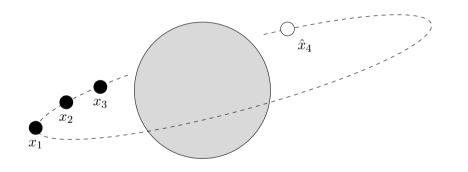
## Differentiable Optimisation in Deep Learning

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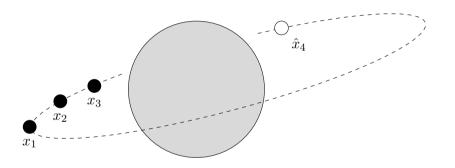
15 December, 2022

# Discovery of Ceres





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- ▶ **logistics and planning:** find the cheapest way to distribute goods from suppliers to consumers across a network
- **statistics/data science:** curve fitting and data visualisation
- ► machine learning and deep learning: minimise loss functions with respect to the parameters of our model

#### Overview

- Introduction to Optimisation
  - Formal definition
  - Least squares
  - Convex sets and functions
  - Convex optimisation problems
  - Lagrangian
  - Optimality conditions
  - Algorithms
- Differentiable Optimisation and Deep Learning
  - ► Machine learning from 10,000ft

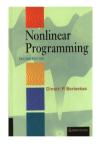
- Automatic differentiation
- Forward and backward passes
- Imperative and declarative nodes
- Bi-level optimisation
- Implicit function theorem
- Differentiable optimisation results
- Examples and Applications
  - Least squares
  - Optimal transport
  - ► Blind perspective-n-point

accompanying lecture notes available at https://users.cecs.anu.edu.au/~sgould

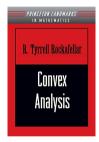
#### lecture 1

### Lecture 1: Introduction to Optimisation



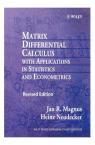


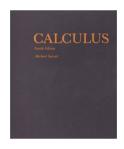




## Assumed Background







## Optimisation Problems

find the assignment to variables that minimises a measure of cost subject to some constraints<sup>1</sup>

<sup>&</sup>lt;sup>1</sup>In these lectures we will be concerned with continuous-valued variables

## Optimisation Problems

 $\begin{array}{ll} \text{minimize (over } x) & \text{objective}(x) \\ \text{subject to} & \text{constraints}(x) \end{array}$ 

## **Optimisation Problems**

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\ldots,p \\ & h_i(x)=0, \quad i=1,\ldots,q \end{array}$$

- $x = (x_1, \dots, x_n) \in \mathbb{R}^n$  optimisation variables
- $f_0: \mathbb{R}^n \to \mathbb{R}$  objective (or cost or loss) function
- $f_i:\mathbb{R}^n \to \mathbb{R}, \ i=1,\ldots,p$  inequality constraint functions
- $ightharpoonup h_i: \mathbb{R}^n o \mathbb{R}, \ i=1,\ldots,q$  equality constraint functions

## Solution and Optimal Value

A point x is **feasible** if  $x \in \text{dom}(f_0)$  and it satisfies the constraints.

A **solution**, or optimal point,  $x^*$  has the smallest value of  $f_0$  among all feasible x.

<sup>&</sup>lt;sup>1</sup>Warning: notation clash between p and  $p^*$ !

## Solution and Optimal Value

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A **solution**, or optimal point,  $x^*$  has the smallest value of  $f_0$  among all feasible x.

The **optimal value** is<sup>1</sup>

$$p^* = \inf_{x \in \mathcal{D}} \left\{ f_0(x) \mid f_i(x) \le 0, \quad i = 1, \dots, p \\ h_i(x) = 0, \quad i = 1, \dots, q \right\}.$$

- $\triangleright p^*$  and is equal to  $f_0(x^*)$  when  $x^*$  exists
- $ightharpoonup p^* = \infty$  if the problem is infeasible (no x satisfies the constraints)
- $p^* = -\infty$  if the problem is unbounded below

<sup>&</sup>lt;sup>1</sup>Warning: notation clash between p and  $p^*$ !

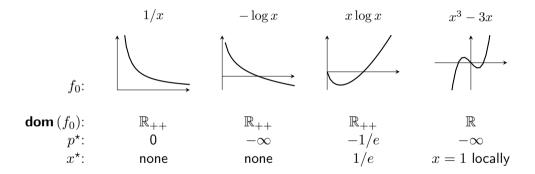
## Locally Optimal Points

A point x is **locally optimal** if there is an R>0 such that z=x is optimal for

```
minimize (over z) f_0(z) subject to f_i(z) \leq 0 \qquad \qquad i=1,\ldots,p h_i(z) = 0 \qquad \qquad i=1,\ldots,q \|z-x\|_2 \leq R.
```

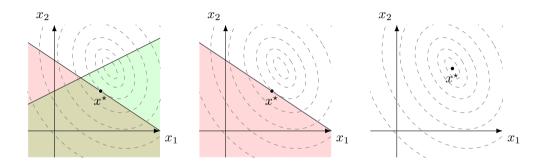
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## Examples (1D)



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# Examples (2D)



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## Least Squares

 $\text{minimize} \quad \|Ax-b\|_2^2$ 

### Least Squares

minimize 
$$||Ax - b||_2^2$$

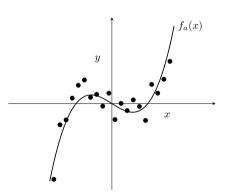
- unique solution if  $A^TA$  is invertible,  $x^* = (A^TA)^{-1}A^Tb$
- $\blacktriangleright$  solution via SVD,  $A=U\Sigma V^T$  , if  $A^T\!A$  not invertible,  $x^\star=V\Sigma^{-1}U^Tb$ 
  - ▶ in fact,  $x^* + w$  for any  $w \in \mathcal{N}(A)$  also a solution
- ightharpoonup solution via QR factorisation,  $x^{\star} = R^{-1}Q^Tb$
- ightharpoonup solved in  $O(n^2m)$  time, less if structured
- typically use iterative solver

## Example: Polynomial Curve Fitting

fit n-th order polynomial  $f_a(x) = \sum_{k=0}^n a_k x^k$  to set of noisy points  $\{(x_i, y_i)\}_{i=1}^m$ 

minimize (over 
$$a$$
)  $\sum_{i=1}^{m} (f_a(x_i) - y_i)^2$ 

special case of convex optimisation

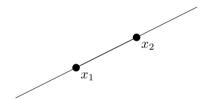


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## Lines and Line Segments

lacktriangle a **line** through two points  $x_1$  and  $x_2$ 

$$x = \theta x_1 + (1 - \theta)x_2, \quad (\theta \in \mathbb{R})$$



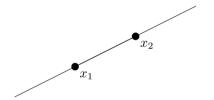
- ➤ an affine set contains the line through any two distinct points in the set
- ▶ an affine hull the set formed by taking all lines through points in a set

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### Lines and Line Segments

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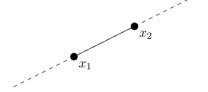
$$x = \theta x_1 + (1 - \theta)x_2, \quad (\theta \in \mathbb{R})$$



- ➤ an affine set contains the line through any two distinct points in the set
- ➤ an **affine hull** the set formed by taking all lines through points in a set

▶ a **line segment** between  $x_1$  and  $x_2$ 

$$x = \theta x_1 + (1 - \theta)x_2, \quad (0 \le \theta \le 1)$$

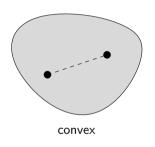


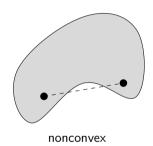
- a convex set contains the line segment between any two distinct points in the set
- ► an **convex hull** the set formed by taking all line segments between points in a set

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#### Convex Sets

$$x_1, x_2 \in \text{convex set } C \implies \theta x_1 + (1-\theta)x_2 \in C \text{ for all } 0 \leq \theta \leq 1$$



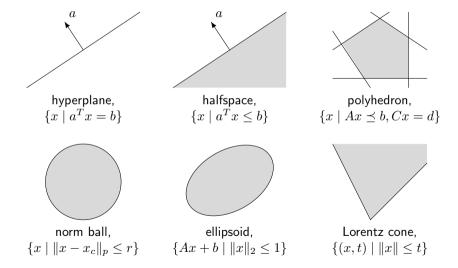


common examples in machine learning:

- ▶ nonnegative orthant,  $\mathbb{R}^n_+ = \{x \mid x_i \geq 0, i = 1, \dots, n\}$
- **p** positive semindefinite matrices,  $\mathbb{S}^n_+ = \{X \mid z^T X z \geq 0, z \in \mathbb{R}^n\}$

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## More Examples



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#### Convex Functions

A function  $f: \mathbb{R}^n \to \mathbb{R}$  is convex if  $\mathbf{dom}(f)$  is a convex set and

$$f(\theta x + (1 - \theta)y) \le \theta f(x) + (1 - \theta)f(y)$$

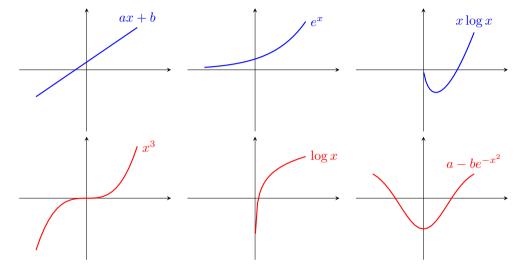
for all  $x, y \in \mathbf{dom}(f), 0 \le \theta \le 1$ .



ightharpoonup f is convex if -f is convex

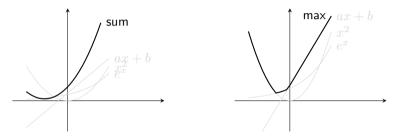
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# Examples



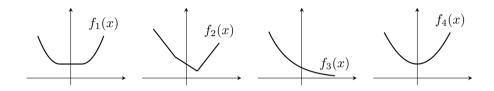
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# Weighted Sum and Pointwise Maximum Preserve Convexity



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## Convex, Strictly Convex, and Strongly Convex



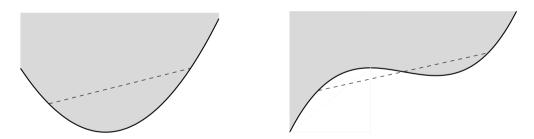
- ▶  $f_1$  is smooth and convex:  $f(\theta x + (1 \theta)y) \le \theta f(x) + (1 \theta)y$
- ▶  $f_2$  is non-differentiable and convex:  $f(\theta x + (1 \theta)y) \le \theta f(x) + (1 \theta)y$
- ▶  $f_3$  is strictly convex:  $f(\theta x + (1 \theta)y) < \theta f(x) + (1 \theta)y$
- ▶  $f_4$  is strongly convex:  $\exists m \text{ s.t. } m(y-x)^2 \leq f(y) f(x)$

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### **Epigraph**

The epigraph of function  $f: \mathbb{R}^n \to \mathbb{R}$  is the set

$$\operatorname{epi}(f) = \{(x,t) \in \mathbb{R}^{n+1} \mid x \in \operatorname{dom}\left(f\right), f(x) \leq t\}.$$



