

# UNIT 5A Recursion: Basics

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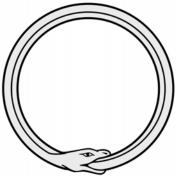
### Feedback in Autolab

- When your CAs grade your programming assignments, they should be leaving feedback if you don't get full credit on a problem.
- Click on the score to see the feedback.
- If you still have questions about why an answer wasn't correct, ask your CA.

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### Recursion

- A "recursive" function is one that calls itself.
- Infinite loop? Not necessarily.



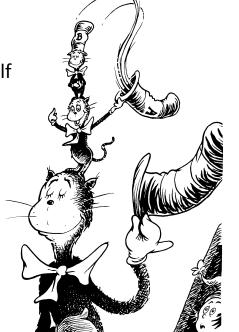
Not like this.

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 The recursive function calls itself on a smaller version of the problem to be solved.

• Recursion looks more like this:



### **Recursive Definitions**

- Every recursive definition includes two parts:
  - Base case (non-recursive)
     A simple case that can be done without solving the same problem again.
  - Recursive case(s)
     One or more cases that are "simpler" versions of the original problem.
    - By "simpler", we sometimes mean "smaller" or "shorter" or "closer to the base case".

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# Example: Factorial

- $N! = N \times (N-1) \times (N-2) \times \cdots \times 1$
- $5! = 5 \times 4 \times 3 \times 2 \times 1$
- $6! = 6 \times 5 \times 4 \times 3 \times 2 \times 1$
- So  $6! = 6 \times 5!$
- And  $5! = 5 \times 4!$
- And  $4! = 4 \times 3!$
- What is the base case? 0! = 1

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# Factorial in Ruby (Recursive)

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# **Tracing Factorial**

```
factorial(5) = 5 * factorial(4)
  factorial(4) = 4 * factorial(3)
    factorial(3) = 3 * factorial(2)
    factorial(2) = 2 * factorial(1)
    factorial(1) = 1 * factorial(0)
    factorial(0) = 1
```

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# **Tracing Factorial**

```
factorial(5) = 5 * factorial(4)
  factorial(4) = 4 * factorial(3)
   factorial(3) = 3 * factorial(2)
    factorial(2) = 2 * factorial(1)
     factorial(1) = 1 * factorial(0) = 1 * 1 = 1
     factorial(0) = 1
```

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## **Tracing Factorial**

```
factorial(5) = 5 * factorial(4)

factorial(4) = 4 * factorial(3)

factorial(3) = 3 * factorial(2)

factorial(2) = 2 * factorial(1) = 2 * 1 = 2

factorial(1) = 1 * factorial(0) = 1 * 1 = 1

factorial(0) = 1
```

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# **Tracing Factorial**

```
factorial(5) = 5 * factorial(4)

factorial(4) = 4 * factorial(3)

factorial(3) = 3 * factorial(2) = 3 * 2 = 6

factorial(2) = 2 * factorial(1) = 2 * 1 = 2

factorial(1) = 1 * factorial(0) = 1 * 1 = 1

factorial(0) = 1
```

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## **Tracing Factorial**

```
factorial(5) = 5 * factorial(4)

factorial(4) = 4 * factorial(3) = 4 * 6 = 24

factorial(3) = 3 * factorial(2) = 3 * 2 = 6

factorial(2) = 2 * factorial(1) = 2 * 1 = 2

factorial(1) = 1 * factorial(0) = 1 * 1 = 1

factorial(0) = 1
```

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## **Tracing Factorial**

```
factorial(5) = 5 * factorial(4) = 5 * 24 = 120

factorial(4) = 4 * factorial(3) = 4 * 6 = 24

factorial(3) = 3 * factorial(2) = 3 * 2 = 6

factorial(2) = 2 * factorial(1) = 2 * 1 = 2

factorial(1) = 1 * factorial(0) = 1 * 1 = 1

factorial(0) = 1
```

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## Recursive vs. Iterative Solutions

- For every recursive function, there is an equivalent iterative solution.
- For every iterative function, there is an equivalent recursive solution.
- But some problems are easier to solve one way than the other way.

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# Factorial Function (Iterative)

```
def factorial (n)
  result = 1
  for i in 1..n do
    result = result * i
  end
  return result
end
```

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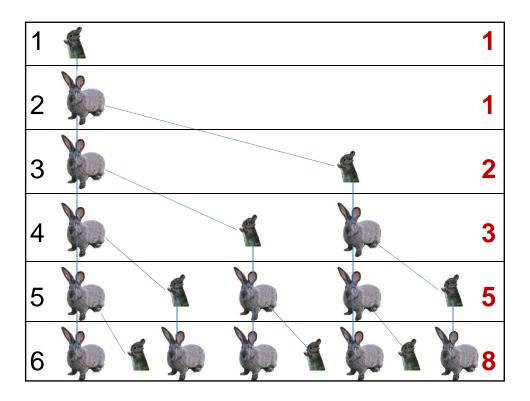
## Fibonacci Sequence

- Start with 1 pair of baby rabbits.
- Babies take 1 month to reach maturity.
- Mature rabbits produce 1 new pair of babies every month.
- After a year, how many rabbits do you have?





