScheduledThreadPool(int n);`
- Aside: see `NetworkServer.java` (later)

# Fork/Join: another common computational pattern

- In a long computation:
  - Fork a thread (or more) to do some work
  - Join the thread(s) to obtain the result of the work

# Fork/Join: another common computational pattern

- In a long computation:
  - Fork a thread (or more) to do some work
  - Join the thread(s) to obtain the result of the work
- The `java.util.concurrent.ForkJoinPool` class
  - Implements `ExecutorService`
  - Executes `java.util.concurrent.ForkJoinTask<V>` or `java.util.concurrent.RecursiveTask<V>` or `java.util.concurrent.RecursiveAction`

# The RecursiveAction abstract class

```
public class MyActionFoo extends RecursiveAction {  
    public MyActionFoo(...) {  
        store the data fields we need  
    }  
  
    @Override  
    public void compute() {  
        if (the task is small) {  
            do the work here;  
            return;  
        }  
  
        invokeAll(new MyActionFoo(...), // smaller  
                  new MyActionFoo(...), // tasks  
                  ...);                  // ...  
    }  
}
```

# A ForkJoin example

- See `PrefixSumsParallelForkJoin.java`
- See the processor go, go go!

# Parallel prefix sums algorithm

- How good is this?
  - Work:  $O(n)$
  - Depth:  $O(\lg n)$
- See `PrefixSumsParallelArrays.java`

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- See `PrefixSumsSequential.java`

# Parallel prefix sums algorithm

- How good is this?
  - Work:  $O(n)$
  - Depth:  $O(\lg n)$
- See `PrefixSumsParallelArrays.java`
- See `PrefixSumsSequential.java`
  - $n-1$  additions
  - Memory access is sequential
- For `PrefixSumsSequentialWithParallelWork.java`
  - About  $2n$  useful additions, plus extra additions for the loop indexes
  - Memory access is non-sequential
- The punchline:
  - Don't roll your own
  - Cache and constants matter

# Coming Thursday...

- Distributed systems

## In-class example for parallel prefix sums

[ 7,      5,      8,   -36,    17,      2,      21,      18 ]