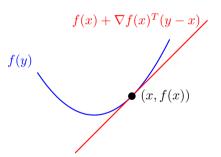
ightharpoonup f is a convex function if and only if epi(f) is a convex set

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#### First-order Condition

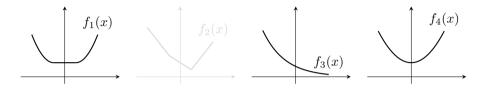
differentiable f with convex domain is convex iff

$$f(y) \geq f(x) + \nabla f(x)^T (y-x) \quad \text{for all } x,y \in \operatorname{dom}\left(f\right)$$



 $\triangleright$  first-order approximation of (convex) f is a global under estimator

#### Second-order Condition



twice differentiable f with convex domain is convex iff

$$\nabla^2 f(x) \succeq 0$$
 for all  $x \in \mathbf{dom}(f)$ 

- ▶ if  $\nabla^2 f(x) \succ 0$  for all  $x \in \mathbf{dom}(f)$ , then f is strictly convex
- ▶ if  $\nabla^2 f(x) \succeq mI$  for some m > 0 and all  $x \in \mathbf{dom}(f)$ , then f is strongly convex

strongly convex functions have a unique minimum

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# Worked Example: log-sum-exp is Convex

$$f(x) = \log \sum_{k=1}^{n} \exp x_k$$

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Proof. Start by computing the gradient and Hessian,

$$\frac{\partial f(x)}{\partial x_i} = \frac{\exp x_i}{\sum_{k=1}^n \exp x_k} \qquad \qquad \text{(derivative of } \log(z), \ z'/z)$$
 
$$\frac{\partial^2 f(x)}{\partial x_i \partial x_j} = \frac{\left(\sum_{k=1}^n \exp x_k\right) \left[\!\!\left[i=j\right]\!\!\right] \exp x_i - \exp x_i \exp x_j}{\left(\sum_{k=1}^n \exp x_k\right)^2} \qquad \qquad \text{(quotient rule, } \frac{v \cdot \mathrm{d} u - u \cdot \mathrm{d} v}{v^2})$$

$$f(x) = \log \sum_{k=1}^{n} \exp x_k$$

Proof. Start by computing the gradient and Hessian,

$$\frac{\partial f(x)}{\partial x_i} = \frac{z_i}{\mathbf{1}^T z} \qquad (z_k = \exp x_k)$$

$$\frac{\partial^2 f(x)}{\partial x_i \partial x_j} = \frac{(\mathbf{1}^T z) \left[ i = j \right] z_i - z_i z_j}{(\mathbf{1}^T z)^2}$$

$$f(x) = \log \sum_{k=1}^{n} \exp x_k$$

Proof. Start by computing the gradient and Hessian,

$$\nabla f(x) = \frac{1}{\mathbf{1}^T z} z \qquad (z_k = \exp x_k)$$

$$\nabla^2 f(x) = \frac{1}{(\mathbf{1}^T z)^2} \left( (\mathbf{1}^T z) \operatorname{diag}(z) - z z^T \right)$$

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$$\nabla^2 f(x) = \frac{1}{\left(\mathbf{1}^T z\right)^2} \left( \left(\mathbf{1}^T z\right) \operatorname{diag}(z) - z z^T \right) \qquad (z_k = \exp x_k)$$

To show that  $\nabla^2 f(x) \succeq 0$ , we must verify that  $v^T \nabla^2 f(x) v \geq 0$  for all v.

$$f(x) = \log \sum_{k=1}^{n} \exp x_k$$

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$$\begin{split} \boldsymbol{v}^T \nabla^2 f(\boldsymbol{x}) \boldsymbol{v} &= \frac{1}{\left(\mathbf{1}^T \boldsymbol{z}\right)^2} \, \boldsymbol{v}^T \Big( (\mathbf{1}^T \boldsymbol{z}) \mathrm{diag}(\boldsymbol{z}) - \boldsymbol{z} \boldsymbol{z}^T \Big) \, \boldsymbol{v} \\ &= \frac{1}{\left(\mathbf{1}^T \boldsymbol{z}\right)^2} \, \Big( (\mathbf{1}^T \boldsymbol{z}) \boldsymbol{v}^T \mathrm{diag}(\boldsymbol{z}) \boldsymbol{v} - \boldsymbol{v}^T \boldsymbol{z} \boldsymbol{z}^T \boldsymbol{v} \Big) \end{split}$$

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$$\begin{split} v^T \nabla^2 f(x) v &= \frac{1}{\left(\mathbf{1}^T z\right)^2} \, v^T \Big( (\mathbf{1}^T z) \mathrm{diag}(z) - z z^T \Big) \, v \\ &= \frac{1}{\left(\mathbf{1}^T z\right)^2} \, \Big( (\mathbf{1}^T z) v^T \mathrm{diag}(z) v - v^T z z^T v \Big) \end{split}$$

Therefore we need to show that  $(\mathbf{1}^T z)v^T \operatorname{diag}(z)v \geq (v^T z)^2$  for all v.

$$f(x) = \log \sum_{k=1}^{n} \exp x_k$$

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$$\left(\sum_{k=1}^{n} z_k\right) \left(\sum_{k=1}^{n} z_k v_k^2\right) \ge \left(\sum_{k=1}^{n} v_k z_k\right)^2$$

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which is true by the Cauchy-Schwarz inequality,  $\|a\|_2^2 \|b\|_2^2 \ge (a^T b)^2$ , with  $a = (\sqrt{z_1}, \dots, \sqrt{z_n})$  and  $b = (\sqrt{z_1}v_1, \dots, \sqrt{z_n}v_n)$ .

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## Convex Optimisation

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, p \\ & a_i^T x = b_i, \quad i = 1, \dots, q \end{array}$$

- $ightharpoonup f_0, f_1, \ldots, f_p$  are convex
- $ightharpoonup h_i(x) riangleq a_i^T x b_i$  are affine, often written as Ax = b

minimise a convex objective over a convex feasible set

#### Local Optima are Global Optima

any local minimum of a convex problem is (globally) optimal

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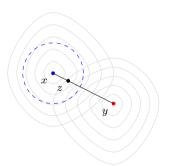
#### Proof Sketch.

- towards contradiction, suppose x is locally optimal, but there exists a feasible y with lower objective
- since x is locally optimally there exists a radius R such that no other point within R of x has lower objective
- (so y must be further than R from x)
- lacktriangle pick a point z on the line segment between x and y and within R of x
- lacktriangle so z must be feasible and have objective no lower than x
- but, by the basic inequality of convex functions,

$$f_0(\theta x + (1 - \theta)y) \le \theta f_0(x) + (1 - \theta)f_0(y),$$

the objective value at z must be between that at x and y, i.e., lower than  $f_0(x)$ 

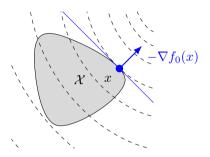
we have a contradiction





# Optimality Criterion for Differentiable $f_0$

x is optimal if and only if it is feasible and  $\nabla f_0(x)^T(y-x) \geq 0$  for all feasible y



if nonzero.

- $\triangleright \nabla f_0(x)$  defines a supporting hyperplane to feasible set  $\mathcal{X}$  at x
- $ightharpoonup f_0$  cannot be improved by moving in a direction where x stays feasible

#### Lagrangian

Standard form problem (not necessarily convex),

```
\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\ldots,p \\ & h_i(x)=0, \quad i=1,\ldots,q \end{array}
```

variable  $x \in \mathbb{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^*$ 

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Standard form problem (not necessarily convex),

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\dots,p \\ & h_i(x) = 0, \quad i=1,\dots,q \end{array}$$

variable  $x \in \mathbb{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^{\star}$ 

**Lagrangian:**  $\mathcal{L}: \mathbb{R}^n \times \mathbb{R}^p \times \mathbb{R}^q \to \mathbb{R}$ , with  $\operatorname{dom}(\mathcal{L}) = \mathcal{D} \times \mathbb{R}^p \times \mathbb{R}^q$ ,

$$\mathcal{L}(x,\lambda,\nu) = f_0(x) + \sum_{i=1}^p \lambda_i f_i(x) + \sum_{i=1}^q \nu_i h_i(x)$$

- weighted sum of objective and constraint functions
- $\triangleright$   $\lambda_i$  is the Lagrange multiplier (dual variable) associated with  $f_i(x) \leq 0$
- $\triangleright$   $\nu_i$  is the Lagrange multiplier (dual variable) associated with  $h_i(x)=0$

→ duality

## Karush-Kuhn-Tucker (KKT) Conditions

The following four conditions are called KKT conditions (for differentiable  $f_i$ ,  $h_i$ ):

- primal feasible:  $f_i(x) \leq 0, \quad i = 1, \dots, p$  $h_i(x) = 0, \quad i = 1, \dots, q$
- ▶ dual feasible:  $\lambda \succeq 0$
- ightharpoonup complementary slackness:  $\lambda_i f_i(x) = 0$  for  $i = 1, \dots, p$
- ightharpoonup gradient of Lagrangian with respect to x vanishes,

$$\nabla f_0(x) + \sum_{i=1}^p \lambda_i \nabla f_i(x) + \sum_{i=1}^q \nu_i \nabla h_i(x) = 0$$

Generalizes optimality condition  $\nabla f_0(x) = 0$  for unconstrained problems.

#### **Gradient Descent**

minimize 
$$f_0(x)$$

- $ightharpoonup f_0$  convex, twice continuously differentiable
- we assume optimal value  $p^* = \inf_x f_0(x)$  is attained (and finite)

#### **Gradient Descent**

minimize 
$$f_0(x)$$

- $ightharpoonup f_0$  convex, twice continuously differentiable
- we assume optimal value  $p^* = \inf_x f_0(x)$  is attained (and finite)

#### Gradient descent:

- 1. **given** a starting point  $x \in \text{dom}(f_0)$
- 2. **repeat**  $x := x t\nabla f_0(x)$ . (choose step size, t)
- 3. **until** stopping criterion satisfied, e.g.,  $\|\nabla f_0(x)\|_2 \leq \epsilon$ .
- ightharpoonup variants of gradient descent define step direction  $\Delta x$  different to  $-\nabla f_0(x)$

#### **Choosing Step Size**

**fixed schedule:** set t to a small constant or decay with each iteration

exact line search:  $t = \operatorname{argmin}_{t>0} f_0(x + t\Delta x)$ 

backtracking line search (with parameters  $\alpha \in (0, 1/2), \beta \in (0, 1)$ )

ightharpoonup starting at t=1 with search direction  $\Delta x$ , repeat  $t:=\beta t$  until

$$f_0(x + t\Delta x) < f_0(x) + \alpha t \nabla f_0(x)^T \Delta x$$

#### **Choosing Step Size**

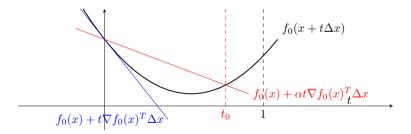
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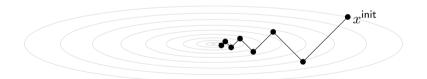
$$f_0(x + t\Delta x) < f_0(x) + \alpha t \nabla f_0(x)^T \Delta x$$



#### Example

Gradient descent (even with exact line search) can be slow. E.g.,

$$f_0(x) = x_1^2 + \gamma x_2^2, \quad \gamma \gg 1$$



#### Newton's Method

$$\Delta x_{\mathsf{nt}} = -\nabla^2 f_0(x)^{-1} \nabla f_0(x)$$

 $ightharpoonup x + \Delta x_{
m nt}$  minimizes the second-order approximation of  $f_0$  at x,

$$\hat{f}(x+v) = f_0(x) + \nabla f_0(x)^T v + \frac{1}{2} v^T \nabla^2 f_0(x) v$$

#### Newton's method:

- 1. **given** a starting point  $x \in \text{dom}(f_0)$ .
- 2. **repeat**  $x := x + t\Delta x_{nt}$ . (choose step size, t)
- 3. until stopping criterion satisfied.