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## Recursive Fibonacci

- Base case: we start with nothing.
  - fib(0) is 0
- In the first month we have 1 baby rabbit:
  - fib(1) is 1
- At n>1 months, the number of rabbits is:
   # of rabbits from last month fib(n-1)

+

# of babies born this month ????

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### Recursive Fibonacci

- How many babies are born in month n?
  - One baby for every adult who was alive at n-1
- How many adults were alive in month n-1?
  - As many as the total number of rabbits at n-2
- Therefore:

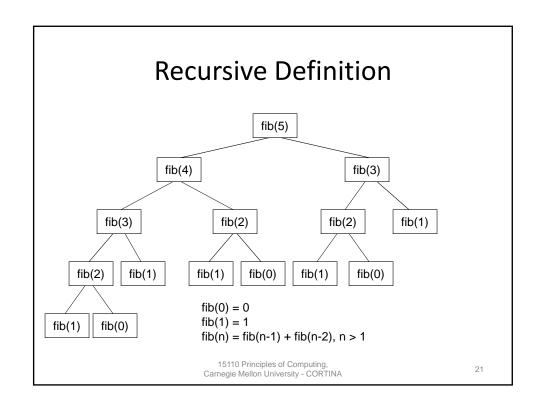
```
fib(n) = fib(n-1) + fib(n-2)

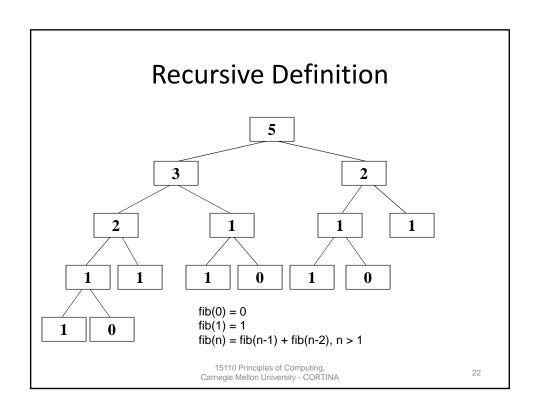
Adults in Babies in month n month n month n

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```

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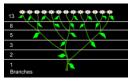
## Recursive Fibonacci in Ruby





### Fibonacci Numbers in Nature

- 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, etc.
- Number of branches on a tree.
- Number of petals on a flower.
- Number of spirals on a pineapple.



















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#### Iterative Fibonacci

```
def fib(n)
  x = 0
  next_x = 1
  for i in 1..n do
    x, next_x = next_x, x+next_x
  end
  return x
end

Much faster than
the recursive
version. Why?
```

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```
def gcd2(x, y)
  if y == 0 then
    return x
  else
    return gcd2(y, x % y)
  end
end

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    def gcd2(x, y)
    base case
    recursive
    case
    (a "simpler"
    version of
    the same
    problem)
```

#### Recursive sum of a list def sumlist(list) n = list.length Base case: if n == 0 then The sum of an empty list is 0. return 0 else return list[0] + sumlist(list[1..n-1]) end end Recursive case: The sum of a list is the first element + the sum of the rest of the list. 15110 Principles of Computing, Carnegie Mellon University - CORTINA 26

## **Towers of Hanoi**

 A puzzle invented by French mathematician Edouard Lucas in 1883.



Towers of Hanoi with 8 discs.

- At a temple far away, priests were led to a courtyard with three pegs and 64 discs stacked on one peg in size order.
  - Priests are only allowed to move one disc at a time from one peg to another.
  - Priests may not put a larger disc on top of a smaller disc at any time.
- The goal of the priests was to move all 64 discs from the leftmost peg to the rightmost peg.
- According to the story, the world would end when the priests finished their work.

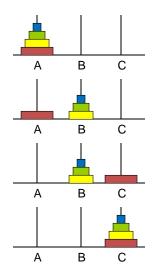
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#### Towers of Hanoi

Problem: Move n discs from peg A to peg C using peg B.

- 1. Move n-1 discs from peg A to peg B using peg C. (recursive step)
- 2. Move 1 disc from peg A to peg C.
- 3. Move n-1 discs from peg B to C using peg A. (recursive step)



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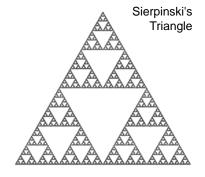
## Towers of Hanoi in Ruby

```
def hanoi (disks, from, temp, to)
  n = disks.length
  if n > 1 then
    towers(disks[1..n-1], from, to, temp)
  end
  print "Move ",disks[0]," from ", from,
        " to ", to, "\n"
  if n > 1 then
    hanoi(disks[1..n-1], temp, from, to)
  end
end
In irb: towers([4,3,2,1], "A", "B", "C")
How many moves do the priests need to move 64 discs?
```

# Geometric Recursion (Fractals)

• A recursive operation performed on successively smaller regions.





http://fusionanomaly.net/recursion.jpg

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