From now on: non-binary field k; non-square  $d \in k$ .

$$E(k) = \{(x, y) \in k \times k : x^2 + y^2 = 1 + dx^2y^2\}$$

is a commutative group with  $(x_1, y_1) + (x_2, y_2) = (x_3, y_3)$  defined by Edwards addition law:

$$x_3 = \frac{x_1y_2 + y_1x_2}{1 + dx_1x_2y_1y_2},$$

$$y_3 = \frac{y_1y_2 - x_1x_2}{1 - dx_1x_2y_1y_2}.$$

Birationally equivalent to  $(1/e)v^2 = u^3 + (4/e - 2)u^2 + u$  where e = 1 - d.

Represent  $(x, y) \in E(k)$ by  $(X : Y : Z) \in \mathbf{P}^{2}(k)$ ; i.e.,  $(X, Y, Z) \in k^{3}$  with  $Z \neq 0$ and  $(X^{2} + Y^{2})Z^{2} = Z^{4} + dX^{2}Y^{2}$ represents  $(X/Z, Y/Z) \in E(k)$ .

10**M** (10 field mults) + 1**S** (1 field squaring) + 1**D** (1 field mult by d) + 7add (7 field additions) to obtain sum  $(X_3:Y_3:Z_3)$ of  $(X_1:Y_1:Z_1)$ ,  $(X_2:Y_2:Z_2)$ .

Don't have to make distinctions for equal inputs, negatives, etc.

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$$= \left(\frac{(x + y)^2}{x^2 + y^2} - 1, \frac{y^2 - x^2}{2 - x^2 - y^2}\right)$$

3M (3 field mults) + 4S (4 field squarings) + 6add (6 field additions) to double  $(X_1 : Y_1 : Z_1)$ :

$$B = (X_1 + Y_1)^2,$$
  
 $C = X_1^2,$   
 $D = Y_1^2,$   
 $E = C + D,$   
 $H = Z_1^2,$   
 $J = E - 2H,$   
 $X_3 = (B - E)J,$   
 $Y_3 = E(C - D),$   
 $Z_3 = EJ.$ 

Comparison of doubling costs if curve parameters are small:

System	Cost
Projective	5M + 6S
Projective if $a = -3$	7M + 3S
Hessian	7M + 1S
Doche/Icart/Kohel 3	2M + 7S
Jacobian	1M + 8S
Jacobian if $a = -3$	3M + 5S
Jacobi quartic	2M + 6S
Jacobi intersection	3M + 4S
Edwards	3M + 4S
Doche/Icart/Kohel 2	2M + 5S

EFD! EFD! EFD! EFD! e.g. Doche/Icart/Kohel paper says 3M+4S for Doche/Icart/Kohel 2.

Jacobian a = -3 vs. Edwards:

	Jac-3	Edwards
Double	3 <b>M</b> +5 <b>S</b>	3 <b>M</b> +4 <b>S</b>
Triple	7M+7S	9M+4S
Add	11M+5S	10M+1S+1D
Readd	10M+4S	10M+1S+1D
Mixed	7M+4S	9M + 1S + 1D
Unified	unclear	$10\mathbf{M} + 1\mathbf{S} + 1\mathbf{D}$

Jac-3 speedup for readd: Chudnovsky/Chudnovsky 1986; "Chudnovsky coordinates" etc.

Edwards tripling:

Bernstein/Birkner/Lange/Peters 2007; independently Hisil/Carter/Dawson 2007.

A sensible ElGamal-type system (van Duin, sci.crypt, 2006):

Everyone knows standard point B, prime order q, on "Curve25519":  $\mathbf{Z}/(2^{255}-19)$ ; d=1-1/121666.

Signer has 32-byte secret key *n*. Everyone knows signer's 32-byte public key: compressed *nB*.

To verify (m, compressed R, t): verify tB = H(R, m)R + nB.

To sign m: generate a secret s; R = sB;  $t = H(R, m)s+n \mod q$ .

Notes: 1. No inversions mod q. 2. Send R, not H(R, m).

Batch verification of many  $t_i B - h_i R_i = S_i$ : check  $\sum_i v_i t_i B - \sum_i v_i h_i R_i - \sum_i v_i S_i = 0$  for random 128-bit  $v_i$ . (Naccache et al., Eurocrypt 1994; Bellare et al., Eurocrypt 1998)

Use subtractive multi-scalar

multiplication algorithm: if  $n_1 \ge n_2 \ge \cdots$  then  $n_1P_1 + n_2P_2 + n_3P_3 + \cdots = (n_1 - qn_2)P_1 + n_2(qP_1 + P_2) + n_3P_3 + \cdots$  where  $q = \lfloor n_1/n_2 \rfloor$ . (credited to Bos and Coster by de Rooij, Eurocrypt 1994; see also tweaks by Wei Dai, 2007)

Verifying 100 signatures requires a 201-scalar mult with 101 256-bit scalars and 100 128-bit scalars.

Subtractive algorithm then uses

- $pprox 24.4 \cdot 256$  readds and
- $\approx 0.8 \cdot 256$  mixed adds.

S/M = 0.8, small parameters:

- $\approx 845 M/signature$  with Jacobian;
- $\approx 695 M/signature$  with Edwards.