### **Equality Constrained Methods**

minimize 
$$f_0(x)$$
  
subject to  $Ax = b$ 

- $ightharpoonup f_0$  convex, twice continuously differentiable
- $lacksquare A \in \mathbb{R}^{q imes n}$  with  $\mathrm{rank}(A) = q$  (and  $b \in \mathrm{range}(A)$ )
- ightharpoonup we assume  $p^*$  is finite and attained

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- $\triangleright$  we assume  $p^*$  is finite and attained

**optimality condition:**  $x^*$  is optimal iff there exists a  $\nu^*$  such that

$$\nabla f_0(x^*) + A^T \nu^* = 0, \quad Ax^* = b$$

## Newton Step for Equality Constrained Optimisation

Newton step  $\Delta x_{\rm nt}$  of  $f_0$  at feasible x is given by solution v of

$$\begin{bmatrix} \nabla^2 f_0(x) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} v \\ w \end{bmatrix} = \begin{bmatrix} -\nabla f_0(x) \\ 0 \end{bmatrix}$$

- second row ensures that x iterates stay feasible
- solves quadratic approximation of optimisation problem

minimize 
$$\hat{f}(x+v) \triangleq f_0(x) + \nabla f_0(x)^T v + \frac{1}{2} v^T \nabla^2 f_0(x) v$$
 subject to 
$$A(x+v) = b$$

solves linear approximation of optimality condition

For inequality constrained problems,

```
\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\dots,p \\ & Ax = b \end{array}
```

For inequality constrained problems,

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\dots,p \\ & Ax = b \end{array}$$

we reformulate using an indicator function,

minimize 
$$f_0(x) + \sum_{i=1}^p I_{\mathbb{R}_-}(f_i(x))$$
  
subject to  $Ax = b$ 

where  $I_{\mathbb{R}_{-}}(u)=0$  if  $u\leq 0$  and  $I_{\mathbb{R}_{-}}(u)=\infty$  otherwise,

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minimize 
$$f_0(x) - \frac{1}{t} \sum_{i=1}^p \log(-f_i(x))$$
 subject to  $Ax = b$ 

to get an equality constrained approximation.

For inequality constrained problems,

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i=1,\dots,p \\ & Ax = b \end{array}$$

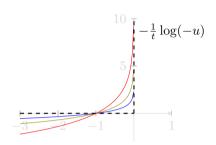
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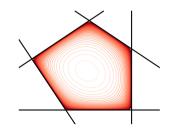
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to get an equality constrained approximation.





## Algorithms for Large Scale Problems

- ▶ for large scale problems, e.g., deep learning, Newton's method is too expensive
- even computing the true gradient may be too expensive
- lacktriangle many loss functions in machine learning decompose over train data  $\{(x_i,y_i)\}_{i=1}^m$ ,

$$L(\theta) = \sum_{i=1}^{m} \ell(f(x_i; \theta), y_i)$$

▶ SGD approximates the gradient on mini-batches  $\mathcal{I} \subseteq \{1, \dots, m\}$ 

$$\widehat{\nabla_{\theta}L} = \sum_{i \in \mathcal{I}} \nabla_{\theta} \ell(f(x_i; \theta), y_i)$$

- $\blacktriangleright$  under mild assumptions  $E\left[\widehat{\nabla_{\theta}L}\right] = \nabla_{\theta}L$
- ▶ for constrained problems can project back onto feasible set

Many, many other schemes and variations!

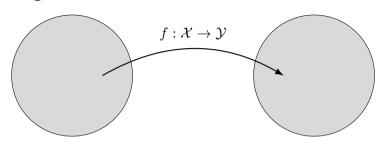
#### lecture 2

## Lecture 2: Differentiable Optimisation and Deep Learning



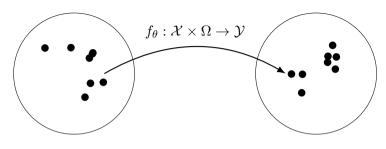
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# Machine Learning from 10,000ft



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## Machine Learning from 10,000ft



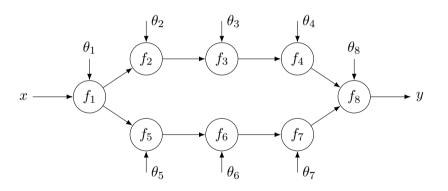
minimize (over  $\theta$ )  $\sum_{(x,y)\sim\mathcal{X}\times\mathcal{Y}} L(f_{\theta}(x),y)$ 

- ightharpoonup loss L what to do
- ightharpoonup model  $f_{\theta}$  how to do it
- optimised by gradient descent

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#### Deep Learning as an End-to-end Computation Graph

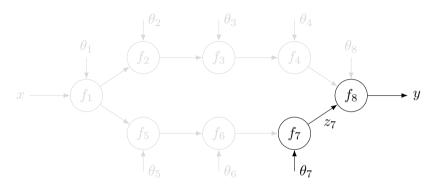
Deep learning does this by defining a function (equiv. computation graph) composed of many simple parametrized functions (equiv. computation nodes).



$$y = f_8(f_4(f_3(f_2(f_1(x)))), f_7(f_6(f_5(f_1(x)))))$$

(parameters  $\theta_i$  omitted for brevity)

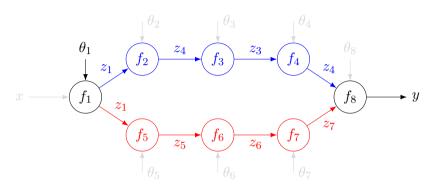
#### **Backward Pass**



#### Example 1.

$$\frac{\partial L}{\partial \theta_7} = \frac{\partial L}{\partial y} \frac{\partial y}{\partial z_7} \frac{\partial z_7}{\partial \theta_7}$$

#### **Backward Pass**

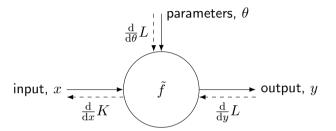


#### Example 2.

$$\frac{\partial L}{\partial \theta_1} = \frac{\partial L}{\partial y} \left( \frac{\partial y}{\partial z_4} \frac{\partial z_4}{\partial z_3} \frac{\partial z_3}{\partial z_2} \frac{\partial z_2}{\partial z_1} + \frac{\partial y}{\partial z_7} \frac{\partial z_7}{\partial z_6} \frac{\partial z_6}{\partial z_5} \frac{\partial z_5}{\partial z_4} \right) \frac{\partial z_1}{\partial \theta_1}$$

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## Deep Learning Node



Forward pass: compute output y as a function of the input x (and model parameters  $\theta$ ).

Backward pass: compute the derivative of the loss with respect to the input x (and model parameters θ) given the derivative of the loss with respect to the output y.

#### Notational Aside (Often Sloppy)

For scalar-valued functions:

total derivative:  $\frac{\mathrm{d}f}{\mathrm{d}x}$ 

partial derivative:  $\frac{\partial f}{\partial x}$ 

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#### Notational Aside (Often Sloppy)

For scalar-valued functions:

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partial derivative:  $\frac{\partial f}{\partial x}$ 

For multi-dimensional scalar-valued functions,  $f: \mathbb{R}^n \to \mathbb{R}$ :

$$\nabla f(x) = \left(\frac{\mathrm{d}f}{\mathrm{d}x_1}, \dots, \frac{\mathrm{d}f}{\mathrm{d}x_n}\right) \in \mathbb{R}^n$$

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For multi-dimensional vector-valued functions,  $f: \mathbb{R}^n \to \mathbb{R}^m$ :

$$\frac{\mathrm{d}}{\mathrm{d}x}f(x) = \begin{bmatrix} \frac{\mathrm{d}f_1}{\mathrm{d}x_1} & \cdots & \frac{\mathrm{d}f_1}{\mathrm{d}x_n} \\ \vdots & \ddots & \vdots \\ \frac{\mathrm{d}f_m}{\mathrm{d}x_n} & \cdots & \frac{\mathrm{d}f_m}{\mathrm{d}x_n} \end{bmatrix} \in \mathbb{R}^{m \times n} \qquad (\frac{\partial}{\partial x}f(x,y) \text{ for partial})$$

Sometimes D and  $D_X$  for  $\frac{d}{dx}$  and  $\frac{\partial}{\partial x}$ , respectively.

## Automatic Differentiation (AD)

- algorithmic procedure that produces code for computing exact derivatives
- assumes numeric computations are composed of a small set of elementary operations that we know how to differentiate
  - ▶ arithmetic, exp, log, trigonometric
- workhorse of modern machine learning that greatly reduces development effort

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## Automatic Differentiation (AD)

- algorithmic procedure that produces code for computing exact derivatives
- assumes numeric computations are composed of a small set of elementary operations that we know how to differentiate
  - ▶ arithmetic, exp, log, trigonometric
- workhorse of modern machine learning that greatly reduces development effort
- two flavours
  - (forward mode) propagates results on the first-order approximation  $x+\Delta x$  forward through the computations
  - (reverse mode) builds a program to compute derivative based on the chain rule re-using computation where applicable

$$\frac{\mathrm{d}L}{\mathrm{d}x} = \frac{\mathrm{d}L}{\mathrm{d}y} \frac{\mathrm{d}y}{\mathrm{d}x}$$

▶ different deep learning frameworks use slightly different approaches (explicit graph construction versus implicit operator tracking)

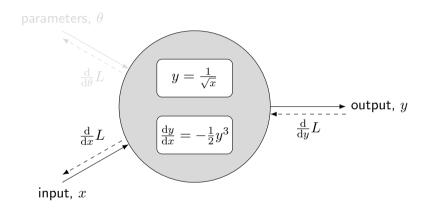
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# Computing $1/\sqrt{x}$

```
float Q_rsqrt( float number )
     long i;
     float x2, y;
    const float threehalfs = 1.5F;
     x2 = number * 0.5F;
     v = number:
     i = 0x5f3759df - (i >> 1); // what the f**k?
10
11
     v = * (float *) &i:
    y = y * (threehalfs - (x2 * y * y)); // 1st iter
     // y = y * (threehalfs - (x2 * y * y )); // 2nd iter, can be removed
13
14
15
     return v:
16
```

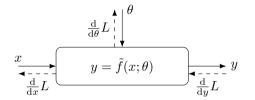
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## Separate Forward and Backward Operations



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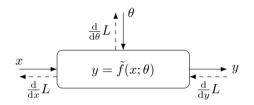
## Imperative vs Declarative Nodes



- imperative node
- input-output relationship explicit,

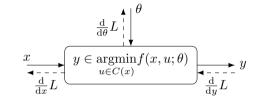
$$y = \tilde{f}(x;\theta)$$

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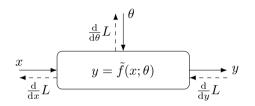
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- declarative node
- input-output relationship specified as solution to an optimisation problem,

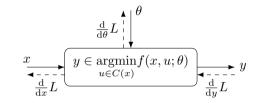
$$y \in \operatorname*{arg\,min}_{u \in C(x)} f(x, u; \theta)$$

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- input-output relationship specified as solution to an optimisation problem,

$$y \in \operatorname*{arg\,min}_{u \in C(x)} f(x, u; \theta)$$

can co-exist in the same computation graph (network)

# Average Pooling Example

$$\{x_i \in \mathbb{R}^m \mid i = 1, \dots, n\} \to \mathbb{R}^m$$

imperative specification

$$y = \frac{1}{n} \sum_{i=1}^{n} x_i$$

declarative specification

$$y = \operatorname{argmin}_{u \in \mathbb{R}^m} \sum_{i=1}^n \|u - x_i\|^2$$

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$$y = \operatorname{argmin}_{u \in \mathbb{R}^m} \sum_{i=1}^n \|u - x_i\|^2$$

can be easily varied, e.g., made robust

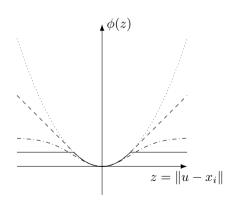
$$y = \operatorname{argmin}_{u \in \mathbb{R}^m} \sum_{i=1}^n \phi(u - x_i)$$

for some penalty function  $\phi$ 

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# Average Pooling Example

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for some penalty function  $\phi$ 

#### Bi-level Optimisation: Stackelberg Games

#### Consider two players, a leader and a follower

- ▶ the market dictates the price it's willing to pay for some goods based on supply, i.e., quantity produced by both players,  $P(q_1 + q_2)$
- ightharpoonup each player has a cost structure associated with producing goods,  $C_i(q_i)$  and wants to maximize profits,  $q_i P(q_1 + q_2) C_i(q_i)$
- ▶ the leader picks a quantity of goods to produce knowing that the follower will respond optimally. In other words, the leader solves

$$\begin{array}{ll} \text{maximize (over } q_1) & q_1P(q_1+q_2)-C_1(q_1) \\ \text{subject to} & q_2 \in \operatorname{argmax}_q qP(q_1+q)-C_2(q) \end{array}$$

```
\label{eq:local_equation} \begin{array}{ll} \text{minimize (over $x$)} & L(x,y) \\ \text{subject to} & y \in \operatorname{argmin}_{u \in C(x)} f(x,u) \end{array}
```

```
 \begin{array}{ll} \text{minimize (over $x$)} & L(x,y) \\ \text{subject to} & y \in \operatorname{argmin}_{u \in C(x)} f(x,u) \end{array}
```

ightharpoonup closed-form solution: substitute for y in upper-level problem (if possible)

```
minimize (over x) L(x, y(x))
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**closed-form solution:** substitute for y in upper-level problem (if possible)

```
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```

 convex lower-level problem: replace lower-level problem with sufficient optimality conditions (e.g., KKT conditions),

```
\begin{array}{ll} \mbox{minimize (over } x,y) & L(x,y) \\ \mbox{subject to} & h(x,y) = 0 \end{array}
```

$$\begin{array}{ll} \text{minimize (over $x$)} & L(x,y) \\ \text{subject to} & y \in \operatorname{argmin}_{u \in C(x)} f(x,u) \end{array}$$

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**pradient descent:** compute gradient of lower-level solution y with respect to x, and use the chain rule to get the total derivative,

$$x \leftarrow x - \eta \left( \frac{\partial L(x, y)}{\partial x} + \frac{\partial L(x, y)}{\partial y} \frac{\mathrm{d}y}{\mathrm{d}x} \right)$$

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**gradient descent:** compute gradient of lower-level solution y with respect to x, and use the chain rule to get the total derivative,

$$x \leftarrow x - \eta \left( \frac{\partial L(x, y)}{\partial x} + \frac{\partial L(x, y)}{\partial y} \frac{\mathrm{d}y}{\mathrm{d}x} \right)$$

by back-propagating through optimisation procedure or implicit differentiation

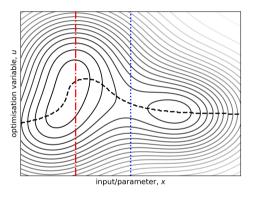
#### Parametrized Optimisation

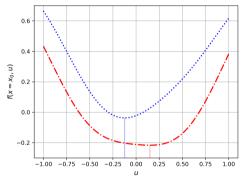
In the context of deep learning the upper-level Stackelberg problem is the **learning problem** and the lower-level Stackelberg problem is the **inference problem**.

A declarative node defines a family of problems indexed by continuous variable  $x \in \mathbb{R}^n$ ,

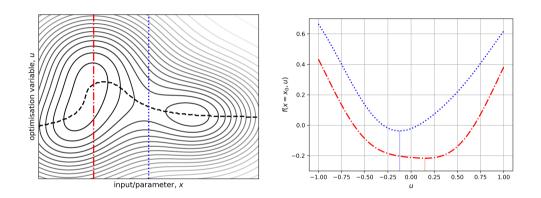
$$\left\{\begin{array}{ll} \text{minimize (over } u \in \mathbb{R}^m) & f_0(x,u) \\ \text{subject to} & f_i(x,u) \leq 0, \quad i=1,\dots,p \\ & h_i(x,u) = 0, \quad i=1,\dots,q \end{array}\right\}_{x \in \mathbb{R}^n}$$

## Parametrized Optimisation Example





## Parametrized Optimisation Example



**Main question:** How do we compute  $\frac{d}{dx} \operatorname{argmin}_u f(x, u)$ ?

#### Dini's Implicit Function Theorem

Consider the solution mapping associated with the equation f(x, u) = 0,

$$Y: x \mapsto \{u \in \mathbb{R}^m \mid f(x, u) = 0\} \text{ for } x \in \mathbb{R}^n.$$

We are interested in how elements of Y(x) change as a function of x.

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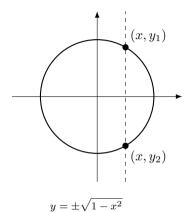
#### **Theorem**

Let  $f: \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}^m$  be differentiable in a neighbourhood of (x,u) and such that f(x,u)=0, and let  $\frac{\partial}{\partial u}f(x,u)$  be nonsingular. Then the solution mapping Y has a single-valued localization y around x for u which is differentiable in a neighbourhood  $\mathcal X$  of x with Jacobian satisfying

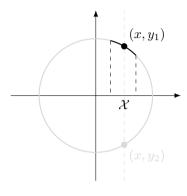
$$\frac{dy(x)}{dx} = -\left(\frac{\partial f(x,y(x))}{\partial y}\right)^{-1} \frac{\partial f(x,y(x))}{\partial x}$$

for every  $x \in \mathcal{X}$ .

#### Unit Circle Example



$$\frac{\mathrm{d}y}{\mathrm{d}x} = \frac{\mp 2x}{2\sqrt{1-x^2}} = -\frac{x}{3}$$



$$f(x,y) = x^2 + y^2 - 1$$

$$\frac{dy}{dx} = -\left(\frac{\partial f}{\partial y}\right)^{-1} \left(\frac{\partial f}{\partial x}\right)$$

$$= -\left(\frac{1}{2y}\right)(2x) = -\frac{x}{y}$$

# Differentiating Unconstrained Optimisation Problems

Let  $f: \mathbb{R} \times \mathbb{R} \to \mathbb{R}$  be twice differentiable and let

$$y(x) \in \operatorname{argmin}_u f(x, u)$$

then for non-zero Hessian

$$\frac{\mathrm{d}y(x)}{\mathrm{d}x} = -\left(\frac{\partial^2 f}{\partial y^2}\right)^{-1} \frac{\partial^2 f}{\partial x \partial y}.$$

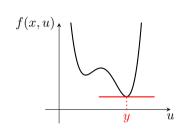
#### Differentiating Unconstrained Optimisation Problems

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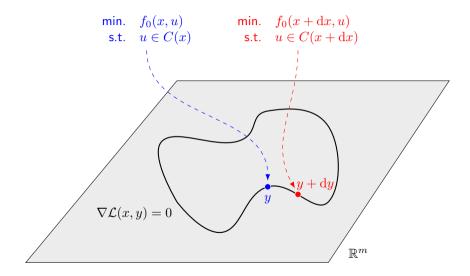


**Proof.** The derivative of f vanishes at (x,y), i.e.,  $y \in \operatorname{argmin}_u f(x,u) \implies \frac{\partial f(x,y)}{\partial y} = 0$ .

$$\begin{array}{ll} \mathsf{LHS}: & \frac{\mathrm{d}}{\mathrm{d}x} \frac{\partial f(x,y)}{\partial y} & = \frac{\partial^2 f(x,y)}{\partial x \partial y} + \frac{\partial^2 f(x,y)}{\partial y^2} \frac{\mathrm{d}y}{\mathrm{d}x} \\ \mathsf{RHS}: & \frac{\mathrm{d}}{\mathrm{d}x} 0 & = 0 \end{array}$$

Equating and rearranging gives the result.

#### Differentiable Optimisation: Big Picture Idea



#### Differentiating Equality Constrained Optimisation Problems

Consider functions  $f: \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}$  and  $h: \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}^q$ . Let

$$y(x) \in \mathop{\arg\min}_{u \in \mathbb{R}^m} f(x, u)$$
  
subject to  $h(x, u) = 0_q$ 

Assume that y(x) exists, that f and h are twice differentiable in the neighbourhood of (x,y(x)), and that  $\operatorname{rank}(\frac{\partial h(x,y)}{\partial y})=q$ .

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## Differentiating Equality Constrained Optimisation Problems

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Assume that y(x) exists, that f and h are twice differentiable in the neighbourhood of (x,y(x)), and that  $\mathrm{rank}(\frac{\partial h(x,y)}{\partial y})=q$ . Then for H non-singular

$$\frac{\mathrm{d}y(x)}{\mathrm{d}x} = H^{-1}A^{T}(AH^{-1}A^{T})^{-1}(AH^{-1}B - C) - H^{-1}B$$

where

$$\begin{split} A &= \frac{\partial h(x,y)}{\partial y} \in \mathbb{R}^{q \times m} \quad B = \frac{\partial^2 f(x,y)}{\partial x \partial y} - \sum_{i=1}^q \nu_i \frac{\partial^2 h_i(x,y)}{\partial x \partial y} \in \mathbb{R}^{m \times n} \\ C &= \frac{\partial h(x,y)}{\partial x} \in \mathbb{R}^{q \times n} \quad H = \frac{\partial^2 f(x,y)}{\partial y^2} - \sum_{i=1}^q \nu_i \frac{\partial^2 h_i(x,y)}{\partial y^2} \in \mathbb{R}^{m \times m} \end{split}$$

and  $\nu \in \mathbb{R}^q$  satisfies  $\nu^T A = \frac{\partial f(x,y)}{\partial y}$ .