## **Use Edwards coordinates!**

Can similar speeds be achieved by genus-2 hyperelliptic curves? Current attempts seem very slow.

We've counted mults (with various **S**/**M**, **D**/**M**) for Edwards, Jac-3, Hessian, et al. in NAF; width-4 sliding windows; JSF; accelerated ECDSA; batch verification, as above; fixed-point scalar mult; and several side-channel situations.

Edwards consistently wins!

Should even beat Montgomery for big single-scalar mult.

Need to measure overheads too. Planning new Edwards software. Expect new speed records.

Dimitrov/Imbert/Mishra 2005, Doche/Imbert 2006:

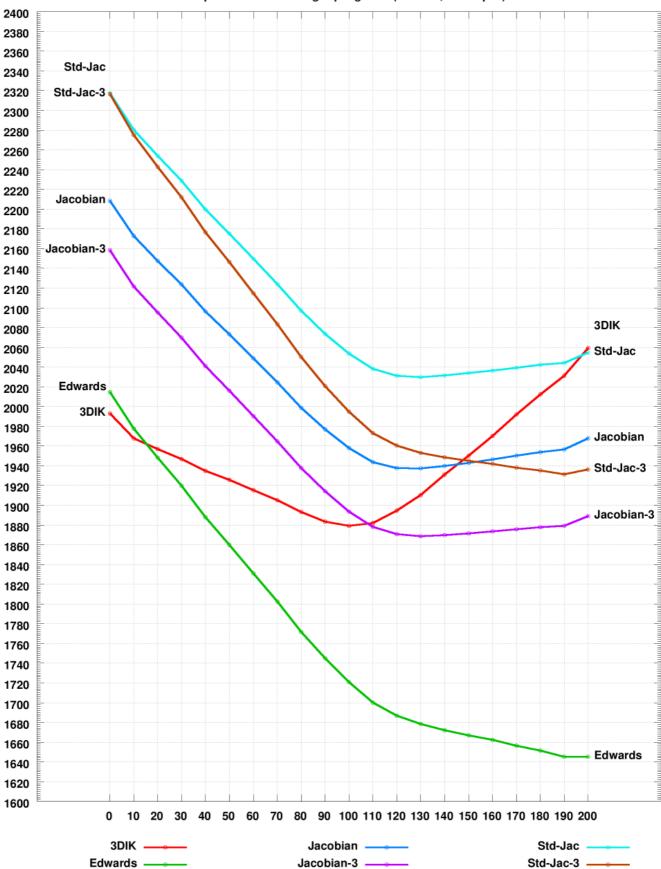
Mix doublings with triplings to gain speed for single-scalar mult.

Bernstein/Birkner/Lange/Peters 2007: Have analyzed Edwards, Jac-3, et al. with 5423 combinations of bit size, doubling/tripling ratio, windowing strategy.

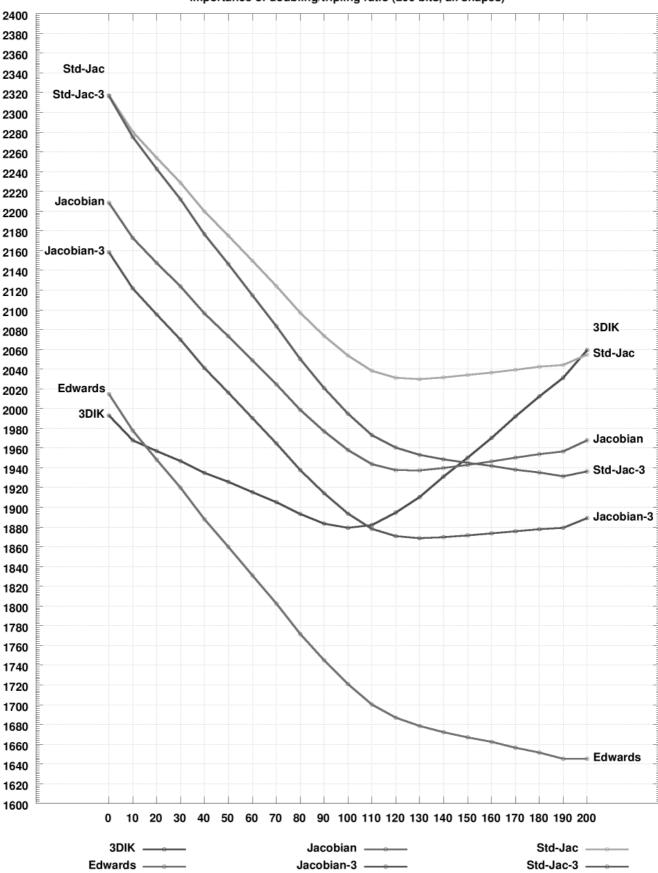
Planning more combinations.

Conclusions: Triplings *are* useful for Jac-3, 3DIK, et al. But Edwards wins solidly.

Importance of doubling/tripling ratio (200 bits, all shapes)



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Edwards for precomputation!

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And beyond ECC:

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And beyond ECC:

Edwards for ECM!

Edwards for ECPP!

Edwards for ECXYZ!

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And beyond ECC:

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Return of the Hyperelliptic!