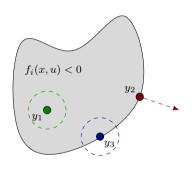
→ derivation

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## Dealing with Inequality Constraints

$$\begin{array}{c} y(x) \in \mathop{\arg\min}_{u \in \mathbb{R}^m} \ f_0(x,u) \\ \text{subject to} & h_i(x,u) = 0, \ i = 1,\dots,q \\ & f_i(x,u) \leq 0, \ i = 1,\dots,p. \end{array}$$

- Replace inequality constraints with log-barrier approximation (see last lecture)
- ► Treat as equality constraints if active  $(y_2 \text{ or } y_3)$  and ignore otherwise  $(y_1 \text{ or } y_3)$ 
  - ▶ may lead to one-sided gradients since  $\lambda \succeq 0$



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#### Automatic Differentiation for Differentiable Optimisation

- At one extreme we can try back propagate through the optimisation algorithm (i.e., unrolling the optimisation procedure using automatic differentiation)
- At the other extreme we can use the implicit differentiation result to hand-craft efficient backward pass code
- ► There are two options in between:
  - ▶ Use automatic differentiation to obtain quantities *A*, *B*, *C* and *H* from software implementations of the objective and (active) constraint functions
  - Implement the optimality condition  $\nabla \mathcal{L} = 0$  in software and automatically differentiate that

(in the next lecture we will see examples of the first two)

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#### Vector-Jacobian Product

For brevity consider the unconstrained optimisation case. The backward pass computes

$$\frac{\mathrm{d}L}{\mathrm{d}x} = \frac{\mathrm{d}L}{\mathrm{d}y} \frac{\mathrm{d}y}{\mathrm{d}x}$$
$$= \underbrace{(v^T)}_{\mathbb{R}^{1 \times m}} \underbrace{(-H^{-1}B)}_{\mathbb{R}^{m \times n}}$$

evaluation order: 
$$-v^T \left(H^{-1}B\right)$$
  $\left(-v^T H^{-1}\right) B$  
$$\cos t^{\dagger} \colon \quad O(m^2 n + mn) \qquad \qquad O(m^2 + mn)$$

 $^{\dagger}$  assumes  $H^{-1}$  is already factored (in  $O(m^3)$  if unstructured, less if structured)

#### Summary and Open Questions

- optimisation problems can be embedded inside deep learning models
- back-propagation by either unrolling the optimisation algorithm or implicit differentiation of the optimality conditions
  - ▶ the former is easy to implement using automatic differentiation but memory intensive
  - ightharpoonup the latter requires that solution be strongly convex locally (i.e., invertible H)
  - but does not need to know how the problem was solved, nor store intermediate forward-pass calculations
  - ightharpoonup computing  $H^{-1}$  may be costly

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### Summary and Open Questions

- optimisation problems can be embedded inside deep learning models
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  - but does not need to know how the problem was solved, nor store intermediate forward-pass calculations
  - ightharpoonup computing  $H^{-1}$  may be costly
- ▶ active area of research and many open questions
  - Are declarative nodes slower?
  - Do declarative nodes give theoretical guarantees?
  - ▶ How best to handle non-smooth or discrete optimization problems?
  - ▶ What about problems with multiple solutions?
  - ▶ What if the forward pass solution is suboptimal?
  - Can problems become infeasible during learning?

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### lecture 3

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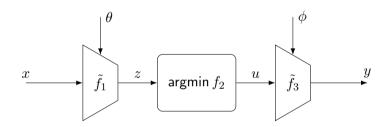
### Lecture 3: Examples and Applications



https://deepdeclarativenetworks.com

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### Common Theme



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### Differentiable Least Squares

Consider our old friend, the least-squares problem,

minimize 
$$||Ax - b||_2^2$$

parameterized by A and b and with closed-form solution  $x^* = (A^T A)^{-1} A^T b$ .

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### Differentiable Least Squares

Consider our old friend, the least-squares problem,

minimize 
$$||Ax - b||_2^2$$

parameterized by A and b and with closed-form solution  $x^* = (A^T A)^{-1} A^T b$ .

We are interested in derivatives of the solution with respect to the elements of A,

$$rac{\mathrm{d}x^{\star}}{\mathrm{d}A_{ij}} = rac{\mathrm{d}}{\mathrm{d}A_{ij}} \left(A^T\!A
ight)^{-1} A^T b \ \in \mathbb{R}^n$$

We could also compute derivatives with respect to elements of b (but not here).

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### Least Squares Backward Pass

The backward pass combines  $\frac{\mathrm{d}x^\star}{\mathrm{d}A_{ij}}$  with  $v^T=\frac{\mathrm{d}L}{\mathrm{d}x^\star}$  via the vector-Jacobian product. After some algebraic manipulation (see lecture notes) we get

$$\left(\frac{\mathrm{d}L}{\mathrm{d}A}\right)^T = wr^T - x^*(Aw)^T \in \mathbb{R}^{m \times n}$$

where  $w^T = v^T (A^T A)^{-1}$ .

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### Least Squares Backward Pass

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where  $w^T = v^T (A^T A)^{-1}$ .

- $(A^TA)^{-1}$  is used in both the forward and backward pass
- ightharpoonup factored once to solve for x, e.g., into A=QR
- $\triangleright$  cache R and re-use when computing gradients

→ derivation

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### Aside: PyTorch and Batched Data

Deep learning frameworks process data in batches, passed as tensors, for stochastic gradient descent. The first dimension of the tensor is the batch dimension.

**Example.** For the operation y = Ax + b we might have

$$X = \{x^{(1)}, \dots, x^{(K)}\}$$
 (input) 
$$Y = \{Ax^{(1)} + b, \dots, Ax^{(K)} + b\}$$
 (output)

Many PyTorch functions are batch-aware, e.g., torch.bmm. For many operations the einsum function and broadcasting are particularly useful, e.g.,

```
y = torch.einsum("ij,kj->ki", A, x) + b
```

computes  $y = Ax^{(k)} + b$  on each element k = 1, ..., K of the batch.

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## PyTorch Implementation: Forward Pass

```
class LeastSquaresFcn(torch.autograd.Function):
       """PvTorch autograd function for least squares."""
       Ostaticmethod
       def forward(ctx, A, b):
           B, M, N = A.shape
           assert b.shape == (B. M. 1)
           with torch.no_grad():
10
               Q, R = torch.linalg.gr(A, mode='reduced')
               x = torch.linalg.solve_triangular(R,
12
13
                   torch.bmm(b.view(B, 1, M), Q).view(B, N, 1), upper=True)
14
           # save state for backward pass
15
           ctx.save for backward(A, b, x, R)
16
           # return solution
           return x
```

$$A = QR$$
 
$$x = R^{-1} \left( Q^T b \right)$$
 (solves  $Rx = Q^T b$ )

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### PyTorch Implementation: Backward Pass

```
Ostaticmethod
        def backward(ctx. dx):
            # check for None tensors
            if dx is None:
5
6
7
8
9
10
                 return None, None
            # unpack cached tensors
            A, b, x, R = ctx.saved_tensors
            B, M, N = A.shape
            dA. db = None. None
13
            w = torch.linalg.solve triangular(R.
14
                 torch.linalg.solve_triangular(torch.transpose(R, 2, 1),
15
                 dx, upper=False), upper=True)
16
            Aw = torch.bmm(A, w)
17
18
            if ctx.needs_input_grad[0]:
19
                 r = b - torch.bmm(A, x)
                 dA = torch.einsum("bi,bj->bij", r.view(B,M), w.view(B,N)) - \
torch.einsum("bi,bj->bij", Aw.view(B,M), x.view(B,N))
20
            if ctx.needs_input_grad[1]:
                 dh = \Delta w
24
            # return gradients
            return dA. db
```

$$w = (A^{T}A)^{-1} v$$

$$= R^{-1} (R^{-T}v)$$

$$r = b - Ax$$

$$\left(\frac{dL}{dA}\right)^{T} = wr^{T} - x(Aw)^{T}$$

$$\left(\frac{dL}{db}\right)^{T} = Aw$$

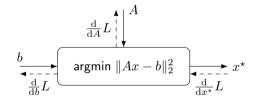
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### Example

Bi-level optimisation problem with lower-level least squares:

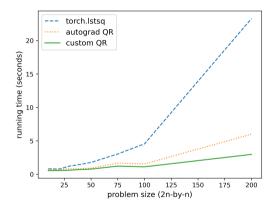
$$\begin{array}{ll} \text{minimize} & \frac{1}{2}\|x^\star - x^{\mathsf{target}}\|_2^2 \\ \text{subject to} & x^\star = \mathrm{argmin}_x \ \|Ax - b\|_2^2 \end{array}$$

with upper-level variable  $A \in \mathbb{R}^{m \times n}$ .



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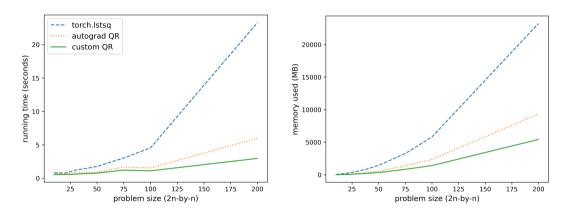
### **Profiling**



(problems with m=2n; run for 1000 iterations on CPU using PyTorch 1.13.0)

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### **Profiling**



(problems with m=2n; run for 1000 iterations on CPU using PyTorch 1.13.0)

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### **Optimal Transport**

One view of optimal transport is as a matching problem

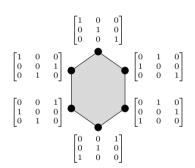
- ightharpoonup from an m-by-n cost matrix M
- ightharpoonup to an m-by-n probability matrix P,

often formulated with an entropic regularisation term,

$$\begin{array}{ll} \text{minimize} & \langle M,P\rangle + \frac{1}{\gamma}\langle P,\log P\rangle \\ \text{subject to} & P\mathbf{1} = r \\ & P^T\mathbf{1} = c \end{array}$$

with 
$$\mathbf{1}^{T}r = \mathbf{1}^{T}c = 1$$

The row and column sum constraints ensure that P is a doubly stochastic matrix (lies within the convex hull of permutation matrices).



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### Solving Entropic Optimal Transport

Solution takes the form

$$P_{ij} = \alpha_i \beta_j e^{-\gamma M_{ij}}$$

and can be found using the Sinkhorn algorithm,

- ▶ Set  $K_{ij} = e^{-\gamma M_{ij}}$  and  $\alpha, \beta \in \mathbb{R}^n_{++}$
- Iterate until convergence,

$$\alpha \leftarrow r \oslash K\beta$$
$$\beta \leftarrow c \oslash K^T \alpha$$

where  $\oslash$  denotes componentwise division

 $\blacktriangleright \ \, \mathsf{Return} \,\, P = \mathbf{diag}(\alpha) K \mathbf{diag}(\beta)$ 

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### Differentiable Optimal Transport

▶ Option 1: back-propagate through Sinkhorn algorithm

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### Differentiable Optimal Transport

- ▶ Option 1: back-propagate through Sinkhorn algorithm
- ▶ Option 2: use the implicit differentiation result

$$\underbrace{\frac{\mathrm{d}L}{\mathrm{d}M}}_{m\text{-by-}n} = \underbrace{\frac{\mathrm{d}L}{\mathrm{d}P}}_{m\text{-by-}n} \underbrace{\frac{\mathrm{d}P}{\mathrm{d}M}}_{m\text{-by-}n}$$

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### Differentiable Optimal Transport

- ▶ Option 1: back-propagate through Sinkhorn algorithm
- ▶ Option 2: use the implicit differentiation result



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### Optimal Transport Gradient

Derivation of the optimal transport gradient is quite tedious (see notes). The result:

$$\begin{split} \frac{\mathrm{d}L}{\mathrm{d}M} &= \frac{\mathrm{d}L}{\mathrm{d}P} \left( H^{-1} \mathbf{A}^T \left( A H^{-1} A^T \right)^{-1} \mathbf{A} H^{-1} - H^{-1} \right) B \\ &= \gamma \frac{\mathrm{d}L}{\mathrm{d}P} \mathrm{diag}(P) \begin{bmatrix} A_1 \\ A_2 \end{bmatrix}^T \begin{bmatrix} \Lambda_{11} & \Lambda_{12} \\ \Lambda_{12}^T & \Lambda_{22} \end{bmatrix} \begin{bmatrix} A_1 \\ A_2 \end{bmatrix} \mathrm{diag}(P) - \gamma \frac{\mathrm{d}L}{\mathrm{d}P} \mathrm{diag}(P) \end{split}$$

where

$$\begin{bmatrix} A_1 \\ A_2 \end{bmatrix} = \begin{bmatrix} \mathbf{0}_n^T & \mathbf{1}_n^T & \dots & \mathbf{0}_n^T \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_n^T & \mathbf{0}_n^T & \dots & \mathbf{1}_n^T \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \qquad \begin{pmatrix} AH^{-1}A^T \end{pmatrix}^{-1} = \frac{1}{\gamma} \begin{bmatrix} \Lambda_{11} & \Lambda_{12} \\ \Lambda_{12}^T & \Lambda_{22} \end{bmatrix} \\ & = \frac{1}{\gamma} \begin{bmatrix} \operatorname{diag}(r_{2:m}) & P_{2:m,1:n} \\ P_{2:m,1:n}^T & \operatorname{diag}(c) \end{bmatrix}^{-1}$$

▶ derivation

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### **Implementation**

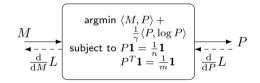
```
@staticmethod
   def backward(ctx, dJdP)
       # unpacked cached tensors
      M, r, c, P = ctx.saved_tensors
       batches, m, n = P.shape
       # initialize backward gradients (-v^T H^{-1} B)
8
       dLdM = -1.0 * gamma * P * dLdP
9
10
       # compute [vHAt1, vHAt2] = -v^T H^{-1} A^T
11
       vHAt1, vHAt2 = sum(dJdM[:, 1:m, 0:n], dim=2), sum(dJdM, dim=1)
13
       # compute [v1, v2] = -v^T H^{-1} A^T (A H^{-1} A^T)^{-1}
14
       P over c = P[:, 1:m, 0:n] / c.view(batches, 1, n)
       lmd_11 = cholesky(diag_embed(r[:, 1:m]) - einsum("bij,bkj->bik", P[:, 1:m, 0:n], P_over_c))
15
16
       lmd_12 = cholesky_solve(P_over_c, lmd_11)
17
       lmd_22 = diag_embed(1.0 / c) + einsum("bji,bjk->bjk", lmd_12, P_over_c)
18
19
       v1 = choleskv_solve(vHAt1.view(batches, m-1, 1), lmd_11).view(batches, m-1) -
20
           einsum("bi,bii->bi", vHAt2, lmd_12)
21
       v2 = einsum("bi,bij->bj", vHAt2, lmd_22) - einsum("bi,bij->bj", vHAt1, lmd_12)
23
       # compute v^T H^{-1} A^T (A H^{-1] A^T)^{-1} A H^{-1} B - v^T H^{-1} B
24
       dLdM[:. 1:m. 0:n] -= v1.view(batches. m-1. 1) * P[:. 1:m. 0:n]
25
       dJdM -= v2.view(batches, 1, n) * P
26
27
       # return gradients
28
       return didM
```

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### Experiment

Bi-level optimisation problem with lower-level optimal transport problem:

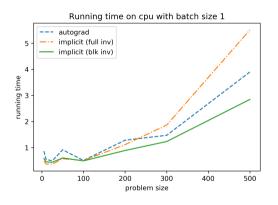
minimize 
$$\frac{1}{2}\|P-P^{\mathsf{target}}\|_F^2$$
 subject to minimize  $\langle M,P\rangle+\frac{1}{\gamma}\langle P,\log P\rangle$  subject to 
$$P\mathbf{1}=\frac{1}{n}\mathbf{1}$$
 
$$P^T\mathbf{1}=\frac{1}{m}\mathbf{1}$$

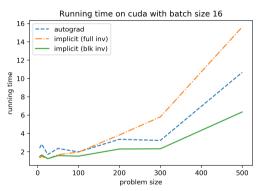


with upper-level variable  $M \in \mathbb{R}^{m \times n}$ .

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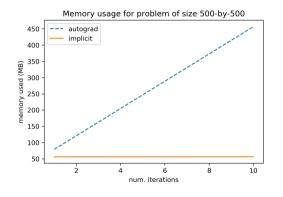
### Results: Running Time

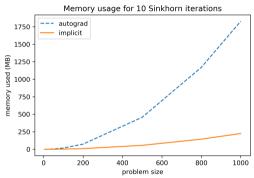




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## Results: Memory Usage





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### Application to Blind Perspective-n-Point



find the location where the photograph was taken

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### Coupled Problem



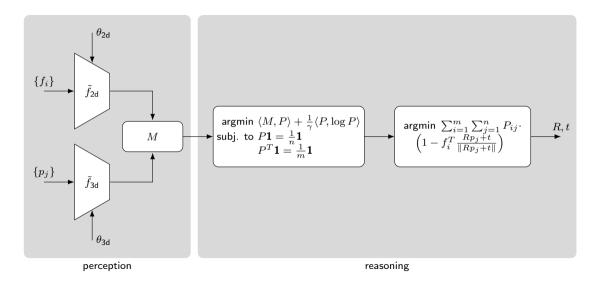
▶ if we knew correspondences then determining camera pose would be easy



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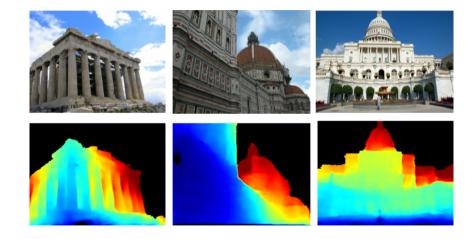
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### Blind Perspective-n-Point Network Architecture



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# Blind Perspective-n-Point Results



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#### Further Resources

#### Where to from here?

- ▶ Deep declarative networks (http://deepdeclarativenetworks.com)
  - lots of small code examples and tutorials
- CVXPyLayers (https://github.com/cvxgrp/cvxpylayers)
- ► Theseus (https://sites.google.com/view/theseus-ai)
- JAXopt (https://github.com/google/jaxopt)

lecture notes available at https://users.cecs.anu.edu.au/~sgould

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### break-out

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### Local Optima are Global Optima Proof Plack

#### any local minimum of a convex problem is (globally) optimal

**Proof.** Suppose that x is locally optimal, but there exists a feasible y with lower objective, i.e.,  $f_0(y) < f_0(x)$ . Local optimality of x means there must be an R > 0 such that

$$z$$
 feasible and  $||z - x||_2 \le R \implies f_0(z) \ge f_0(x)$ 

Consider  $z=\theta y+(1-\theta)x$  with  $\theta=\frac{R}{2\|y-x\|_2}$ . We have that  $\|y-x\|_2>R$  since we assumed  $f_0(y)< f_0(x)$ , so  $0<\theta<1/2<1$ . Therefore z is a convex combination of two feasible points, hence also feasible. Moreover,  $\|z-x\|_2=R/2$  (from our choice of  $\theta$ ) and therefore  $f_0(z)\geq f_0(x)$  by our assumption that x is locally optimal. But

$$f_0(z) \le \theta f_0(y) + (1 - \theta) f_0(x) < \theta f_0(x) + (1 - \theta) f_0(x) = f_0(x)$$

where the first inequality is by the definition of convex function and the second inequality is from our assumption that  $f_0(y) < f_0(x)$ . We have a contradiction. Therefore every locally optimal point is globally optimal.

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### automatic differentiation

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### Toy Example: Babylonian Algorithm back

Consider the following implementation for a forward operation:

```
1: procedure \operatorname{FWDFCN}(x)

2: y_0 \leftarrow \frac{1}{2}x

3: for t = 1, \dots, T do

4: y_t \leftarrow \frac{1}{2}\left(y_{t-1} + \frac{x}{y_{t-1}}\right)

5: end for

6: return y_T

7: end procedure
```

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### Toy Example: Babylonian Algorithm back



Consider the following implementation for a forward operation:

```
1: procedure FWDFCN(x)
       y_0 \leftarrow \frac{1}{2}x
      for t = 1, \dots, T do y_t \leftarrow \frac{1}{2} \left( y_{t-1} + \frac{x}{y_{t-1}} \right)
         end for
          return u_T
  end procedure
```

Automatic differentiation algorithmically generates the backward code:

```
1: procedure BCKFCN(x, y_T, \frac{dL}{dy_T})
2: \frac{\mathrm{d}L}{\mathrm{d}x} \leftarrow 0
3: \mathbf{for} \ t = T, \dots, 1 \ \mathbf{do}
4: \frac{\mathrm{d}L}{\mathrm{d}x} \leftarrow \frac{\mathrm{d}L}{\mathrm{d}x} + \frac{\mathrm{d}L}{\mathrm{d}y_t} \left(\frac{1}{2y_{t-1}}\right)
5: \frac{\mathrm{d}L}{\mathrm{d}y_{t-1}} \leftarrow \frac{\mathrm{d}L}{\mathrm{d}y_t} \left(\frac{1}{2} - \frac{x}{2y_{t-1}^2}\right)
    7: \frac{\mathrm{d}L}{\mathrm{d}x} \leftarrow \frac{\mathrm{d}L}{\mathrm{d}x} + \frac{\mathrm{d}L}{\mathrm{d}u_0} \frac{1}{2}
                   return \frac{\mathrm{d}L}{\mathrm{d}z}
      9: end procedure
```

### Toy Example: Babylonian Algorithm •• back

Consider the following implementation for a forward operation:

```
1: procedure \operatorname{FWDFCN}(x)

2: y_0 \leftarrow \frac{1}{2}x

3: for t=1,\ldots,T do

4: y_t \leftarrow \frac{1}{2}\left(y_{t-1} + \frac{x}{y_{t-1}}\right)

5: end for

6: return y_T

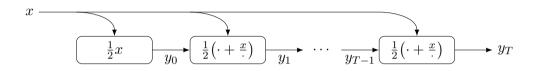
7: end procedure
```

- ightharpoonup computes  $y = \sqrt{x}$
- derivative computed directly is  $\frac{dy}{dx} = \frac{1}{2\sqrt{x}} = \frac{1}{2y}$

Automatic differentiation algorithmically generates the backward code:

```
1: procedure BCKFCN(x, y_T, \frac{dL}{dy_T})
2: \frac{\mathrm{d}L}{\mathrm{d}x} \leftarrow 0
3: \mathbf{for} \ t = T, \dots, 1 \ \mathbf{do}
4: \frac{\mathrm{d}L}{\mathrm{d}x} \leftarrow \frac{\mathrm{d}L}{\mathrm{d}x} + \frac{\mathrm{d}L}{\mathrm{d}y_t} \left(\frac{1}{2y_{t-1}}\right)
5: \frac{\mathrm{d}L}{\mathrm{d}y_{t-1}} \leftarrow \frac{\mathrm{d}L}{\mathrm{d}y_t} \left(\frac{1}{2} - \frac{x}{2y_{t-1}^2}\right)
                                                                                                              \partial u_{\pm}/\partial u_{\pm} = 1
     7: \frac{dL}{dx} \leftarrow \frac{dL}{dx} + \frac{dL}{du_0} \frac{1}{2}
                  return \frac{dL}{dx}
             end procedure
```

# Computation Graph for Babylonian Algorithm Pack



$$y_T = f(x, f(x, f(x, \dots f(x, \frac{1}{2}x))))$$
 with  $f(x, y) = \frac{1}{2} \left(y + \frac{x}{y}\right)$ 

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## duality

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## Lagrange Dual Function Pack

Define Lagrange dual function,  $g: \mathbb{R}^p \times \mathbb{R}^q \to \mathbb{R}$ , as

$$g(\lambda, \nu) = \inf_{x \in \mathcal{D}} \mathcal{L}(x, \lambda, \nu)$$
$$= \inf_{x \in \mathcal{D}} \left( f_0(x) + \sum_{i=1}^p \lambda_i f_i(x) + \sum_{i=1}^q \nu_i h_i(x) \right)$$

- ightharpoonup q is concave (always), can be  $-\infty$  for some  $\lambda, \nu$
- ▶ lower bound property: if  $\lambda \succeq 0$ , then  $g(\lambda, \nu) \leq p^*$  (since for feasible x we have  $f_i(x) \leq 0$  and  $h_i(x) = 0$ )

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#### The Dual Problem • back

The Lagrange dual problem is to maximise the dual function

maximize 
$$g(\lambda, \nu)$$
 subject to  $\lambda \succeq 0$ 

- $\triangleright$  finds the best lower bound on  $p^*$ , obtained from Lagrange dual function
- ightharpoonup a convex optimisation problem with optimal value denoted by  $d^{\star}$
- $\triangleright$   $\lambda, \nu$  are dual feasible if  $\lambda \succeq 0$  and  $(\lambda, \nu) \in \mathbf{dom}(g)$
- original problem is known as the primal problem

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## Weak and Strong Duality back

#### weak duality: $d^{\star} \leq p^{\star}$

- always holds (for convex and nonconvex problems)
- can be used to find nontrivial lower bounds for difficult problems

#### strong duality: $d^* = p^*$

- does not hold in general
- ► (usually) holds for convex problems
- conditions that guarantee strong duality on convex problems are called constraint qualifications

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#### differentiating equality constrained problems

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#### Abridged Derivation back

Forming the Lagrangian at optimal y for fixed x we have

$$\mathcal{L}(x, y, \nu) = f(x, y) - \sum_{i=1}^{q} \nu_i h_i(x, y).$$

Since  $\frac{\partial h(x,y)}{\partial y}$  is full rank we have that y is a regular point. Then there exists a  $\nu$  such that the Lagrangian is stationary at the point  $(y,\nu)$ . Thus

$$\begin{bmatrix} \frac{\partial \mathcal{L}}{\partial Y}^T \\ \frac{\partial \mathcal{L}}{\partial \nu}^T \end{bmatrix} = \begin{bmatrix} \left( \frac{\partial f(x,y)}{\partial y} - \sum_{i=1}^q \nu_i \frac{\partial h_i(x,y)}{\partial y} \right)^T \\ h(x,y) \end{bmatrix} = \mathbf{0}_{m+q}$$

which we can differentiate with respect to x,

$$\frac{\mathrm{d}}{\mathrm{d}x} \left[ \left( \frac{\partial f(x,y)}{\partial y} \right)^T - \sum_{i=1}^q \nu_i \left( \frac{\partial h_i(x,y)}{\partial y} \right)^T \right] = \mathbf{0}_{(m+q) \times n}$$

to get (after some re-arranging in matrix form)

$$\begin{bmatrix} \frac{\partial^2 f(x,y)}{\partial y^2} - \sum_{i=1}^q \nu_i \frac{\partial^2 h_i(x,y)}{\partial y^2} & -(\frac{\partial h(x,y)}{\partial y})^T \\ \frac{\partial h(x,y)}{\partial y} & \mathbf{0}_{q \times q} \end{bmatrix} \begin{bmatrix} \frac{\mathrm{d}y(x)}{\mathrm{d}x} \\ \frac{\mathrm{d}\nu(x)}{\mathrm{d}x} \end{bmatrix} = -\begin{bmatrix} \frac{\partial^2 f(x,y)}{\partial x \partial y} - \sum_{i=1}^q \nu_i \frac{\partial^2 h_i(x,y)}{\partial x \partial y} \\ \frac{\partial}{\partial x} h(x,y) \end{bmatrix}.$$

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#### Abridged Derivation Pack

Forming the Lagrangian at optimal y for fixed x we have

$$\mathcal{L}(x, y, \nu) = f(x, y) - \sum_{i=1}^{q} \nu_i h_i(x, y).$$

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to get (after some re-arranging in matrix form)

$$\begin{bmatrix} H & -A^T \\ A & \mathbf{0}_{q \times q} \end{bmatrix} \begin{bmatrix} \frac{\mathrm{d}y(x)}{\mathrm{d}x} \\ \frac{\mathrm{d}\nu(x)}{\mathrm{d}x} \end{bmatrix} = - \begin{bmatrix} B \\ C \end{bmatrix}.$$

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## Abridged Derivation (cont.) Phack

(from last slide:)

$$\begin{bmatrix} H & -A^T \\ A & \mathbf{0}_{q \times q} \end{bmatrix} \begin{bmatrix} \frac{\mathrm{d}y(x)}{\mathrm{d}x} \\ \frac{\mathrm{d}\nu(x)}{\mathrm{d}x} \end{bmatrix} = - \begin{bmatrix} B \\ C \end{bmatrix}$$

We can solve this system of equations directly or solve by variable elimination. Multiplying out we have

$$H\frac{\mathrm{d}y(x)}{\mathrm{d}x} - A^T \frac{\mathrm{d}\nu(x)}{\mathrm{d}x} = -B \tag{1}$$

$$A\frac{\mathrm{d}y(x)}{\mathrm{d}x} = -C\tag{2}$$

Substituting  $\frac{dy(x)}{dx}$  from (1) into (2) gives,

$$\overbrace{AH^{-1}(A^T \frac{d\nu(x)}{dx} - B)}^{\frac{dy(x)}{dx}} = -C$$

$$\therefore \frac{d\nu(x)}{dx} = \left(AH^{-1}A^T\right)^{-1} \left(AH^{-1}B - C\right)$$

Then substituting back into (1) we get the result

$$\frac{\mathrm{d}y(x)}{\mathrm{d}x} = H^{-1}A^{T} \left(AH^{-1}A^{T}\right)^{-1} \left(AH^{-1}B - C\right) - H^{-1}B$$

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#### least squares

#### 

Differentiating  $x^*$  with respect to single element  $A_{ij}$ , we have

$$\frac{\mathsf{d}}{\mathsf{d}A_{ij}}x^* = \frac{\mathsf{d}}{\mathsf{d}A_{ij}} \left(A^T A\right)^{-1} A^T b$$

$$= \left(\frac{\mathsf{d}}{\mathsf{d}A_{ij}} \left(A^T A\right)^{-1}\right) A^T b + \left(A^T A\right)^{-1} \left(\frac{\mathsf{d}}{\mathsf{d}A_{ij}} A^T b\right)$$

Using the identity  $\frac{d}{dz}Z^{-1} = -Z^{-1}\left(\frac{d}{dz}Z\right)Z^{-1}$  we get, for the first term,

$$\frac{d}{dA_{ij}} (A^T A)^{-1} = -(A^T A)^{-1} \left( \frac{d}{dA_{ij}} (A^T A) \right) (A^T A)^{-1}$$
$$= -(A^T A)^{-1} (E_{ij}^T A + A^T E_{ij}) (A^T A)^{-1}$$

where  $E_{ij}$  is a matrix with one in the (i, j)-th element and zeros elsewhere. Furthermore, for the second term,

$$\frac{\mathsf{d}}{\mathsf{d}A_{ij}}A^Tb = E_{ij}^Tb$$

## Least Squares Backward Pass Derivation (cont.)

Plugging these back into parent equation we have

$$\frac{d}{dA_{ij}}x^* = -(A^TA)^{-1}(E_{ij}^TA + A^TE_{ij})(A^TA)^{-1}A^Tb + (A^TA)^{-1}E_{ij}^Tb$$

$$= -(A^TA)^{-1}(E_{ij}^TA + A^TE_{ij})x^* + (A^TA)^{-1}E_{ij}^Tb$$

$$= -(A^TA)^{-1}(E_{ij}^T(Ax^* - b) + A^TE_{ij}x^*)$$

$$= -(A^TA)^{-1}((a_i^Tx^* - b_i)e_j + x_j^*a_i)$$

where  $e_j = (0, 0, \dots, 1, 0, \dots) \in \mathbb{R}^n$  is the j-th canonical vector, i.e., vector with a one in the j-th component and zeros everywhere else, and  $a_i^T \in \mathbb{R}^{1 \times n}$  is the i-th row of matrix A.

#### Least Squares Backward Pass Derivation (cont.)

Let  $r = b - Ax^*$  and let  $v^T$  denote the backward coming gradient  $\frac{d}{dx^*}L$ . Then

$$\frac{dL}{dA_{ij}} = v^T \frac{dx^*}{dA_{ij}}$$

$$= v^T (A^T A)^{-1} (r_i e_j - x_j^* a_i)$$

$$= w^T (r_i e_j - x_j^* a_i)$$

$$= r_i w_j - w^T a_i x_j^*$$

where  $w = (A^T A)^{-1} v$ . We can compute the entire matrix of  $m \times n$  derivatives efficiently as the sum of outer products

$$\left(\frac{\mathrm{d}L}{\mathrm{d}A}\right)^T = \left[\frac{\mathrm{d}L}{\mathrm{d}A_{ij}}\right]_{\substack{i=1,\dots,m\\i=1,\dots,n}} = wr^T - x^*(Aw)^T$$

#### optimal transport

## Objective and Constraint Functions back

$$f(M, P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij}$$

$$h(M,P) = \begin{bmatrix} \mathbf{1}_n^T & \mathbf{0}_n^T & \dots & \mathbf{0}_n^T \\ \mathbf{0}_n^T & \mathbf{1}_n^T & \dots & \mathbf{0}_n^T \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_n^T & \mathbf{0}_n^T & \dots & \mathbf{1}_n^T \\ I_{n\times n} & I_{n\times n} & \dots & I_{n\times n} \end{bmatrix} \begin{bmatrix} P_{11} \\ P_{12} \\ \vdots \\ P_{1n} \\ P_{21} \\ \vdots \\ P_{mn} \end{bmatrix} - \begin{bmatrix} \mathbf{r_1} \\ r_2 \\ \vdots \\ r_m \\ c_1 \\ \vdots \\ c_n \end{bmatrix}$$

(one constraint is redundant—a linear combination of the others—and removed to ensure  ${\bf rank}(A)=q)$ 

$$f(M,P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij} \qquad h(M,P) = \begin{bmatrix} \mathbf{0}_{n}^{1} & \mathbf{1}_{n}^{1} & \dots & \mathbf{0}_{n}^{1} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \vec{P} - \begin{bmatrix} r_{2} \\ \vdots \\ r_{m} \\ c \end{bmatrix}$$

$$f(M,P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij} \qquad h(M,P) = \begin{bmatrix} \mathbf{0}_{n}^{-} & \mathbf{1}_{n}^{-} & \dots & \mathbf{0}_{n}^{-} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \vec{P} - \begin{bmatrix} r_{2} \\ \vdots \\ r_{m} \\ c \end{bmatrix}$$

$$\frac{dP}{dM} = \left(H^{-1}A^{T}(AH^{-1}A^{T})^{-1}AH^{-1} - H^{-1}\right)B$$

$$A = \frac{\mathrm{d}}{\mathrm{d}P}h \in \mathbb{R}^{(m+n-1)\times mn} \qquad B = \frac{\mathrm{d}^2}{\mathrm{d}M^2P}f \in \mathbb{R}^{mn\times nn} \quad H = \frac{\mathrm{d}^2}{\mathrm{d}P^2}f \in \mathbb{R}^{mn\times mn}$$

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 $= egin{bmatrix} oldsymbol{0}_n^T & oldsymbol{1}_n^T & \dots & oldsymbol{0}_n^T \ oldsymbol{0}_n^T & oldsymbol{0}_n^T & \dots & oldsymbol{1}_n^T \ I_{n imes n} & I_{n imes n} & \dots & I_{n imes n} \end{bmatrix}$ 

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$$f(M,P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij} \qquad h(M,P) = \begin{bmatrix} \mathbf{0}_{n}^{T} & \mathbf{1}_{n}^{T} & \dots & \mathbf{0}_{n}^{T} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \vec{P} - \begin{bmatrix} r_{2} \\ \vdots \\ r_{m} \\ c \end{bmatrix}$$

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 $= egin{bmatrix} oldsymbol{0}_n^T & oldsymbol{1}_n^T & \dots & oldsymbol{0}_n^T \ oldsymbol{0}_n^T & oldsymbol{0}_n^T & \dots & oldsymbol{1}_n^T \ oldsymbol{I}_{n imes n} & I_{n imes n} & \dots & I_{n imes n} \end{bmatrix} \qquad B = rac{\mathrm{d}^2}{\mathrm{d}M\partial P} oldsymbol{f} \in \mathbb{R}^{mn imes mn} \qquad H = rac{\mathrm{d}^2}{\mathrm{d}P^2} oldsymbol{f} \in \mathbb{R}^{mn imes mn} \qquad H = rac{\mathrm{d}^2}{\mathrm{d}P^2} oldsymbol{f} \in \mathbb{R}^{mn imes mn}$ 

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 $= egin{bmatrix} oldsymbol{0}_n^T & oldsymbol{1}_n^T & \dots & oldsymbol{0}_n^T \ oldsymbol{0}_n^T & oldsymbol{0}_n^T & \dots & oldsymbol{1}_n^T \ I_{n imes n} & I_{n imes n} & \dots & I_{n imes n} \end{bmatrix} \qquad = I_{mn imes mn}$ 

$$f(M,P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij} \qquad h(M,P) = \begin{bmatrix} \mathbf{0}_{n}^{1} & \mathbf{1}_{n}^{1} & \dots & \mathbf{0}_{n}^{1} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \vec{P} - \begin{bmatrix} r_{2} \\ \vdots \\ r_{m} \\ c \end{bmatrix}$$

$$\frac{dP}{dM} = \left(H^{-1}A^{T}(AH^{-1}A^{T})^{-1}AH^{-1} - H^{-1}\right)B$$

$$A = \frac{\mathrm{d}}{\mathrm{d}P} \overset{\pmb{h}}{h} \in \mathbb{R}^{(m+n-1)\times mn} \qquad B = \frac{\mathrm{d}^2}{\mathrm{d}M\partial P} f \in \mathbb{R}^{mn\times nn} \qquad H = \frac{\mathrm{d}^2}{\mathrm{d}P^2} f \in \mathbb{R}^{mn\times mn}$$

$$= \begin{bmatrix} \mathbf{0}_n^T & \mathbf{1}_n^T & \dots & \mathbf{0}_n^T \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_n^T & \mathbf{0}_n^T & \dots & \mathbf{1}_n^T \\ I_{n\times n} & I_{n\times n} & \dots & I_{n\times n} \end{bmatrix} \qquad = I_{mn\times mn} \qquad H_{ij,kl} = \begin{cases} \frac{1}{\gamma P_{ij}} & \text{if } ij = kl \\ 0 & \text{otherwise} \end{cases}$$

$$f(M,P) = \sum_{i=1}^{m} \sum_{j=1}^{n} M_{ij} P_{ij} + \frac{1}{\gamma} \sum_{i=1}^{m} \sum_{j=1}^{n} P_{ij} \log P_{ij} \qquad h(M,P) = \begin{bmatrix} \mathbf{0}_{n}^{T} & \mathbf{1}_{n}^{T} & \dots & \mathbf{0}_{n}^{T} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix} \vec{P} - \begin{bmatrix} r_{2} \\ \vdots \\ r_{m} \\ c \end{bmatrix}$$

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# Computing $(AH^{-1}A^T)^{-1}$ back

$$H^{-1} = \gamma \operatorname{diag}(\vec{P}) \qquad A = \begin{bmatrix} \mathbf{0}_n^T & \mathbf{1}_n^T & \dots & \mathbf{0}_n^T \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_n^T & \mathbf{0}_n^T & \dots & \mathbf{1}_n^T \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix}$$

$$\frac{\mathrm{d}P}{\mathrm{d}M} = \left(H^{-1}A^T \left(AH^{-1}A^T\right)^{-1}AH^{-1} - H^{-1}\right)B$$

## Computing $(AH^{-1}A^T)^{-1}$

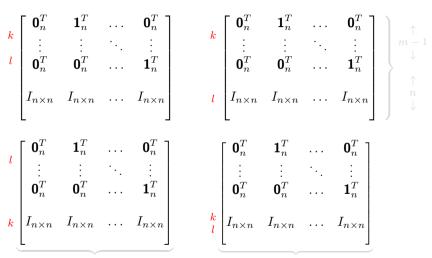
$$H^{-1} = \gamma \operatorname{diag}(\vec{P}) \qquad A = \begin{bmatrix} \mathbf{0}_{n}^{T} & \mathbf{1}_{n}^{T} & \dots & \mathbf{0}_{n}^{T} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_{n}^{T} & \mathbf{0}_{n}^{T} & \dots & \mathbf{1}_{n}^{T} \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix}$$

$$\frac{\mathrm{d}P}{\mathrm{d}M} = \left(H^{-1}A^T \left(AH^{-1}A^T\right)^{-1}AH^{-1} - H^{-1}\right)B$$

The (k, l)-th entry of  $AH^{-1}A^T$  for  $k, l \in 1, ..., m+n-1$  is

$$(AH^{-1}A^{T})_{kl} = \sum_{i=1}^{m} \sum_{j=1}^{n} \frac{A_{k,ij}A_{l,ij}}{H_{ij,ij}} = \gamma \sum_{i=1}^{m} \sum_{j=1}^{n} A_{k,ij}A_{l,ij}P_{ij}$$

## Interpreting $A_{k,ij}A_{l,ij}$ back



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# Evaluating $(AH^{-1}A^T)_{kl}=\gamma\sum_{i=1}^m\sum_{j=1}^nA_{k,ij}A_{l,ij}P_{ij}$ where

$$0 \leq l \leq m-1 \qquad m \leq l \leq m+n-1$$
 
$$0 \leq k \leq m-1 \qquad \begin{cases} \gamma \sum_{j=1}^{n} P_{k+1,j} & \text{if } k=l \\ 0 & \text{otherwise} \end{cases} \qquad \gamma P_{k+1,l-m+1}$$
 
$$m \leq k \leq m+n-1 \qquad \gamma P_{l+1,k-m+1} \qquad \begin{cases} \gamma \sum_{i=1}^{m} P_{i,k-m+1} & \text{if } k=l \\ 0 & \text{otherwise} \end{cases}$$

# Computing $(AH^{-1}A^T)^{-1}$ back

$$H^{-1} = \gamma \operatorname{diag}\left(\vec{P}\right)$$
 
$$A = \begin{bmatrix} \mathbf{0}_n^T & \mathbf{1}_n^T & \dots & \mathbf{0}_n^T \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}_n^T & \mathbf{0}_n^T & \dots & \mathbf{1}_n^T \\ I_{n \times n} & I_{n \times n} & \dots & I_{n \times n} \end{bmatrix}$$

$$\frac{\mathrm{d}P}{\mathrm{d}M} = \left(H^{-1}A^T \left(AH^{-1}A^T\right)^{-1}AH^{-1} - H^{-1}\right)B$$

$$AH^{-1}\!A^T = \gamma \begin{bmatrix} \operatorname{diag}(r_{2:m}) & P_{2:m,1:n} \\ P_{2:m,1:n}^T & \operatorname{diag}(c) \end{bmatrix} \qquad \left(AH^{-1}\!A^T\right)^{-1} = \frac{1}{\gamma} \begin{bmatrix} \Lambda_{11} & \Lambda_{12} \\ \Lambda_{12}^T & \Lambda_{22} \end{bmatrix}$$

$$\begin{split} &\Lambda_{11} = \left( \mathsf{diag} \Big( r_{2:m} - P_{2:m,1:n} \mathsf{diag}(c)^{-1} \, P_{2:m,1:n}^T \Big) \right)^{-1} \\ &\Lambda_{12} = -\Lambda_{11} P_{2:m,1:n} \mathsf{diag}(c)^{-1} \\ &\Lambda_{22} = \mathsf{diag}(c)^{-1} - \mathsf{diag}(c)^{-1} \, P_{2:m,1:n}^T \Lambda_{12} \end{split}$$

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#### end

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